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PRODUCTIVE AND INDUCTIVE CONSTRUCTIONS OF GRAPHS*

Jiří Vinárek

0. Introduction

In [5], there is given a characterization of systems of anti-reflexive graphs in which any induced subgraph of a subdirectly irreducible (SI) graph is again SI. In the present paper we give a full characterization of hereditary subdirect irreducibility for graphs.

Importance of investigation this topic is following : having a list of SI \underline{C} -graphs one can construct any \underline{C} -graph using only products and restrictions to induced subgraphs. If a class \underline{C} of graphs is hereditary subdirectly irreducible (HSI) then the subdirect dimension coincides with the product dimension (for the definitions see [4]).

1. Notations and known facts

1.1. For the description of HSI graphs we shall use some symbols from [5] and also introduce some new ones.

If A is an induced subgraph of B we shall write $A \longleftrightarrow B$. For an arbitrary graph G denote $V(G)$ its set of vertices and $E(G)$ its set of edges. $L(G) \longleftrightarrow G$ is a graph such that $V(L(G)) = \{v \in V(G) ; (v,v) \in E(G)\}$. $N(G) \longleftrightarrow G$ is a graph with $V(N(G)) = V(G) - V(L(G))$. (Edges of $L(G)$ are denoted as LL -edges, edges of $N(G)$ as NN -edges, edges from $L(G)$ to $N(G)$ as LN -edges and edges from $N(G)$ to $L(G)$ as NL -edges.)

For any ordinal n denote $K_n = (n, \{(i,j) ; i, j \in n, i \neq j\})$ (i.e. the complete antireflexive graph with n vertices),

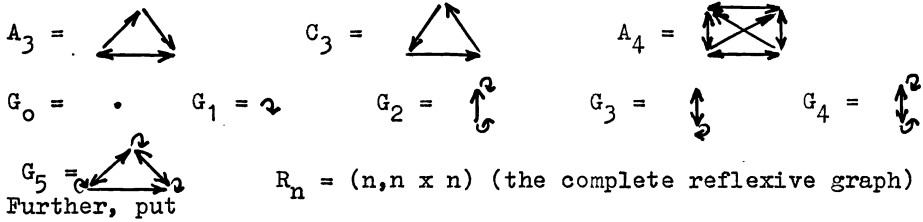
$$K_n^+ = (n, \{(i,j) ; i, j \in n, i \neq j, (i,j) \neq (0,1)\}),$$

$$L_n^n = (n, \{(i,j) ; i, j \in n, i < j\}),$$

$$L_n^+ = (n, \{(i,j) ; i, j \in n, i < j\} \cup \{(1,0)\}),$$

$$L_n^- = (n, \{(i,j) ; i, j \in n, i > j\} \cup \{(0,1)\}),$$

*) This paper is in final form and no version of it will be submitted for publication elsewhere,



- $\underline{K} = \{K_n ; n \in \text{Ord}\},$
- $\underline{K}' = \{K'_n ; n \in \text{Ord}\},$
- $\underline{L}^+ = \{L^+_n ; n \in \text{Ord}\},$
- $\underline{L}^- = \{L^-_n ; n \in \text{Ord}\},$

Set = $\{(X, \emptyset) ; X \text{ is a set}\}$ (the class of sets = discrete graphs),

$\underline{T} = \{G ; |\{(x,y), (y,x)\} \cap E(G)| = 1 \text{ for any } x \neq y \in V(G)\}$ (the class of all antireflexive tournaments)

$\underline{U} = \{(n,R) ; n \leq 6, |R| = \binom{n}{2} + \lfloor \frac{n}{2} \rfloor, x \neq y \Rightarrow |\{(x,y), (y,x)\} \cap R| \geq 1 \text{ and } (n,R) \text{ contains neither } K_3 \text{ nor } A_3 \text{ as an induced subgraph}\},$

$\underline{V} = \{(n,R) ; n \leq 4, x \neq y \Rightarrow |R \cap \{(x,y)\}| \geq 1, R \supseteq \{(0,1), (1,0), (2,3), (3,2)\} \cap n \times n \text{ and } (n,R) \text{ does not contain } K_3 \text{ as an induced subgraph}\},$

$\underline{W} = \{A ; \text{any induced subgraph of } G \text{ with 3 vertices is either isomorphic to } A_3, \text{ or to } L_3\},$

$\underline{X}_0 = \{X ; V(X) = V \cup \{v\}, E(X) = E \cup \{(v,v)\} \text{ where } (V,E) \in \underline{X}\}$
for $\underline{X} \in \{\underline{K}, \underline{K}', \underline{L}^+, \underline{L}^-, \underline{T}, \underline{U}, \underline{V}, \underline{W}\},$

$\underline{\text{Sym}}_5 = \{A ; A \text{ is reflexive symmetric, } |V(A)| \leq 5\}.$

1.2. By a product of graphs we mean the categorical product (i.e.

$$\prod_{i \in I} (V_i, E_i) = (\prod_{i \in I} V_i, E) \text{ where } ((x_i)_I, (y_i)_I) \in E \text{ iff } (x_i, y_i) \in E_i \text{ for any } i \in I.$$

1.3. Let \underline{C} be a class of graphs. Then $A \in \underline{C}$ (i.e. a \underline{C} -graph A) is said to be subdirectly irreducible if, whenever an isomorphic copy A' of A is contained as an induced subgraph in a product $\prod_{i \in I} B_i$ with $B_i \in \underline{C}$ and $p_j(A') = B_j$ for all the projections, there is a j such that the restriction of p_j to A' is an isomorphism onto B_j . (This formulation is due to A.Pultr - see [2].)

1.4. A class \underline{C} of graphs is said to be hereditary with respect to subdirect irreducibility (HSI) if any induced subgraph of a SI \underline{C} -graph is again SI (see [5]).

1.5. If $V(A) = V(B)$ then the meet of graphs $A \wedge B$ denotes the graph $(V(A), E(A) \cap E(B))$. If $C = A \wedge B, C \neq A, B$ then C is subdirectly reducible in \underline{C} (see [3]).

1.6. Let \underline{D} be a family of graphs. Then $\text{SP}(\underline{D})$ denotes (similarly as

in [1]) the class of all the graphs which can be embedded as induced subgraphs into products of graphs from \underline{D} .

1.7. Theorem (see[5]). Let \underline{C} be a productive hereditary class of antireflexive graphs (i.e. a class closed to categorical products and to induced subgraphs). Then \underline{C} is HSI iff either $\underline{C} = \underline{Set}$ or $\underline{C} = SP(\underline{D})$ where \underline{D} satisfies one of the following conditions :

- (i) $\underline{D} \subseteq \underline{K} \cup \underline{K}'$
- (ii) $\underline{D} \subseteq \underline{K} \cup \{K_3', A_4\}$
- (iii) $\underline{D} \subseteq \underline{K} \cup \underline{I}_7^+ \cup \underline{T}$
- (iv) $\underline{D} \subseteq \underline{K} \cup \underline{I}_7^- \cup \underline{T}$
- (v) $\underline{D} \subseteq \underline{K} \cup \underline{U}$
- (vi) $\underline{D} \subseteq \underline{K} \cup \underline{V}$
- (vii) $\underline{D} \subseteq \underline{K} \cup \underline{W}$

2. Hereditary subdirect irreducibility

Before giving the general characterization theorem for HSI in graphs we shall consider partial cases discussing possibilities for reflexive and antireflexive parts of graphs and for LL-,LN- and NL-edges.

Throughout this chapter, \underline{C} denotes a productive hereditary class of graphs which is HSI.

2.1.Lemma. If any reflexive subgraph of a SI \underline{C} -graph A is complete then $|L(A)| \leq 2$.

Proof. Since \underline{C} is HSI, any L(A) is SI whenever A is SI. Any reflexive complete graph is an induced subgraph of a power R_2^k for some k. Hence, $|L(A)| \leq 2$.

2.2.Lemma. If some LN-edge of a SI \underline{C} -graph A_0 is $\sigma \leftrightarrow$ then no edge of a SI \underline{C} -graph A is \leftrightarrow . Moreover, if some LL-edge of A_0 is $\sigma \rightarrow \tau$ then no edge of A is \rightarrow .

Proof. $\leftrightarrow = \sigma \leftrightarrow \wedge \leftrightarrow \tau$, hence it is reducible and it cannot be an induced subgraph of a SI graph A. If $\sigma \rightarrow \tau \in \underline{C}$ then $\rightarrow = \sigma \rightarrow \wedge \sigma \rightarrow \tau$ is subdirectly reducible and hence it cannot be an induced subgraph of A.

2.3.Lemma. Let A_0, A be SI \underline{C} -graphs. If some LN-edge of A_0 is $\sigma \rightarrow$ ($\sigma \leftarrow$, resp.) then either any NN-edge of A is \leftrightarrow or any NN-edge of A is \rightarrow .

Proof. $\rightarrow = \sigma \rightarrow \wedge \leftarrow = \sigma \leftarrow \wedge \leftarrow$. Hence, graphs \leftrightarrow and \rightarrow cannot be both SI in \underline{C} .

2.4.Lemma. If $\sigma \leftrightarrow$ and $\sigma \rightarrow$ ($\rightarrow \tau$, resp.) are both \underline{C} -graphs then $|N(A)| \leq 1$ for any SI \underline{C} -graph A.

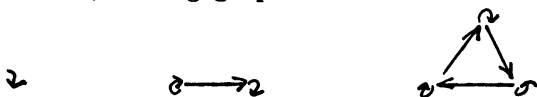
Proof. There is $\leftrightarrow = \sigma \leftrightarrow \wedge \leftrightarrow \tau$, $\leftarrow = \sigma \leftarrow \wedge \leftarrow \tau =$

$= \leftarrow 2 \wedge \leftarrow$, $\cdot \cdot = \leftarrow \wedge \leftarrow = \leftarrow 2 \wedge \leftarrow$, hence any antireflexive graph with at least 2 vertices is subdirectly reducible in \underline{C} and $|N(A)| \leq 1$ for any SI \underline{C} -graph A.

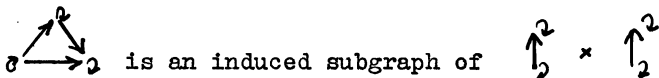
2.5.Lemma. If $\leftarrow 2$ and \leftarrow are both \underline{C} -graphs then any NN-edge of a SI \underline{C} -graph A is \leftarrow and no LN-edge of a SI \underline{C} -graph is \downarrow .

Proof. There is $\leftarrow = \leftarrow 2 \wedge \leftarrow$, $\cdot \cdot = \leftarrow \wedge \leftarrow$. Hence, any NN-edge of a SI \underline{C} -graph is \leftarrow . Moreover, $\downarrow \cdot = \leftarrow \wedge \leftarrow$, hence no LN-edge of a SI \underline{C} -graph is $\downarrow \cdot$.

2.6.Lemma. If a reflexive tournament A is SI in \underline{C} then A is one of the following graphs :



Proof.



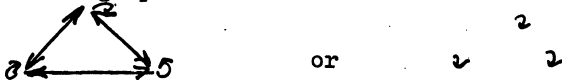
hence it is not SI. Any reflexive tournament with at least 4 vertices contains



as an induced subgraph, hence it is not SI.

2.7.Lemma. If a symmetric reflexive graph A is SI in \underline{C} then $|V(A)| \leq 5$.

Proof. Using Dirichlet principle, one can check that any symmetric reflexive graph G with at least 6 vertices contains either



as an induced subgraph. Since



G is not SI in \underline{C} .

2.8.Proposition. If a reflexive graph A is SI in \underline{C} then $|V(A)| \leq 9$.

Proof. If A is symmetric then $|V(A)| \leq 5$ due to Lemma 2.7. If A is not symmetric then it contains $\downarrow \rightarrow$ as an induced subgraph. Hence,

$\downarrow \downarrow = \downarrow \rightarrow \wedge \leftarrow$ is not SI in \underline{C} . Therefore, $\downarrow \downarrow$ is not

an induced subgraph of A. Using Ramsey theorem one can prove that any reflexive graph with at least 9 vertices which does not contain $\downarrow \downarrow$ as an induced subgraph contains either



or a tournament with 4 vertices as an induced subgraph. Using Lemmas 2.1 and 2.6 one proves that G is not SI.

2.9. Lemma. If a reflexive graph A is SI in \underline{C} and contains G_5 as an induced subgraph, then $A = G_5$.

Proof. Using Proposition 4.6 from [3] one can prove that any reflexive graph is a \underline{C} -graph. Since any reflexive graph with at least 4 vertices is - due to [3] - subdirectly reducible and A is SI in \underline{C} , there is $A = G_5$.

2.10. Lemma. If G_5 is not a \underline{C} -graph and G_2 is a \underline{C} -graph then any reflexive SI \underline{C} -graph has the following property : For an $(x,y) \notin E(A)$ define $U(x,y)$ as the smallest subset of $V(A) \times V(A) - E(A)$ containing (x,y) and such that $(V(A), E(A) \cup U(x,y)) \in \underline{C}$. Then for any two $(x_0, y_0), (x_1, y_1) \in (V(A) \times V(A)) - E(A)$, $U(x_0, y_0) \cap U(x_1, y_1) \neq \emptyset$, and for every morphism $\varphi: A \rightarrow B$ with $|V(B)| < |V(A)|$ there is an $(x,y) \in (V(A) \times V(A)) - E(A)$ with $(\varphi(x), \varphi(y)) \in E(B)$.

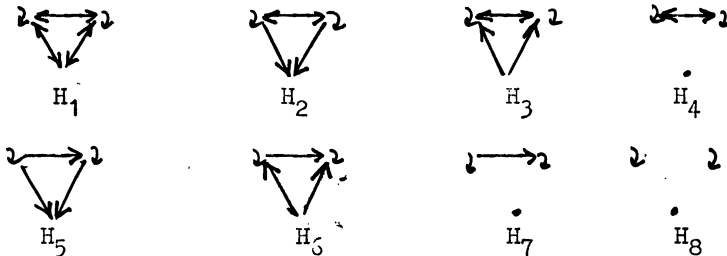
Proof follows directly from [3], Lemma 6.8.

2.11. Proposition. If a reflexive graph A is SI in \underline{C} then it satisfies one of the following conditions :

- (i) $A \in \text{Sym}_5$
- (ii) $A \leftrightarrow G_5$
- (iii) $|V(A)| \leq 9$ and A satisfies the conditions from Lemma 2.10. (Denote this class of graphs by Ref_9 .)

Proof follows from 2.7-2.10.

2.12. Lemma. The following graphs are subdirectly reducible in \underline{C} whenever they are \underline{C} -graphs:



Proof. $H_1 \leftrightarrow G_3 \times R_2^2$, $H_2 \leftrightarrow \downarrow \times R_2^2$, $H_3 \leftrightarrow \uparrow \times R_2^2$, $H_4 \leftrightarrow \cdot \times R_2^2$, $H_5 \leftrightarrow \downarrow \times G_2^2$, $H_6 \leftrightarrow \uparrow \times G_2^2$, $H_7 \leftrightarrow \cdot \times G_2^2$, $H_8 \leftrightarrow \cdot \times \cdot$.

2.13. Proposition. If A is SI in \underline{C} , $L(A) \neq \emptyset$ is a complete graph and $(l,n), (n,l) \in E(A)$ whenever $l \in L(A)$, $n \in N(A)$ then either A is one of the following graphs : G_1, G_3, G_4 , or $|L(A)| = 1$ and $N(A)$ is an antireflexive tournament. (Denote this class of graphs by \mathbb{T}_1 .)

Proof. According to 2.1, $|L(A)| \leq 2$. Hereditary subdirect irreducibility implies that $N(A)$ satisfies conditions of Characterization Theorem 1.7. Consider two cases :

- 1. $|L(A)| = 1$. Consider possibilities for $N(A)$ using 1.7.

(i) If $N(A) \in \underline{K} \cup \underline{K}'$ then by 2.2 there is $|N(A)| \leq 2$. For the case $|N(A)| = 0$ one obtains $A = G_1$, for the case $|N(A)| = 1$ there is $A = G_3$ and for the case $|N(A)| = 2$ there is $N(A) = K_2'$ which is an antireflexive tournament on 2 points.

(ii) If $N(A) = A_4$ then one obtains a contradiction using 2.2.

(iii) If $N(A) \in \underline{L}^+ \cup \underline{T}$ then $N(A)$ is an antireflexive tournament.

(iv) If $N(A) \in \underline{L}^- \cup \underline{T}$ then $N(A)$ is an antireflexive tournament as well.

(v) If $N(A) \in \underline{U}$ then one obtains \longleftrightarrow as an induced subgraph which contradicts 2.2.

(vi) If $N(A) \in \underline{V}$ then one obtains a contradiction with 2.2, too.

(vii) If $N(A) \in \underline{W}$ then using 2.2 one obtains that $N(A)$ is an antireflexive tournament.

2. $|L(A)| = 2$.

If $N(A) = \emptyset$ then $A = G_4$.

If $N(A) \neq \emptyset$ then A contains H_1 as an induced subgraph which contradicts 2.12.

2.14. Proposition. If A is SI in \underline{C} , $L(A) \neq \emptyset$ is a complete graph, $N(A) = \emptyset$ and $(l,n) \in E(A)$, $(n,l) \notin E(A)$ whenever $l \in L(A)$, $n \in N(A)$, then $|L(A)| = 1$ and either

(i) $N(A) = K_n$ for some n ,

or

(ii) $N(A)$ is an antireflexive tournament.

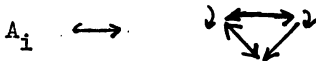
(Denote the class of graphs satisfying (i) ((ii), resp.) by $\underline{K}_1^{\downarrow}$ ($\underline{T}_1^{\downarrow}$, resp.).

Proof. Lemmas 2.1 and 2.12 imply that $|L(A)| = 1$. Using 1.7 and 2.3 one obtains that $N(A)$ is either antireflexive complete or an antireflexive tournament.

2.15. Proposition. If A is SI in \underline{C} , $L(A) \neq \emptyset$ is complete, $N(A) \neq \emptyset$ and $(l,n) \notin E(A)$, $(n,l) \in E(A)$ whenever $l \in L(A)$, $n \in N(A)$, then $|L(A)| = 1$ and either $N(A) = K_n$ for some n or $N(A)$ is an antireflexive tournament. (Denote the corresponding class of graphs by $\underline{K}_1^{\uparrow} \cup \underline{T}_1^{\uparrow}$.)

Proof is similar to the proof of Proposition 2.14.

2.16. Proposition. If $A_i (i=0,1)$ are SI in \underline{C} , $L(A_i) \neq \emptyset$ are complete, $N(A_0) \neq \emptyset$, $(l,n) \in E(A_i)$ whenever $l \in L(A_i)$, $n \in N(A_i)$ and if there exist $n_1, n_2 \in N(A_0)$, $l_1, l_2 \in L(A_0)$ such that $(n_1, l_1) \in E(A_0)$, $(n_2, l_2) \notin E(A_0)$ then



Proof. By 2.2 and 2.3, \leftrightarrow and \rightarrow (and also $\circ \circ = \rightarrow \wedge \leftarrow$) are subdirectly reducible. Hence, $|N(A_1)| \leq 1$. By 2.1 and the assumptions of Proposition, $|L(A_1)| \leq 2$. Therefore,

$$A_0 = \begin{array}{c} \swarrow \quad \searrow \\ \circ \quad \circ \\ \swarrow \quad \searrow \end{array}, A_1 \leftrightarrow A_0.$$

2.17. Proposition. If $A_i (i = 0, 1)$ are SI in \underline{C} , $L(A_1) \neq \emptyset$ are complete, $N(A_0) \neq \emptyset$, $(n, l) \in E(A_1)$ whenever $l \in L(A_1)$, $n \in N(A_1)$ and if there exist $n_1, n_2 \in N(A_0)$, $l_1, l_2 \in L(A_0)$ such that $(l_1, n_1) \in E(A_0)$, $(l_2, n_2) \in E(A_0)$ then

$$A_1 \leftrightarrow \begin{array}{c} \swarrow \quad \searrow \\ \circ \quad \circ \\ \swarrow \quad \searrow \end{array}$$

Proof is similar to the proof of 2.16.

2.18. Proposition. If $A_i (i = 0, 1)$ are SI in \underline{C} , $L(A_0) \neq \emptyset$ is either complete or a tournament, $N(A_0) \neq \emptyset$, $|\{(n, l), (l, n)\} \cap E(A_1)| = 1$ for any $(n, l) \in N(A_1) \times L(A_1)$ and if there exist $n_1, n_2 \in N(A_0)$, $l_1, l_2 \in L(A_0)$ such that $(l_1, n_1) \in E(A_0)$, $(n_2, l_2) \in E(A_0)$ then $N(A_1) = K_n$ for some n , $(l', n) \in E(A_1) \Leftrightarrow (n, l'') \in E(A_1)$ whenever $l' \neq l'' \in L(A_1)$. (Denote the class of graphs A_1 satisfying these conditions by \underline{L} .)

Proof. By 2.5, any NN-edge of A_1 is \leftrightarrow . Hence, $N(A_1) = K_n$ for some n . By 2.12, A_1 does not contain neither H_2 nor H_3 as an induced subgraph.

2.19. Proposition. If A is SI in \underline{C} , $L(A) \neq \emptyset$, $N(A) \neq \emptyset$ and if there are no NL- and LN-edges in A then $|L(A)| = 1$.

Proof follows from 2.12 because A cannot contain H_4 , H_7 and H_8 as induced subgraphs.

2.20. Proposition. If $A_i (i = 0, 1)$ are SI in \underline{C} , $L(A_1) \neq \emptyset$ is complete, $N(A_1) \neq \emptyset$, $\emptyset \neq L(A_0) \times N(A_0) \cap E(A_0) \neq L(A_0) \times N(A_0)$, and for any $(l, n) \in L(A_1) \times N(A_1)$ there is $(l, n) \in E(A_1) \Leftrightarrow (n, l) \in E(A_1)$, then $|L(A_1)| \leq 2$ and $N(A_1)$ is an antireflexive tournament. Moreover, if $|L(A_1)| = 2$ then $L(A_1) = \{l', l''\}$ such that $(l', n) \in E(A_1) \Leftrightarrow (l'', n) \in E(A_1)$ for any $n \in N(A_1)$.

(Denote the class of graphs satisfying these conditions by \underline{M} .)

Proof follows from 1.7, 2.1, 2.2 and 2.12.

2.21. Proposition. If $A_i (i = 0, 1)$ are SI in \underline{C} , $L(A_0) \neq \emptyset$, $N(A_0) \neq \emptyset$, $\emptyset \neq L(A_0) \times N(A_0) \cap E(A_0) \neq L(A_0) \times N(A_0)$ and $N(A_1) \times L(A_1) \cap E(A_1) = \emptyset$, then $N(A_1)$ is either an antireflexive tournament or an antireflexive complete graph. Moreover, if $l' \neq l'' \in L(A_1)$ then $(l', n) \in E(A_1) \Leftrightarrow (l'', n) \notin E(A_1)$ for any

$n \in N(A_i)$. (Denote the class of graphs satisfying these conditions by \underline{N} .)

Proof follows from 1.7, 2.3 and 2.12.

2.22. Proposition. If $A_i (i = 0, 1)$ are SI in \underline{C} , $L(A_0) \neq \emptyset$, $N(A_0) \neq \emptyset \neq N(A_0) \times L(A_0) \cap E(A_0) \neq N(A_0) \times L(A_0)$ and $L(A_1) \times N(A_1) \cap \cap E(A_1) = \emptyset$ then $N(A_1)$ is either an antireflexive tournament, or an antireflexive complete graph. Moreover, if $l' \neq l'' \in L(A)$ then $(n, l') \in E(A_1) \Leftrightarrow (n, l'') \notin E(A_1)$ for any $n \in N(A_1)$. (Denote the class of graphs satisfying these conditions by \underline{P} .)

Proof follows from 1.7, 2.3 and 2.12.

2.23. Lemma. If there are $(l_1, n_1), (l_2, n_2), (l_3, n_3) \in L(A) \times N(A)$ such that $(l_1, n_1), (n_1, l_1), (n_2, l_2), (l_3, n_3) \notin E(A)$, $(l_2, n_2), (n_3, l_3) \in E(A)$ then A is subdirectly reducible in \underline{C} .

Proof. Since $\cdot \cdot, \cdot \rightarrow$ and $\cdot \leftarrow$ are induced subgraphs of A , $\cdot \cdot = \cdot \rightarrow \wedge \cdot \leftarrow$ and \underline{C} is HSI, A cannot be SI.

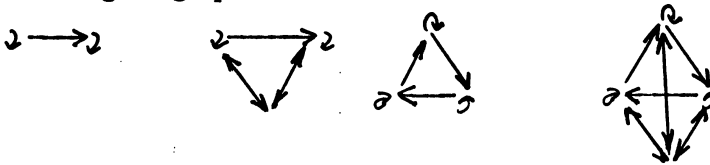
2.24. Lemma. If $L(A)$ is complete and if there are $(l_1, n_1), (l_2, n_2), (l_3, n_3) \in L(A) \times N(A)$ such that $(l_1, n_1), (l_3, n_3), (n_3, l_3) \in E(A)$ and $(n_2, l_2) \notin E(A), (l_2, n_2) \in E(A)$ ($(n_2, l_2) \in E(A), (l_2, n_2) \notin E(A)$ resp.) then A is subdirectly reducible in \underline{C} .

Proof. Suppose A be SI. Then 2.1 implies that $|L(A)| \leq 2$, 2.4 implies that $|N(A)| \leq 1$ and there are at most 2 LN- and 2 NL-edges which contradicts assumptions of Lemma.

2.25. Proposition. If A is a SI \underline{C} -graph which contains G_2 as an induced subgraph, $L(A) \times N(A) \subset E(A)$, $N(A) \times L(A) \subset E(A)$, then $|N(A)| \leq 1$. (Denote the class of graphs satisfying these conditions by \underline{Q} .)

Proof. Lemma 2.2 implies that no edge of A is neither \leftrightarrow nor \rightarrow . Since $\cdot \cdot = \rightarrow \wedge \leftarrow$ there is $|N(A)| \leq 1$.

2.26. Corollary. For the case that $L(A)$ is a tournament we obtain the following SI graphs :



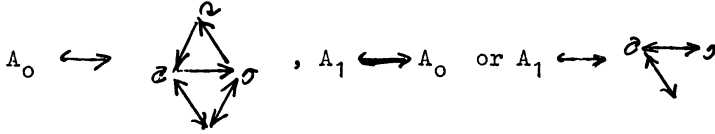
2.27. Lemma. If $G_2 \leftrightarrow A$, $N(A) \neq \emptyset$, $L(A) \times N(A) \subseteq E(A)$, $N(A) \times L(A) \cap E(A) = \emptyset$ ($L(A) \times N(A) \cap E(A) = \emptyset, N(A) \times L(A) \subseteq E(A)$ resp.), then A is subdirectly reducible in \underline{C} .

Proof follows from 2.12.

2.28. Lemma. If A contains G_2, G_3 and $\cdot \rightarrow$ (\rightarrow_2 resp.) as induced subgraphs then A is not SI in \underline{C} .

Proof. $\sigma \rightarrow = \sigma \leftarrow \wedge \sigma \rightarrow$, $\rightarrow 2 = \leftarrow 2 \wedge \rightarrow 2$, hence A is not SI in \underline{C} .

2.29.Proposition. If $A_i (i = 0, 1)$ are SI in \underline{C} , $G_2 \leftrightarrow A_0$, $\emptyset \neq L(A_0) \times N(A_0) \cap E \neq L(A_0) \times N(A_0)$ and if for any $(l, n) \in L(A_1) \times N(A_1)$ there is $(l, n) \in E(A_1) \iff (n, l) \in E(A_1)$ then



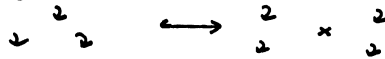
Proof follows from 2.12 and 2.2.

2.30.Lemma. If a reflexive graph A is SI in \underline{C} and $\rightarrow 2$ is also SI in \underline{C} then A is symmetric.

Proof. $\rightarrow 2 = \sigma \rightarrow \wedge \sigma \leftarrow$, hence A cannot contain G_2 as an induced subgraph.

2.31.Proposition. If $A_i (i = 0, 1)$ and $\rightarrow 2$ are SI in \underline{C} , $L(A_0) \times N(A_0) \cap E(A_0) \neq \emptyset$ and $(n, l) \in E(A_1) \iff (l, n) \in E(A_1)$ for any $(n, l) \in L(A_1) \times N(A_1)$ then $L(A_1) \times N(A_1) \subseteq E(A_1)$ and $N(A_1)$ is an antireflexive tournament. Moreover, $L(A_1) \leftrightarrow \rightarrow 2$ (Denote the class of graphs satisfying these conditions by \underline{T}_2 .)

Proof. According to 2.30, $L(A_1)$ is symmetric. By 2.2, $N(A_1)$ is an antireflexive tournament. Since $\rightarrow 2 = \rightarrow 2 \wedge \sigma \leftarrow$ there is $L(A_1) \times N(A_1) \subseteq E(A_1)$. Lemma 2.12 implies that A_1 cannot contain H_1 as an induced subgraph. Hence, any LL-edge of A_1 is $\rightarrow 2$. Since



there is $L(A_1) \leftrightarrow \rightarrow 2$.

2.32.Proposition. If A and $\rightarrow 2$ are SI in \underline{C} , $L(A) \times N(A) \cap E(A) \neq \emptyset$ and $N(A) \times L(A) \cap E(A) = \emptyset$ ($L(A) \times N(A) \cap E(A) = \emptyset$ and $N(A) \times L(A) \cap E(A) \neq \emptyset$, resp.) then $L(A) \times N(A) \subseteq E(A)$ ($N(A) \times L(A) \subseteq E(A)$ resp.) and either $N(A) = K_n$ for some n or $N(A)$ is an antireflexive tournament and $L(A) \leftrightarrow \rightarrow 2$.

(Denote the class of graphs satisfying these conditions by $K_2 \downarrow \cup \underline{T}_2 \uparrow$ ($K_2 \uparrow \cup \underline{T}_2 \uparrow$ resp.).)

Proof. By 1.7 and 2.3, $N(A)$ is either complete or a tournament. Since $\rightarrow 2 = \rightarrow 2 \wedge \sigma \rightarrow = \rightarrow 2 \wedge \sigma \leftarrow$, there is $L(A) \times N(A) \subseteq E(A)$ ($N(A) \times L(A) \subseteq E(A)$ resp.). Lemma 2.12 implies that A cannot contain neither H_5 nor H_6 as an induced subgraph. Hence, any LL-edge of A is $\rightarrow 2$ and $L(A) \leftrightarrow \rightarrow 2$.

2.33. Proposition. If $A_i (i = 0, 1)$ are SI in \underline{C} , $\rightsquigarrow \rightsquigarrow \longleftrightarrow A_0$,
 $\rightsquigarrow \longleftrightarrow \longleftrightarrow A_0, (L(A_0) \times N(A_0), \cup N(A_0) \times L(A_0)) \cap E(A_0) \neq E(A_0)$,
 $|\{(1, n), (n, 1)\} \cap E(A_i)| \geq 1$ for any $(1, n) \in L(A_i) \times N(A_i)$, then
 $|N(A_i)| \leq 1, |L(A_i)| \leq 5, L(A_i)$ is symmetric and for any $n \in N(A_i)$
and $1', 1'' \in L(A_i)$ such that $(1', 1'') \in E(A_i)$ there does not
hold both $(1', n) \in E(A_i) \Leftrightarrow (1'', n) \in E(A_i)$ and $(n, 1') \in E(A_i) \Leftrightarrow$
 $\Leftrightarrow (n, 1'') \in E(A_i)$. (Denote the corresponding class of graphs
by \underline{R} .)

Proof. By 2.30, $L(A_i)$ are symmetric. By 2.7, $|L(A_i)| \leq 5$. Lemma
2.4 implies that $|N(A_i)| \leq 1$. Lemma 2.12 - which implies that
 A_i cannot contain H_1, H_2, H_3 as induced subgraphs - finishes
the proof.

2.34. Proposition. If A and $\rightsquigarrow \rightsquigarrow$ are SI in \underline{C} and if for any
 $(1, n) \in L(A) \times N(A)$ there is $(1, n) \in E(A) \Leftrightarrow (n, 1) \notin E(A)$ then
 $L(A)$ is symmetric, $|L(A)| \leq 5, N(A) = K_n$ for some n and for any
 $1', 1'' \in L(A)$ such that $(1', 1'') \in E(A)$ there is $(1', n) \in E(A) \Leftrightarrow$
 $\Leftrightarrow (n, 1'') \in E(A)$ for any $n \in N(A)$. (Denote the corresponding
class of graphs by \underline{S} .)

Proof. By 2.30, $L(A)$ is symmetric. By 2.7, $|L(A)| \leq 5$. Lemma 2.5
implies that $N(A)$ is a tournament and $\rightsquigarrow \circ$ is not an induced
subgraph of A . By 2.12, for any $n \in N(A)$ and $1', 1'' \in L(A)$ such
that $(1', 1'') \in E(A)$ there is $(1', n) \in E(A) \Leftrightarrow (n, 1'') \in E(A)$.

3. Characterization Theorem

Now, we can prove the following :

3.1. Theorem. Let \underline{C} be a productive hereditary class of graphs.
Then \underline{C} is HSI iff either $\underline{C} = \underline{Set}$ or $\underline{C} = SP(\underline{D})$ where \underline{D} satisfies
the following conditions :

- (i) $\underline{D} \subseteq \underline{K} \cup \underline{K}' \cup \underline{K}_0 \cup \underline{K}'_0$
- (ii) $\underline{D} \subseteq \underline{K} \cup \{K_3, A_4\} \cup \underline{K}_0 \cup \{(K_3)_0, (A_4)_0\}$
- (iii) $\underline{D} \subseteq \underline{K} \cup \underline{L}^+ \cup \underline{T} \cup \underline{K}_0 \cup \underline{L}_0^+ \cup \underline{T}_0$
- (iv) $\underline{D} \subseteq \underline{K} \cup \underline{L}^- \cup \underline{T} \cup \underline{K}_0 \cup \underline{L}_0^- \cup \underline{T}_0$
- (v) $\underline{D} \subseteq \underline{K} \cup \underline{U} \cup \underline{K}_0 \cup \underline{U}_0$
- (vi) $\underline{D} \subseteq \underline{K} \cup \underline{V} \cup \underline{K}_0 \cup \underline{V}_0$
- (vii) $\underline{D} \subseteq \underline{K} \cup \underline{W} \cup \underline{K}_0 \cup \underline{W}_0$
- (viii) $\underline{D} \subseteq \underline{Sym}_5 \cup \{G_5\} \cup \underline{Ref}_9$
- (ix) $\underline{D} \subseteq \{G_1, G_3, G_4\} \cup \underline{T}_1$
- (x) $\underline{D} \subseteq \underline{T}_1 \cup \underline{K}_1$
- (xi) $\underline{D} \subseteq \underline{T}_1 \cup \underline{K}_1$
- (xii) $\underline{D} \subseteq \{ \begin{matrix} \longleftrightarrow \\ \searrow \end{matrix} \}$

- (xiii) $\underline{D} \subseteq \{ \overset{\curvearrowright}{\leftarrow} \overset{\curvearrowleft}{\rightarrow} \}$
- (xiv) $\underline{D} \subseteq \underline{L}$
- (xv) $\underline{D} \subseteq \underline{M}$
- (xvi) $\underline{D} \subseteq \underline{N}$
- (xvii) $\underline{D} \subseteq \underline{P}$
- (xviii) $\underline{D} \subseteq \underline{Q}$
- (xix) $\underline{D} \subseteq \{ \text{graph 1}, \text{graph 2} \}$
- (xx) $\underline{D} \subseteq \mathbb{T}_2^{\curvearrowright}$
- (xxi) $\underline{D} \subseteq \mathbb{K}_2^{\curvearrowright} \cup \mathbb{T}_2^{\curvearrowleft}$
- (xxii) $\underline{D} \subseteq \mathbb{K}_2^{\curvearrowleft} \cup \mathbb{T}_2^{\curvearrowright}$
- (xxiii) $\underline{D} \subseteq \underline{R}$
- (xxiv) $\underline{D} \subseteq \underline{S}$

Proof. I* If \underline{C} is HSI then we can consider the following cases :

- a) All \underline{C} -graphs are antireflexive. Then by 1.7, $\underline{C} = SP(\underline{D})$ where \underline{D} satisfies one of the conditions (i)-(vii).
- b) All \underline{C} -graphs are reflexive. Then by 2.11, $\underline{C} = SP(\underline{D})$ where \underline{D} satisfies (viii).
- c) There exist \underline{C} -graphs with loops and also \underline{C} -graphs without loops. Then we can divide the proof of Theorem discussing possibilities for LL-, LN- and NL-edges of SI \underline{C} -graphs:

LL-edges	LN- and NL-edges	see	$\underline{C}=SP(\underline{D})$ where \underline{D} satisfies :
$\overset{\curvearrowright}{\leftarrow} \overset{\curvearrowleft}{\rightarrow}$	$\overset{\curvearrowright}{\leftarrow}$	2.13	(ix)
$\overset{\curvearrowright}{\leftarrow} \overset{\curvearrowleft}{\rightarrow}$	$\overset{\curvearrowright}{\leftarrow}$	2.14	(x)
$\overset{\curvearrowright}{\leftarrow} \overset{\curvearrowleft}{\rightarrow}$	$\overset{\curvearrowleft}{\rightarrow}$	2.15	(xi)
$\overset{\curvearrowright}{\leftarrow} \overset{\curvearrowleft}{\rightarrow}$	$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow}$	2.16	(xii)
$\overset{\curvearrowright}{\leftarrow} \overset{\curvearrowleft}{\rightarrow}$	$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow}$	2.17	(xiii)
$\overset{\curvearrowright}{\leftarrow} \text{ or } \overset{\curvearrowleft}{\rightarrow}$	$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow}$	2.18	(xiv)
arbitrary	$\overset{\curvearrowright}{\leftarrow}$	2.19	(i)-(vii)
$\overset{\curvearrowright}{\leftarrow} \overset{\curvearrowleft}{\rightarrow}$	$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow}$	2.20	(xv)
arbitrary	$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow}$	2.21	(xvi)
arbitrary	$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow}$	2.22	(xvii)
arbitrary	$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow} \quad \overset{\curvearrowright}{\leftarrow}$	2.1	contradiction
arbitrary	$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow} \quad \overset{\curvearrowleft}{\rightarrow}$	2.23	contradiction
arbitrary	$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow} \quad \overset{\curvearrowright}{\leftarrow}$	2.24	contradiction
arbitrary	$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow} \quad \overset{\curvearrowleft}{\rightarrow}$	2.24	contradiction
$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow}$	$\overset{\curvearrowright}{\leftarrow}$	2.25	(xviii)
$\overset{\curvearrowright}{\leftarrow} \quad \overset{\curvearrowleft}{\rightarrow}$	$\overset{\curvearrowright}{\leftarrow}$	2.27	contradiction

