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NOTE ON A SPLITTING APPROACH TO ILL-CONDITIONED
LEAST SQUARES PROBLEMS

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1. INTRODUCTION

For an $m \times n$ real matrix A of rank n and a real m -vector b , the least squares problem for the linear system $Ax = b$ is to determine the n -vector \tilde{x} such that

$$\|b - A\tilde{x}\|_2 \leq \|b - Ax\|_2$$

for all n -vectors x . The unique solution is given by

$$\tilde{x} = (A^T A)^{-1} A^T b$$

where “ T ” denotes the transpose. The traditional approach to the problem is to solve the “normal system”

$$(1.1) \quad A^T A x = A^T b$$

by some standard direct procedure such as Gaussian or Cholesky elimination. However, using the computer this approach is often poor, since the condition number of $A^T A$ is the square of the condition number $\kappa(A)$, of A . (Here we take the condition number to be the ratio of the largest to the smallest singular value of the matrix.) In fact, using t -digit binary arithmetic one is not able to obtain even an approximate solution to (1.1) unless $\kappa(A) \leq 2^{t/2}$.

Several authors have suggested alternatives based on the orthogonal decomposition of A into

$$A = Q^T \begin{bmatrix} R \\ 0 \end{bmatrix}$$

where Q is orthogonal and R is $n \times n$ upper triangular. Writing

$$(1.2) \quad Q = \begin{bmatrix} Q_1 \\ Q_2 \end{bmatrix}$$

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where Q_1 is $n \times m$, the least squares solution is given by

$$(1.3) \quad \tilde{x} = R^{-1}Q_1b.$$

One approach of this kind was given by GOLUB [5] and BUSINGER and GOLUB [3] using Householder transformations on A and another employs the Gram-Schmidt algorithm and its modification as suggested by BJÖRCK [2] and by PETERS and WILKINSON [9]. Iterative refinements of these methods can be found in [2] and [6]. Further extensions of the Householder method were given by HANSON and LAWSON [7]. More recently an elimination approach has been suggested by CLINE [4].

In [10] an iterative procedure was suggested for approximating \tilde{x} , based on the splitting of the coefficient matrix A into

$$(1.4) \quad A = M - N,$$

where M and A have the same range. This method was developed further and convergence criteria similar to those for the nonsingular case were given in [1]. In this note the use of (1.4) is suggested as a possible direct approach to avoiding the ill-conditioned properties of the normal system (1.1).

2. A SPLITTING APPROACH

First notice that in the splitting (1.4), $M(M^TM)^{-1}M^TN = N$ since the range of N is contained in the range of M and since $M(M^TM)^{-1}M^T$ is the orthogonal projector on the range of M . Also, 1 is not an eigenvalue of $(M^TM)^{-1}M^TN$ since the null space of A is zero. Thus $I - (M^TM)^{-1}M^TN$ is nonsingular. In particular

$$A = M[I - (M^TM)^{-1}M^TN]$$

so that the least squares solution of $Ax = b$ is given by

$$(2.1) \quad \tilde{x} = [I - (M^TM)^{-1}M^TN]^{-1}(M^TM)^{-1}M^Tb,$$

since the right hand side of (2.1) reduces to $(A^TA)^{-1}A^Tb$.

One such choice of M is Q_1^T , where the matrix Q_1 is given in (1.2), for then $Q_1^T = AR^{-1}$ and $R = Q_1A$ since $Q_1Q_1^T = I$, and so A and Q_1^T have the same ranges. In this case $N = Q_1^T - A$ and (2.1) becomes

$$\begin{aligned} \tilde{x} &= [I - (Q_1Q_1^T)^{-1}Q_1(Q_1^T - A)]^{-1}(Q_1Q_1^T)^{-1}Q_1b = \\ &= [I - I + Q_1A]^{-1}Q_1b = R^{-1}Q_1b. \end{aligned}$$

Thus the familiar form given for \tilde{x} by (1.3) is a special case of (2.1).

Quite obviously the value of this splitting approach for a particular ill-conditioned least squares problem depends upon one's being able to choose the matrix M in such a way that its condition number is somewhat less than the condition number of A and with the property that the matrix $I - (M^T M)^{-1} M^T N$ is not difficult to invert. Theoretically the best choice of M is given by $M = Q_1^T$, as mentioned above. However in some cases M can be chosen by observation and by taking into account the structure of A . An example illustrating this situation is given in the next section.

3. EXAMPLE

Let ε be some small positive real number and consider the $m \times n$ matrix

$$A = \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ \varepsilon & 0 & 0 & \dots & 0 \\ 0 & \varepsilon & 0 & \dots & 0 \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & 0 & 0 & \dots & \varepsilon \end{bmatrix},$$

originally due to LÄUCHLI [8], with the resulting system $Ax = b$. The matrix $A^T A$ then has the form

$$A^T A = \begin{bmatrix} 1 + \varepsilon^2 & 1 & \dots & 1 \\ 1 & 1 + \varepsilon^2 & \dots & 1 \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot \\ 1 & 1 & \dots & 1 + \varepsilon^2 \end{bmatrix}.$$

As was pointed out in [5] and in [11, p. 135], if a solution of the normal system (1.1) is attempted with $\varepsilon^2 < \beta_0$, the machine precision, then $A^T A$ becomes the singular matrix of all 1's and the resulting least squares problem is rendered unsolvable. In particular, when ε is small and higher powers of ε are ignored the condition number of A is approximated by

$$\kappa(A) = \frac{\sqrt[n]{n}}{\varepsilon}.$$

Here one would like to choose M with the same range as A such that $\kappa(M)$ does not involve ε in the denominator. Let C^k denote the k^{th} column of an arbitrary matrix C . Then one can construct M with M^{j+1} independent of ε and having two nonzero entries by setting

$$M^{j+1} = \frac{1}{\varepsilon} (A^{j+1} - A^j)$$

for $j = 1, 2, \dots, n - 1$. Then taking $M^1 = A^1$,

$$M = \begin{bmatrix} 1 & 0 & 0 & \dots & 0 \\ \varepsilon & -1 & 0 & \dots & 0 \\ 0 & 1 & -1 & \dots & 0 \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & -1 \\ 0 & 0 & 0 & \dots & 1 \end{bmatrix}$$

with $\kappa(M) \approx \sqrt{n}$ and with A and M having the same range. In this case the matrix M is much better conditioned than A for small ε and $M^T M$ is the tri-diagonal matrix

$$(3.1) \quad M^T M = \begin{bmatrix} 1 + \varepsilon^2 & -\varepsilon & 0 & \dots & 0 & 0 \\ -\varepsilon & 2 & -1 & \dots & 0 & 0 \\ 0 & -1 & 2 & \dots & 0 & 0 \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & 0 & 0 & \dots & -1 & \\ 0 & 0 & 0 & \dots & -1 & 2 \end{bmatrix}.$$

Moreover, $I - (M^T M)^{-1} M^T N$ is the upper triangular matrix

$$(3.2) \quad (I - M^T M)^{-1} M^T N = \begin{bmatrix} 1 & 1 & 1 & \dots & 1 \\ 0 & \varepsilon & \varepsilon & \dots & \varepsilon \\ 0 & 0 & \varepsilon & \dots & \varepsilon \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ \cdot & \cdot & \cdot & \cdot & \cdot \\ 0 & 0 & 0 & \dots & \varepsilon \end{bmatrix}$$

Inversion of (3.1) and (3.2) and then post-multiplication by $(M^T M)^{-1} M^T b$ yields the least squares solution \tilde{x} to $Ax = b$. In practice, one determines \tilde{x} by elimination and back-substitution so that no matrix inversions are necessary here.

4. REMARKS

(a) It should be mentioned that the example in Section 3 is given only to show that it is sometimes possible to avoid the ill-conditioned problem associated with the normal system $A^T Ax = A^T b$ by a judicious choice of M in the splitting (1.4). In general the matrix M may not be so easy to obtain.

(b) Whenever A and M are $m \times n$ with rank n and have the same ranges, it follows that

$$(M^T A)^{-1} M^T = (A^T A)^{-1} A^T.$$

Thus \tilde{x} can be found by computing

$$(M^T A)^{-1} M^T b .$$

That this may not always be the best approach is well illustrated by the example in Section 3. Here $\varkappa(M^T A) \approx (A^T A)$ for this particular choice of M .

(c) In the case where the spectral radius of $(M^T M)^{-1} M^T N$ is less than one, the iteration

$$x^{(k+1)} = (M^T M)^{-1} M^T N x^{(k)} + (M^T M)^{-1} M^T b$$

converges to \tilde{x} for each $x^{(0)}$ [10]. In the example above the iteration converges if and only if $0 < \varepsilon < 2$. This iterative method may be useful in solving large, sparse least squares problems.

(d) If the linear system $Ax = b$ is under-determined and A has full row rank then one can compute the solution \tilde{y} of minimum norm in the following manner. The matrix A is split into $A = M - N$ where A and M have the same null spaces. Then \tilde{y} is given by

$$\tilde{y} = M^T (M M^T)^{-1} [I - N M^T (M M^T)^{-1}]^{-1} b .$$

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