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# CUBIC MOORE GR APHS 

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By a tied graph of type ( $d, k$ ) we understand - in accordance with [1] a regular graph with a (finite or infinite) degree $d$ and with a finite diameter $k$, not containing any circuit of length $\leqslant 2 k$. Finite tied graphs (i. e., tied graphs of finite degree - so-called Moore graphs) were studied in [1], [2], [3]. In the present paper except in the last $\S 4-$ we shall consider only tied graphs of type ( $3, k$ ), that is cubic Moore graphs. Obviously, there is no Moore graph of type $(3,0)$ and there exists up to isomorphism exactly one Moore graph of type (3,1) (tetrahedron). It is known [2] that there exists up to iso morphism just one Moore graph of type (3,2) (the Petersen graph) and no Moore graph of type ( 3,3 ). In this paper we prove the non-existence of Moore graphs of type ( $3, k$ ), where $3 \leqslant k \leqslant 8$. (1) For $k \geqslant 9$ the question of the existence of Moore graphs of type ( $3, k$ ) remains open. In $\S 4$ we give a survey of known results on the existence and the uniqueness of tied graphs of a given type.

## § 1. BASIC PIROPER'TIES OF CUBIC MOORE GRAPHN

Let $G_{k}$ be a Moore graph of type $(3, k)$ where $k \geqslant 3$. Pick a vertex $w$ of $G_{k}$ As $G_{k}$ is a cubic graph, $w$ is adjacent to three vertices $a, b$ and $c$ of $G_{k}$ (Fig. 1) The distance of vertices $x$ and $y$ in $G_{k}$ will be denoted by $r(x, y)$. Vertices $x$ such that $r(x, w) \quad k$, will be called $w$-vertices of $G_{k}$, edges joining such ver tices - $w$-edges of $G_{k}$. As $r(x, w)-k$, the vertex $x$ is adjacent to a vertex $y$ such that $r(y, w) \quad k \quad 1$. Considering the fact that $G_{k}$ does not contain any circuit of length $\leqslant 2 k$, the remaining two vertices, adjacent to $x$, are $w$-vertices Therefore the $w$-vertices and the $w$-edges form a quadratic subgraph of $G_{h}$, the circuits of which it consists are called $w$-circuits of $G_{k}$. Evidently, $G_{k}$ con tains exactly $3.2^{k}{ }^{1} w$-vertices and the same number of $w$-edges. Further, $G_{k}$ has
${ }^{(1)}$ This iesult was presented at the Colloquium on Graph Theory in Manebach (G.D R.) in May 1967.

$$
1+\sum_{i 1}^{k} 3.2^{k}-3.2^{k}-2
$$

vertices and

$$
\frac{3}{2}\left(3.2^{k}-2\right) \quad 3\left(3.2^{k} 1 \quad 1\right)
$$

edges. If we omit all $w$-edges from $G_{k}$, the graph $T(w)$ obtained in this way will be also connected (from every vertex there is a path to $w$ in $T(w)$ ). As $T(w)$ has $3.2^{k} 2$ vertices and $3\left(3.2^{k} 1 \quad\right.$ l) $3.2^{k}{ }^{1}-3.2^{k}-3$ edges, $T(w)$ is a tree, namely a spanning tree of $G_{k}$. The symbol $r_{w}(x, y)$ denotes the distance of vertices $x$ and $y$ in $T(w)$ and the symbol $x,:, y$ denotes the unique path comnecting $x$ and $y$ in $T(w)$. Obviously, $r_{w}(x, y) \geqslant r(x, y)$ and $r(w, x) \quad r_{u}(w, x)$.

Fig. 1


Suppose $r(x, w) \geqslant 2$. Evidently, there exists a unique vertex $y$ such that $r(x, w) \quad r(y, w)$ and $r_{w}(x, y)-2$. This vertex will be denoted by $y \quad \alpha x$ Obviously, $\alpha^{2} x \quad x$.

Let $i \in\{0,1,2, \ldots, k\}$ and $r(w, x) \geqslant i$. Then there is exactly one vertex y for which $r(w, y) \quad i, r(w, y)+r(y, x) \quad r(w, x)$. It will be denoted by $y \quad \beta, x$ Instead of $\beta_{3}$ we shall write briefly $\beta$. Evidently, if $k \geqslant 4$ and $x$ is a $w$-vertex of $G_{k}$, then $\beta \alpha x \quad \beta x . G_{k}$ is a tied graph, therefore it contains no multiple edges. Thus we may denote the edge joining vertices $x$ and $y$ by ( $x, y$ ) and the path with vertices $x_{1}, x_{2}, \ldots, x_{n}$ by $\left[x_{1}, x_{2}, \ldots, x_{n}\right]$.

Lemma 1. Let $\left(x, y_{1}\right)$ and $\left(x, y_{2}\right)$, where $y_{1} \neq y_{2}$, are $w$-edges of $G_{k}$. Then thr following equality of sets holds:

$$
\left\{\beta_{1} x, \beta_{1} y_{1}, \beta_{1} y_{2}\right\} \quad\{a, b, c\}
$$

Proof. Evidently, each of the elements $\beta_{1} x, \beta_{1} y_{1}, \beta_{1} y_{2}$ belongs to the set $\{a, b, c\}$. If the assertion of Lemma 1 were false, two of elements $\beta_{1} x, \beta_{1} y_{1}$, $\beta_{1} y_{2}$ would coincide. If $\beta_{1} x \quad \beta_{1} y_{i}(i \in\{1,2\})$, there exists in $G_{k}$ a circuit $\left\lfloor x,:, y_{i}, x\right\rfloor$ of length $\leqslant 2 k \quad$, which is in contradiction to the definition of a tied graph. If $\beta_{1} y_{1} \quad \beta_{1} y_{2}$, there exists in $G_{k}$ a circuit $\left[y_{1},:, y_{2}, x, y_{1}\right]$ of length $\leqslant 2 k$, a contradiction again. The lemma follows.

Now we can assign to every $w$-vertex $x$ of $G_{k}$ a $w$-vertex !/ $q x$ adjacent to $x$ so that

$$
\begin{aligned}
& \text { if } \beta_{1} x-a \text {, then } \beta_{1} y \quad b, \\
& \text { if } \beta_{1} x=b \text {, then } \beta_{1} y=c \text {, } \\
& \text { if } \beta_{1} x \quad c \text {, then } \beta_{1} y \quad a \text {. }
\end{aligned}
$$

Lemma 1 guarantces the existence and uniqueness of $q x$.

## § .. AUNILIARY RESULTs

Henceforth we shall use notation introduced in § 1 .
Lemma 2. Let $x$ be a w-verter of $G_{k}$. We have:
(a) $\beta_{1} \varphi^{i} x \quad \beta_{1} \varphi^{j} x$ if and only if $i \quad j(\bmod 3)$.
(b) The elements $\beta x, \alpha \beta x, \beta \varphi^{3} x, \alpha \beta \phi^{3} x$ are mutually different.
(c) $\beta q^{6} x \quad \alpha, \beta x, \beta q^{9} x \quad \alpha, \beta q^{3} x, \beta q^{12} x \quad \beta x$.
(d) The elements $\beta x, \beta q x, \beta q^{2} x, \beta q^{3} x, \ldots, \beta q^{11} x$ are mutually diffferent.

Proof. (a) follows from Lemma 1.
(b) From the definition of $\alpha$ it follows that $\beta x \neq \alpha \beta x$, and $\beta q^{3} x \neq \alpha_{1}{ }^{3} q^{3}, x$ If $\beta x \quad \beta \varphi^{3} x$, then there exists a circuit $\left[x, \varphi x, \varphi^{2} x, \varphi^{3} x,:, x\right]$ in $G_{k}$ of length $<2 k-3$, which is in contradiction to the definition of a tied graph. If $\beta q^{3} \cdot x$
$\alpha \beta x$, we have a circuit $\left[x, \varphi x, \varphi^{2} x, \gamma^{3} x,:, x\right]$ of length $2 k \quad 1$, a contradic tion again. If $\beta x \quad \alpha \beta q^{3} x$, then $\alpha \beta x-\alpha^{2} \beta q^{3} x \quad \beta p^{3} x$, and we have the case treated above. If $\alpha \beta x \quad \alpha \beta q^{3} x$, then $\alpha^{2} \beta x \quad \alpha^{2} \beta q^{3} x$, i. e. $\beta x \quad \beta \varphi^{3} x$, which is also impossible.
(c) According to (b) the elements $\beta x, \alpha \beta x, \beta q^{3} x, \alpha \beta q^{3} x$ are mutually different But from Lemma 1 it follows that $\beta_{1} \beta x-\beta_{1} \beta \varphi^{3} x \quad \beta_{1} \alpha \beta x \quad \beta_{1} \alpha \cdot \beta p^{3} x \quad \beta_{1} \beta_{q} q_{x}$ Thercfore $\beta q^{6} x \in\left\{\beta x, \beta \varphi^{3} x, \alpha, \beta x, \alpha \beta q^{3} x\right\}$. If $\beta q^{6} x \quad \beta x$, then a circuit $\mid x, q x$, $\varphi^{2} x, \varphi^{3} x, \varphi^{4} x, q^{5} x, q^{6} x,:, x \mid$ of length $\leqslant 2 k$ would exist in $G_{k}$, which is a con tradiction. If $\beta q^{6} x \quad \beta q^{3} x$, for $y \quad q^{3} x$ we should have $\beta q^{3} y \quad \beta y$, which contradicts (b). If $\beta \varphi^{6} x \quad \alpha \beta \varphi^{3} x$, then analogously we have $\beta q^{3}!y \quad \alpha \beta y$, again in contradiction to (b). Therefore $\beta \varphi^{6} x-\alpha \beta x$. Using this relation we obtain $\beta q^{9} x \quad \beta \varphi^{6}\left(q^{3} x\right) \quad \beta \varphi^{6} y \quad \alpha \beta y \quad \alpha \beta q^{3} x$. Further, $\beta \varphi^{12} x \quad \beta q^{6}\left(\varphi^{6} x\right) \quad \alpha \beta\left(q^{6} \cdot x^{\prime}\right)-$ $\alpha^{2} / \beta x \quad \beta x$.
(d) Let $\beta q^{i} x \quad \beta q^{j} x, i, j \in\{0,1, \varrho, \ldots, 1 \mathrm{l}\}, i \quad j$. Evidently, $\beta_{1 q^{i}} x \quad \beta_{1 q}{ }^{j} x$; according to (a), we have $i \quad j(\bmod 3)$, i. e. we can write $j \quad i \quad 3 t, t \in\{1, \geq, 3\}$. Put y $\quad \phi^{i} x$. We have: $\beta$ ! $\quad \beta \varphi^{i} x \quad \beta q^{j} x \quad \beta q^{i}{ }^{3 t} x \quad \beta q^{3 t} y$. But from (b) and (c) it follows that $\beta y \neq \beta \varphi^{3 t} y$, which is impossible.

Lemma 3. The length of every w-circuit of $G_{k}$ is a multiple of 1 ㅇ.
Proof follows from (c) and (d) of Lemma $\because$.

Lemma 4. Let $M$ be a set of $w$-vertices of $G_{k}, k \geqslant 5$. If $M$ has more than $2^{h-5}$ loments and for every $y_{1}, y_{2} \in M$ we have $\beta y_{1} \quad \beta y_{2}$, then there exist $x_{1}, x_{2} \in M$, ${ }_{1} \quad x_{2}$ such that $r_{w}\left(x_{1}, x_{2}\right) \leqslant 4$.

Proof. Form the set $N \quad\left\{\beta_{k}{ }_{2} x\right\}_{x \in M}$. The set $N$ evidently cannot have more than $2^{k}{ }^{5}$ elements; therefore for some $x_{1}, x_{2} \in M, x_{1} \neq x_{2}$ we have $i_{h} \quad 2 x_{1} \quad \beta_{h}{ }_{2} x_{2}$. i. e. $r_{w}\left(x_{1}, x_{2}\right) \leqslant 4$.

Lemma 5. Let $x$ and $y$ be w-vertices of $G_{k}$. If $\beta x \quad \beta y$, then $r_{\mu}(x, y) \leqslant 2 k \quad 6$ If $\beta, x \quad \alpha \beta y$, then $n r_{u}(x, y) \quad 2 k \quad 4$.

Proof. The path $[x,:, y\rceil$ has evidently the length $\leqslant 2 k \quad 6$ in the first dse and the length $2 k \quad 4$ in the second case.

Lemma 6. If $x \neq y$ are such w-vertices of $G_{k}$ that $\beta x \quad \beta y$ and $\beta q x \quad \beta \varphi y$, the $r_{w}(x, y)>6$.

Proof. If the assertion of the lemma were not truc, then $r_{u}(x, y)<4$ B. Lemma 5 we have $r_{w}(\varphi x, \varphi y) \leqslant 2 k-6$. But then $\lfloor x, q x,:, q!, y,:, x\rfloor$ would be a circuit of length $\leqslant 2 k$, which is impossible.

Lemma 7. Let $x$ bc a w-vertex of $G_{k}, k \geqslant 4$. Then we have.

$$
\begin{array}{ll}
\beta \varphi^{2} \alpha x & \alpha \beta q x, \\
\beta \phi^{1} \alpha x & \alpha \beta \gamma^{2} x, \\
\beta \varphi \alpha x & \alpha \beta q{ }^{2} x, \\
\beta \phi^{2} \alpha x & \alpha \beta \gamma^{1} x . \tag{4}
\end{array}
$$

Proof. First we prove (3). As $\beta x \quad \beta \alpha x$, consequently $\beta_{1} x \quad \beta_{1} \alpha x$, and dso $\beta_{1} \varphi x \quad \beta_{1} q \alpha x$. According to (d) of Lemma 2 the elements $\beta(q \alpha x), \beta \varphi^{3}(q \alpha x)$ $\beta q^{\prime \prime}(q \alpha x), \beta q^{9}(p \alpha x)$ are mutually different. By (a) of Lemma 2 we have $\beta_{1}(q \propto x)$
$\beta_{1} q^{3}\left(q \alpha x^{x}\right) \quad \beta_{1} q^{6}(q \alpha x) \quad \beta_{1} \varphi^{9}(q \alpha x)$. Since $\beta_{1}(q \alpha x) \quad \beta_{1}(q x)$. the element $\beta(q \alpha x)$ equals one of the elements $\beta(q x), \beta q^{3}(q x) \quad \beta p^{6} q{ }^{2} x, \beta q^{6}(q x), \beta q^{9}(q x)$
$\beta q^{12} q^{2} x$, hence with respect to (c) of Lemma ${ }^{2} \quad \beta(q \alpha x)$ is equal to ,ome of the elements $\beta \varphi x, \alpha \beta \varphi{ }^{2} x, \alpha \beta q x, \beta q{ }^{2} x$.

If $\beta q \alpha, x \quad \beta q x$, then the circuit $\lfloor q x, x,:, \alpha x, q \alpha x,:, q x \mid$ has the length $\leqslant \varphi k$
2 . because $r_{\mu}\left(x, \alpha_{x} x\right) \quad 2$ and according to Lemma $\sigma_{r}(q \alpha x, q x) \leqslant 2 k \quad 6$. If fof $\alpha x \quad \alpha \beta q x$, the circuit $[\varphi x, x,:, \alpha x, \varphi \alpha x,:, \psi x]$ has the length $2 k$, for Lemma ; vields $r_{u}(q \alpha x, \varphi x) \quad 2 k \quad 4$. If $\beta q \alpha x \quad \beta \varphi^{2} x$, the circuit [ $\varphi^{2} x$, ${ }_{q}{ }^{1}, x, x,:, \alpha x^{\prime}, q_{x} x,:, \phi^{2} x \mid$ has the length $\leqslant 2 k \quad$ I, because Lemma 5 implies , ( $\left.\boldsymbol{q}^{2} x, q \alpha x\right)<2 k$ 6. Therefore only the last possibility, i. e. (3), can be valid.

The proof of (2) is ,,dual" to that of (3) it is sufficient to replace $q^{2}, \gamma, q{ }^{1}$ ind $q{ }^{2}$ by $q^{2}, q{ }^{1}, q$ and $q^{2}$, reepectively.

If in (3) we replace $x$ by $\alpha x$, we obtain

$$
\beta ; \alpha^{2} \cdot x \quad \alpha, \beta \eta \quad{ }^{2} \alpha x,
$$

whence, as $\alpha^{2}$ is an identical mapping, it follows that

$$
\beta \varphi^{2} \alpha x=\alpha^{2} \beta \varphi^{-2} \alpha x=\alpha \beta p \alpha^{2} x=\alpha \beta \varphi x,
$$

that is, the relation (1).
The proof of (4) is ,,dual" to that of (1).
Lemma 8. Let $x$ be a w-vertex of $G_{k}$, where $k \geqslant 4$. Then we have:

$$
\begin{aligned}
& \beta q^{4} \alpha x=\beta \varphi x, \\
& \beta \varphi^{5} \alpha x=\beta \varphi^{2} x, \\
& \beta \varphi^{6} \alpha x=\alpha \beta x, \\
& \beta \varphi^{7} \alpha x=\beta \varphi^{2} x, \\
& \beta \varphi^{8} \alpha x=\beta \varphi^{1} x, \\
& \beta \varphi^{10} \alpha x=\alpha \beta \varphi x \\
& \beta \varphi^{11} \alpha x=\alpha \beta \varphi^{2} x, \\
& \beta \varphi^{12} \alpha x=\beta x \\
& \beta \varphi^{13} \alpha x=\alpha \beta \varphi{ }^{2} x .
\end{aligned}
$$

The proof follows from (c) of Lemma 2 and Lemma 7, for instance:

$$
\begin{gathered}
\beta \varphi^{4} \alpha x=\beta \varphi^{6}\left(\varphi^{2} \alpha x\right)-\alpha \beta\left(\varphi^{2} \alpha x\right)=\alpha\left(\beta \varphi^{2} \alpha x\right) \quad \alpha(\alpha \beta \varphi x) \quad \beta \varphi x, \\
\beta \varphi^{5} \alpha x \quad \beta \varphi^{6}\left(\varphi^{1} \alpha x\right)=\alpha \beta\left(\varphi^{1} \alpha x\right)-\alpha\left(\beta \varphi^{1} \alpha x\right) \quad \beta \varphi^{2} x, \\
\beta \varphi^{6}(\alpha x)-\alpha \beta(\alpha x)-\alpha \beta x, \text { etc. }
\end{gathered}
$$

## § 3. MAIN RESULTS

Lemma 9. There is no Moore graph of type (3, 3). ( ${ }^{2}$ )
Proof. Let $\mathrm{G}_{3}$ be a Moore graph of type (3,3). Then for any $w$-vertex $x$ of $\mathrm{G}_{3}$ we have $\beta x \quad x$. (c) of Lemma 2 yields $\alpha x-\alpha \beta x \quad \beta \psi^{6} x \quad \varphi^{6} x, \alpha \varphi x \quad \alpha \beta(\varphi x)$
$\beta \varphi^{6}(\varphi x)-\varphi^{7} x$. Therefore $\mathrm{G}_{3}$ contains a hexagon $\left[x, q x,:, \varphi^{7} x, \varphi^{6} x,:, x \mid\right.$ which contradicts the definition of a Moore graph.

Lemma 10. There is no Moore graph of type (3, 4).
Proof. Let $G_{4}$ be a Moore graph of type (3,4). Let $x$ be a $w$-vertex in $\boldsymbol{G}_{4}$ Evidently $G_{4}$ has just $24 w$-vertices, so that, according to Lemma 6, in $G_{1}$ there is either one single $w$-circuit with 24 vertices or two $w$-circuits, each with 12 vertices. In the first case $G_{4}$ contains a hexagon $\left[x, \varphi x,:, \varphi^{13} x, \varphi^{12} x,:, x\right]$. and we have a contradiction. In the second case from (c) of Lemma 2 and Lemma 7 it follows that
${ }^{(2)}$ This result follows also from [2].

$$
\begin{aligned}
& \beta \varphi^{8} x=\beta \varphi^{6}\left(\varphi^{2} x\right)-\alpha \beta \varphi^{2} x=\beta \varphi^{-1} \alpha x, \\
& \beta \varphi^{7} x-\beta \varphi^{6}(\varphi x)-\alpha \beta \varphi x=\beta \varphi^{2} \alpha x,
\end{aligned}
$$

therefore $G_{4}$ contains a hexagon $\left[\varphi^{7} x, \varphi^{8} x,:, \varphi^{1} \alpha x, \varphi^{-2} \alpha x,:, \varphi^{7} x\right]$, thus we have arrived at a contradiction again.

Lemma 11. The length of any $w$-circuit in $G_{k}(k \geqslant 5)$ is at most $3.2^{k}{ }^{5}$.
Proof. Let $C$ be a $w$-circuit in $G_{k}$ of the length $12 s$ (see Lemma 3). Pirk a vertex $v$ of $C$. Denote $\beta q^{2} v \quad d, \beta \varphi^{6} v-e$. Let $Z$ be the set of all vertices of $C$ of the form $\varphi^{12 n} v$, where $n=0,1,2, \ldots, \mathrm{~s} \quad$. Let $z \in Z$. From (c) of Lemma 2 it easily follows that $\beta \varphi^{2} z-d, \alpha \beta z=e$.

Define the functions $\delta_{1}, \delta_{2}, \delta_{3}, \delta_{4}$ thus ( $x$ runs through the set of all $w$-vertices):

$$
\begin{aligned}
& \delta_{1}(x)=\varphi^{5} \alpha x, \\
& \delta_{2}(x)=\varphi \alpha \varphi \alpha x, \\
& \delta_{3}(x)=\alpha \varphi^{2} x, \\
& \delta_{4}(x)=\varphi^{10} \alpha \varphi^{5} \alpha \varphi^{2} x .
\end{aligned}
$$

Let us prove that $\beta \delta_{i}(z)=d, \beta \varphi \delta_{i}(z)=e$ for $i=1,2,3$ and 4 . By systematic using of (c) of Lemma 2 and of Lemmas 7 and 8 we obtain:

```
    \(\beta \delta_{1}(z) \quad \beta \varphi^{5} \alpha z \quad \beta \varphi^{2} z-d\),
\(\beta \delta_{2}(z) \quad \beta(\varphi \alpha \varphi \alpha z) \quad \beta \varphi \alpha(\varphi \alpha z)=\alpha \beta \varphi^{2}(p \alpha z)=\alpha\left(\beta \varphi^{-1} \alpha z\right)=\alpha\left(\alpha \beta \varphi^{2} z\right)=\)
    \(\beta q^{2} z=d\),
\(\beta \delta_{3}(z) \quad \beta \alpha\left(\varphi^{2} z\right)=\beta \varphi^{2} z=d\),
\(\beta \delta_{4}(z) \quad \beta \varphi^{10} \alpha\left(\varphi^{4} \alpha \varphi^{2} z\right)=\alpha \beta \varphi\left(\varphi^{5} \alpha \varphi^{2} z\right)-\alpha\left(\beta \varphi^{6} \alpha\left(\varphi^{2} z\right)\right)=\alpha\left(\alpha \beta\left(\varphi^{2} z\right)\right)=\)
    \(\beta \varphi^{2} z-d\),
\(\beta \varphi \delta_{1}(z)=\beta \varphi^{6} \alpha z-\alpha \beta z=e\),
\(\beta \varphi \delta_{2}(z) \quad \beta q^{2} \alpha(p \alpha z)-\alpha \beta p^{-1}(\varphi \alpha z)=\alpha \beta(\alpha z)=\alpha \beta z=e\),
\(\beta \varphi \delta_{3}(z)-\beta \varphi \alpha\left(\varphi^{2} z\right) \quad \alpha \beta \varphi^{-2}\left(\varphi^{2} z\right)=\alpha, \beta z-e\),
\(\beta \varphi \delta_{4}(z) \quad \beta \varphi^{11} \alpha\left(\varphi^{5} \alpha \varphi^{2} z\right)=\alpha \beta \varphi^{2}\left(\varphi^{5} \alpha \varphi^{2} z\right) \quad \alpha \beta \varphi^{7} \alpha\left(\varphi^{2} z\right)=\alpha \beta \varphi^{-2}\left(\varphi^{2} z\right)-\)
    \(\alpha \beta z-e\).
```

Evidently, for every $z \in Z$ and $i \in\{1,2,3,4\}$ the edge $\left[\delta_{i}(z), \varphi \delta_{i}(z)\right]$ is a $w$-edge of $G_{k}$. We shall prove that all such edges are mutually different. Suppose that $\left[\delta_{i_{1}}\left(z_{1}\right), \varphi \delta_{i_{1}}\left(z_{1}\right)\right] \quad\left[\delta_{i_{2}}(2), \varphi \delta_{i_{2}}\left(z_{2}\right)\right]$, where $i_{1}, i_{2} \in\{1,2,3,4\}: z_{1}, z_{2} \in Z$. 'There are two possibilities:
I. $\delta_{i_{1}}\left(z_{1}\right) \quad \varphi \delta_{i_{2}}\left(z_{2}\right)$. But then we have $\beta \varphi^{2} v=d=\beta \delta_{i_{1}}\left(z_{1}\right) \quad \beta \varphi \delta_{i_{2}}\left(z_{2}\right)$ $e \quad \beta \varphi^{6} v$, which contradicts (d) of Lemma 2.
II. $\delta_{i_{1}}\left(z_{1}\right) \quad \delta_{i_{2}}\left(z_{2}\right)$. We first prove that $i_{1}=i_{2}$. By using (c) of Lemma 2, Lemma 7 and Lemma 8 we obtain for any $w$-vertex $x$
$\beta \varphi^{1} \delta_{1}(x) \quad \beta \varphi^{4} \alpha x \quad \beta \varphi x$,

$$
\begin{array}{rllll}
\beta \varphi^{1} \delta_{2}(x) & \beta \alpha q \alpha x=\beta q \alpha x-\alpha \beta\left(\varphi^{2} x\right) & \beta \varphi^{6}\left(\varphi^{2} x\right) & \beta q^{4} x, \\
\beta \varphi^{1} \delta_{3}(x) & \beta \varphi^{1} \alpha \varphi^{2} x \quad \alpha \beta q^{4} x \quad \beta q^{10} x, & & \\
\beta q^{2} \delta_{1}(x) & \beta q^{7} \alpha x \quad \beta \varphi{ }^{2} x \quad \beta q^{10} x, & & \\
\beta \varphi^{2} \delta_{3}(x) & \beta \varphi^{2} \alpha \varphi^{2} x-\alpha \beta q x \quad \beta q^{7} x, & & \\
\beta \varphi^{2} \delta_{4}(x) & \beta \varphi^{12} \alpha\left(\varphi^{5} \alpha \varphi^{2} x\right) & \beta \varphi^{5} \alpha \cdot\left(\varphi^{2} x\right) & \beta \varphi^{2}\left(\varphi^{2} x\right) & \beta q^{4} x .
\end{array}
$$

According to (d) of Lemma 2 the elements $\beta \varphi x, \beta \varphi^{4} x, \beta \varphi^{7} x, \beta p^{10} x$ are mutuall! different. From the equality $\delta_{i_{1}}\left(z_{1}\right) \quad \delta_{i_{2}}\left(z_{2}\right)$ it follows that $\beta \varphi{ }^{1} \delta_{r_{1}}\left(\tilde{\sim}_{1}\right)$
$\beta \varphi^{1} \delta_{i_{2}}\left(z_{2}\right)$ and $\beta q^{2} \delta_{i_{1}}\left(z_{1}\right) \quad \beta \varphi^{2} \delta_{i_{2}}\left(z_{2}\right)$. Bdt this is possible only if $i_{1} \quad i_{2}$ o1 if $\left\{i_{1}, i_{2}\right\} \quad\{2,4\}$. First analyse the second possibility. Let, e. g., $i_{1}-2, i_{2} 4$ i. e., $\delta_{2}\left(z_{1}\right)-\delta_{4}\left(z_{2}\right)$. Puty $\quad \alpha q \alpha z_{1}$. Wehave: $\beta y \quad \beta \alpha \varphi \alpha z_{1} \quad \beta q \alpha z_{1} \quad \alpha \beta \varphi{ }^{2} \tilde{z}_{1}$
$\beta \varphi^{4} z_{1}-\beta \varphi^{4} v, \beta \varphi^{3} y \quad \beta \varphi^{2}\left(\varphi \alpha \varphi \alpha z_{1}\right) \quad \beta \varphi^{2} \delta_{2}\left(z_{1}\right)-\beta \varphi^{2} \delta_{4}\left(z_{2}\right) \quad \beta \varphi^{4} z_{2}-\beta \varphi^{4} v$ Thus we obtain that $\beta y \quad \beta \varphi^{3} y$, which contradicts (d) of Lemma 2. Thereforc only the possibility $i_{1} i_{2}$ remains. Put $i \quad i_{1}-i_{2}$ so that $\delta_{i}\left(\tilde{z}_{1}\right) \quad \delta\left(\tau_{2}\right.$ $\alpha$ and $\varphi$ are one-to-one functions. Consequently also every $\delta_{i}$ is a one to ond function and from the equality $\delta_{i}\left(z_{1}\right) \quad \delta_{i}\left(z_{2}\right)$ it follows that $z_{1} \quad z_{2}$.

Thus we proved that all edges of a form $\left[\delta_{i}(z), \varphi \delta_{i}(z)\right]$, where $i \in\{1,2,3,4\}$ $\approx \in\left\{v, \varphi^{12} v, \varphi^{24} v, \ldots, \varphi^{12(s)^{1)}} v\right\}$ are mutually different. Hence we have $4 s \hookrightarrow 1 c h$ edges, and always $\beta \delta_{i}(z) \quad d, \beta \varphi \delta_{i}(z) \quad e$. According to Lemma 6 any two of the vertices $\delta_{i}(z)$ have their distance $r_{w}$ at least 6 . But from Lemma 4 it follows that we can have at most $2^{k}{ }^{5}$ such vertices. Therefore $4 s \leqslant 2^{k} 5$ i. e. the length of C is $12 s \leqslant 3.2^{k 5}$.

Theorem. There is no $M$ sore graph of type $(3, k)$, where $3 \leqslant k \leqslant 9$.
Proof. Let $G_{k}$ be a Moore graph of type ( $3, k$ ), $3 \leqslant k \leqslant 8$. Lemmas 9 and 10 imply that $k \geqslant 5$. From Lemma 3 we know that the length of any w-circuit in $G_{k}$ is a multiple of 12 . According to Lemma 11 this is possible only if $k \quad 7$ But $G_{k}$ contains no circuits of length $\leqslant 14$, especially no 12 -gons. From Lemma 11 it follows that $k-8$ and all $w$-circuits in $G_{k}$ are 24 -gons. Choone a w-circuit $C$, a vertex $v$ of $C$ and construct by the method from the proof of Lemma 11 (for $s \quad$ 2) $8 w$-edges of a form $\left(\delta_{i}(z), \varphi \delta_{i}(z)\right.$ ), where $\beta \delta_{i}(z) \quad d$ $\beta \varphi \delta_{i}(z) \quad e$. Consider the $9^{\text {th }}$ edge ( $\varphi^{1} \alpha \varphi^{6} v, \alpha \varphi^{6} v$ ). By Lemma ${ }^{2}$, (c), Lemma 7 and Lemma 8 it is easy to prove that $\beta \varphi^{1} \alpha q^{6} v-d, \beta \alpha \varphi^{6} v \quad e, \beta^{3}{ }^{2} \alpha q^{4} v$
$\beta \varphi v, \beta \varphi \alpha \varphi^{6} v \quad \beta \varphi^{10} v$. From the proof of Lemma 11 it follows that if this edge equals one of the former 8 edges, we necessarily have $i \quad$, i. e. $\varphi{ }^{1}{ }_{\alpha q}{ }^{6}{ }_{c}$, $\delta_{1}(z)$. As $C$ is a 24 -gon, either $z-v$ or $z \quad q^{12} v$. In the first case in $G_{k}$ there exists a path $\left[v, \varphi v, \varphi^{2} v, \varphi^{3} v, \varphi^{4} v, \varphi^{5} v, \varphi^{6} v,:, \alpha \varphi^{6} v \quad \varphi^{6} \alpha v, \varphi^{5} \alpha v, \varphi^{4} \alpha v, \varphi^{3} \alpha v\right.$, $\left.\varphi^{2} \alpha v, \varphi \alpha v, \alpha v,:, v\right]$; in the second case there is in $G_{k}$ a path $\left[\psi^{6} v, \varphi^{7} v, \varphi^{8} v, \varphi^{9} r^{r}\right.$, $q^{10} v, \varphi^{11} v, \varphi^{12} v,:, \alpha q^{12} v, \varphi \alpha \cdot \varphi^{12} v, \varphi^{2} \alpha \varphi^{12} v, \varphi^{3} \alpha \varphi^{12} v, \varphi^{4} \alpha \varphi^{12} v, \varphi^{5} \alpha q^{12} v, q^{6} \alpha \varphi^{12} v$ $\left.\alpha \varphi^{6} v,:, \varphi^{6} v\right]$. Both these paths contain a circuit of length $\leqslant 16$, which is in $G^{8}$ impossible. Therefore in $G^{8}$ there exist 9 edges of type $(\delta, \varepsilon)$, where $\beta \delta \quad d$
$\beta \varepsilon \quad \epsilon, \varepsilon \quad \gamma \delta$. According to Lemma 4 at least two of the vertices of type $\delta$ say $\delta^{\prime}$ and $\delta^{\prime \prime}$ have the distance $r_{w}\left(\delta^{\prime}, \delta^{\prime \prime}\right) \leqslant 4$. But this contradicts Lemma 6 . The theorem follows.

## § 4. A SURVEY OF TIED GRAP'Hs

Results of [1], [2] and our Theorem make it possible to summarize the known results on the existence and uniqueness of tied graphs of type ( $d, l$ ) moto Table 1.

Table 1


Here the symbol ? means that neither the existence nor the uniqueness of a tied graph of type $(d, k)$ has been proved. The symbol / means that there is no tied graph of the corresponding type, the symbol $E$ denotes that so far only the existence (but not the uniqueness) for a given type has been proved. In the remaining cases there exists (up to isomorphism) exactly one tied graph as indicated in the table, where $K_{n}$ is the complete graph with $n$ veitices, $C_{n}$ is the circuit with $n$ vertices, $R_{n}$ is the graph consisting of one vortex - nd $n$ loops, $P$ is the Petersen graph and $H S$ denotes the Moore gr. ty pe $(7,2)$ with 50 vertices constructed by Hoffman and Singleton in $|2|$ The ,.non trivial" part of the table is strongly framed

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