

Juhani Nieminen

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OPERATIONS ON GRAPHS DETERMINING CONGRUENCES ON GRAPHS

JUHANI NIEMINEN, Oulu

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The purpose of this paper is to characterize by means of concepts and operations of graph theory partitions of the elements of a finite modular lattice H that determine congruence relations on H . By the aid of the characterization we construct thereafter a class of congruence relations on graphs. We recall first some concepts of graph theory and apply thereafter them to the Hasse diagram of H in order to obtain the characterization.

We shall consider finite undirected and connected graphs $G = (P(G), L(G))$ only without loops and multiple lines, where $P(G)$ is the set of points of G and $L(G)$ its set of lines. SP is a mapping $P(G) \times P(G) \rightarrow 2^{P(G)}$ defined as follows:

$$SP(x, y) = \{z \mid z \in P(G) \text{ and } z \text{ is on a shortest path joining } x \text{ and } y \text{ in } G\}.$$

We shall call SP a binary operation on $P(G)$, although the mapping induced by the operation is a one-to-many mapping, as the name operation helps us to find some useful analogies we shall apply. In particular, $\{x, y\} \subseteq SP(x, y)$ and $SP(x, x) = \{x\}$, $x, y \in P(G)$. In general, let U and W be two subsets of $P(G)$, then $SP(U, W)$ denotes the union of the sets $SP(u, w)$, where $u \in U$ and $w \in W$; formally $SP(U, W) = \{z \mid z \in SP(u, w) \text{ for some } u \text{ and } w, u \in U \text{ and } w \in W\}$. A set $U \subset P(G)$ is called an ideal of G , if $U \neq \emptyset$ and $SP(U, U) = U$. By the notation $SP^n(x, y)$ we denote the operation $SP(SP^{n-1}(x, y), SP^{n-1}(x, y))$. Thus $SP^2(x, y) = SP(SP(x, y), SP(x, y))$. As we consider finite graphs only, there is for any pair $x, y \in P(G)$ a value of n such that $SP^n(x, y)$ is an ideal of G . The graph of Figure 1 illuminates the case where $SP^2(x, y)$ is not an ideal of G but $SP^3(x, y)$ is. It is important to construct from a pair $x, y \in P(G)$ an ideal of G by means of sequential applying of the SP -operation and in order to use a brief notation, $SU(x, y)$ denotes the ideal obtained from x, y by applying the SP -operation enough many times.

Ideals of graphs and the SP -operation were introduced in [4] and briefly considered in [5]. These concepts are natural generalizations of corresponding concepts defined for trees by Nebeský in [3].

In this paper we consider the Hasse diagram of a lattice H as an undirected graph and denote it by G_H . Lemma 1 and Theorem 1 are proved in a more general form than we need later. A lattice H is locally finite, if its every interval is finite.

Lemma 1. *Let H be a locally finite lattice. Then $SU(x, y) = [x \wedge y, x \vee y]$ for any two elements $x, y \in H$ if and only if H is modular.*

Proof. If H is modular, then according to the metric properties of finite modular lattices, $x \wedge y, x \vee y \in SP(x, y)$ (see e.g. Draškovičová [1]). As the lengths of any two chains between a and b in a finite modular lattice are equivalent when $a < b$, each $z \in [x \wedge y, x \vee y]$ belongs to a shortest path from $x \wedge y$ to $x \vee y$ and so $z \in SU(x, y)$. Obviously $SU(x, y) \subseteq [x \wedge y, x \vee y]$, and thus $SU(x, y) = [x \wedge y, x \vee y]$.

Let H satisfy the condition of the lemma for any pair $x, y \in H$. If H were non-modular, then it contains the well known non-modular sublattice (in Figure 2 the sublattice of elements a, b, c, d, e), where the set $\{a, b, c\} = SU(b, c) \neq [a, e] = [b \wedge c, b \vee c]$. This completes the proof.

Now we are ready to prove the characterization.

Theorem 1. *Let H be a locally finite modular lattice and $\mathfrak{C} = \{C_1, \dots, C_m\}$ a partition of its elements. \mathfrak{C} is a congruence partition of H with respect to the operations \vee and \wedge on H if and only if the condition (A) holds.*

(A) *If $x, y \in C_i$ and $a, b \in C_j$ in \mathfrak{C} , then $SU(x, a) \cap C_k \neq \emptyset$ holds for some k in G_H if and only if $SU(y, b) \cap C_k \neq \emptyset$ holds, $1 \leq k \leq m$.*

Proof. Assume that \mathfrak{C} is a partition of the points $P(G_H)$ such that $SU(x, a) \cap C_k \neq \emptyset$ if and only if $SU(y, b) \cap C_k \neq \emptyset$. We show that R is a latticecongruence on H , with the classes C_1, \dots, C_m . Clearly R is reflexive, symmetric and transitive. Thus it remains to show the compatibility of R , i.e. to show that xRy implies $x \wedge zRy \wedge z$ and $x \vee zRy \vee z$ for any $z \in H$. Moreover, if $qRp \Leftrightarrow q \wedge pRq \vee p$, we may assume that $x \leq y$.

Let $x < y$ (the case $x = y$ is trivial), xRy and $z \in H$. Thus $x \vee z \leq y \vee z$. We assume that in the partition \mathfrak{C} of $Hx \vee z$ and $y \vee z$ belong to different sets of \mathfrak{C} . As $x \leq y \wedge (x \vee z) \leq y$ and $yRx, y \wedge (x \vee z)Ry$ holds, too. The relations $y \vee zRy \vee z$ and yRx imply that (A) holds for $SU(y \vee z, y)$ and $SU(y \vee z, x)$. $x \vee z \in SU(y \vee z, x)$ and we assume that $x \vee z \in C_h$. As $x \vee z < y \vee z, x \vee z \notin SU(y \vee z, y)$. Then according to (A), $SU(y \vee z, y) \cap C_h \neq \emptyset$, and let t be the greatest element of the set $SU(y \vee z, y) \cap C_h$; such an element exists as $SU(y \vee z, y)$ is finite and for any two elements of $SU(y \vee z, y)$ (of C_h), $SU(y \vee z, y) \cap C_h$ contains the join of these elements. But $x \vee z \vee t \in C_h$ and $x \vee z \vee t \leq y \vee z$, whence $x \vee z \vee t \in SU(y \vee z, y)$. Thus we can assume that $x \vee z \leq t$, and as $t \in SU(y \vee z, y)$, $t \geq y$. But then $y \vee x \vee z = y \vee z \leq t$, whence $y \vee z, x \vee z \in C_h$, which is a contradiction. Hence $y \vee z, x \vee z \in C_k$ for some value k of i . The proof is similar for $y \wedge zRx \wedge z$.

Conversely, we assume that \mathfrak{C} generates a latticecongruence on H . Let $x, y \in C_i$,

$a, b \in C_j$ and $i \neq j$. Accordingly, we may assume that $x \leq y$ and $a \leq b$. As R is a congruence relation on H , $y \vee bRx \vee a$. If there is an element $q \in C_k$, $y \leq q \leq \leq y \vee b$, then $x \leq q \wedge (x \vee a) \leq x \vee a$, $q \wedge (x \vee a) Rq$ and thus $q \wedge (x \vee a) \in C_j$. By applying this technique to the intervals $[y \wedge b, y \vee b]$ and $[x \vee a, x \wedge a]$ we see that the condition (A) holds for $SU(y, b)$ and $SU(x, a)$. The proof is similar for $SU(x, b)$ and $SU(y, a)$. This completes the proof.

As the example of Figure 2 shows, a partition of a non-modular lattice H satisfying the condition (A) need not be either a \wedge -congruence or a \vee -congruence on H .

In the next theorem we show how the condition (A) generalizes by a natural way the construction of compatible tolerances on graphs introduced by Zelinka in [7].

We call a binary, reflexive, symmetric and transitive relation R on a graph a SU -compatible congruence relation on G when aRb and xRy imply $SU(a, x)RSU(b, y)$. The notation $SU(a, x)RSU(b, y)$ means that for any $z \in SU(a, x)$ there is a point $u \in SU(b, y)$ such that zRu , and for any $w \in SU(b, y)$ there is a point $v \in SU(a, x)$ such that vRw .

Theorem 2. *Let \mathfrak{C} be a partition of the pointset $P(G)$ of a graph $G = (P(G), L(G))$. The relation R given by \mathfrak{C} determines a SU -compatible congruence relation on G if and only if \mathfrak{C} satisfies the condition (A).*

Proof. If \mathfrak{C} is a partition of $P(G)$ such that the relation R given by \mathfrak{C} satisfies the condition (A), the SU -compatibility of R follows directly from (A). The converse proof follows similarly directly from the definition of the SU -compatibility.

By using the terminology of Theorem 2, we can say, according to Theorem 1, that R is a latticecongruence on a finite modular lattice H if and only if R is a SU -compatible congruence relation on G_H .

We obtain also a characterization of finite modular lattices as given in the next theorem.

Theorem 3. *Let H be a finite lattice and \mathfrak{C} a partition of H determining a SU -compatible relation R on G_H . H is modular if and only if each R defined above is a latticecongruence on H .*

Proof. If H is modular, then the assertion follows from Theorems 1 and 2. Thus let each R of the theorem be a congruence relation on H . If H is non-modular, it contains as a sublattice the lattice of the elements a, b, c, d, e in Figure 2, where the subset $\{a, b, c\}$ of the partition $\mathfrak{C} = \{\{a, b, c\}, \{d, e\}\}$ shows that \mathfrak{C} does not determine a congruence relation on H although R is SU -compatible on G_H .

As a model for constructing a SU -compatible congruence on G were latticecongruences on a finite modular lattice H . This model is used in the following theorem where an analogy is presented between SU -compatible congruences on G and congruences on algebras. Its proof is a direct copy of the corresponding proof for algebras given e.g. in [6, Thm. 96 and its supplement], and hence we omit it.

Theorem 4. Let G be a given graph. G is a Cartesian product of two non-trivial graphs G_1 and G_2 , i.e. $G = G_1 \times G_2$, if and only if there are two non-trivial SU-compatible congruences $R_1, R_2 \in H(G)$ which are permutable and complements of each other in $H(G)$. $H(G)$ is the lattice of SU-compatible congruences on G .

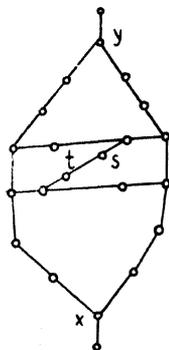


Fig. 1

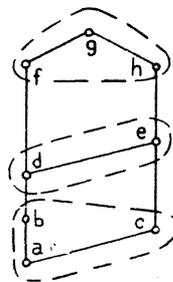


Fig. 2

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Juhani Nieminen
University of Oulu
Dept. Math. in Faculty of Technology
90101 OULU 10
Finland