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*Czechoslovak Mathematical Journal*, Vol. 30 (1980), No. 2, 332–335

Persistent URL: <http://dml.cz/dmlcz/101681>

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## JOIN GRAPHS OF TREES

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(Received October 10, 1978)

In [3] the intersection graph of a tree was defined. The intersection graph of a tree  $T$  is an undirected graph whose vertices are in a one-to-one correspondence with all proper subtrees of  $T$  and in which two vertices are adjacent if and only if the corresponding subtrees have a non-empty intersection.

Analogously a join graph of a tree can be defined. The join graph  $J(T)$  of a tree  $T$  is a graph whose vertices are in a one-to-one correspondence with all proper subtrees of  $T$  and in which two vertices are adjacent if and only if the join of the corresponding subtrees is not equal to  $T$ . (The join of two subtrees of  $T$  is the least subtree of  $T$  which contains both these subtrees.) A graph consisting only of one vertex is also considered a tree.

This definition was formulated in this form in order that the join graph might be a dual concept to the intersection graph. Nevertheless, in the following we shall study the complement  $\bar{J}(T)$  of  $J(T)$ , i.e. the graph in which two vertices are adjacent if and only if the join of the corresponding subtrees is equal to  $T$ .

We shall define some auxiliary concepts.

By  $IG(n)$  we shall denote the intersection graph of the family of all non-empty proper subsets of a set with  $n$  elements.

If  $G$  is an undirected graph without loops and multiple edges, let  $\sim$  be a binary relation on the vertex set of  $G$  such that  $a \sim b$  if and only if  $\Gamma(a) = \Gamma(b)$ . (By the symbol  $\Gamma(x)$  we denote the set of all vertices of  $G$  which are adjacent to  $x$ .) Evidently this relation is an equivalence; we shall call it the *adjacency equivalence* on  $G$ . If  $a \sim b$ ,  $a \neq b$ , then  $a$  and  $b$  are not adjacent in  $G$ ; otherwise we should have  $b \in \Gamma(a)$  which would imply  $b \in \Gamma(b)$ , i.e. the existence of a loop. If  $a \sim a'$ ,  $b \sim b'$ , then evidently  $a$  is adjacent to  $b$  if and only if  $a'$  is adjacent to  $b'$ .

If we identify all adjacency-equivalent pairs of vertices of  $G$ , we obtain a graph  $A(G)$  which will be called the *adjacency reduct* of  $G$ . Evidently there exists a discrete homomorphism [2] of  $G$  onto  $A(G)$ .

**Lemma 1.** *Let  $M$  be a set with  $n$  elements. Let  $G$  be a graph whose vertices are in a one-to-one correspondence with all non-empty proper subsets of  $M$  and in which two vertices are adjacent if and only if the union of the corresponding sets is equal to  $M$ . Then  $G$  is isomorphic to the complement of  $IG(n)$ .*

*Proof.* Let  $\varphi$  be a mapping of the vertex set of  $G$  onto the vertex set of the complement of  $IG(n)$  such that if  $x$  is a vertex of  $G$  corresponding to the subset  $X$  of  $M$ , then  $\varphi(x)$  is the vertex of  $IG(n)$  corresponding to the set  $M - X$ ; this mapping is evidently a bijection. By De Morgan's formula we have  $X \cup Y = M$  if and only if  $(M - X) \cap (M - Y) = \emptyset$ . Therefore vertices  $\varphi(x), \varphi(y)$  are adjacent in the complement of  $IG(n)$  if and only if  $x, y$  are adjacent in  $G$  and  $\varphi$  is an isomorphism.

**Theorem 1.** *Let  $T$  be a finite tree with  $n \geq 3$  vertices and with  $k$  terminal vertices, let  $J(T)$  be its join graph, let  $\bar{J}(T)$  be the complement of  $J(T)$ . Then the adjacency reduct of  $\bar{J}(T)$  is isomorphic to the complement of  $IG(k)$  with one isolated vertex added.*

*Proof.* Let  $K$  be the set of all terminal vertices of  $T$ . If  $L \subset K$ , then we denote by  $\mathcal{T}(L)$  the family of all subtrees of  $T$  which contain all vertices of  $L$  and no vertex of  $K - L$ . Now let  $T_1, T_2$  be two proper subtrees of  $T$ , let  $L_1, L_2$  be such subsets of  $K$  that  $T_1 \in \mathcal{T}(L_1), T_2 \in \mathcal{T}(L_2)$ . If  $L_1 \cup L_2 = K$ , then the join of  $T_1$  and  $T_2$  contains all terminal vertices of  $T$  and this is possible if and only if this join is equal to  $T$ . If  $L_1 \cup L_2 \neq K$ , let  $x \in K - (L_1 \cup L_2)$ . The vertex set of the join of  $L_1$  and  $L_2$  consists of all vertices of  $L_1$  and  $L_2$  and eventually also of all inner vertices of a path in  $T$  connecting a vertex of  $L_1$  with a vertex of  $L_2$ . The vertex  $x$  belongs neither to  $L_1$  nor to  $L_2$  and, being a terminal vertex of  $T$ , it cannot be an inner vertex of a path in  $T$ . Therefore the join of  $T_1$  and  $T_2$  does not contain  $x$  and is not equal to  $T$ . We see that the vertices of  $\bar{J}(T)$  corresponding to  $T_1$  and  $T_2$  are adjacent if and only if  $L_1 \cup L_2 = K$ . Evidently if  $T_1 \in \mathcal{T}(L_1), T_2 \in \mathcal{T}(L_2)$ , then  $T_1 \sim T_2$  (as vertices of  $\bar{J}(T)$ ) if and only if  $L_1 = L_2$ . The adjacency reduct of  $\bar{J}(T)$  is isomorphic to the graph whose vertices are all proper subsets of  $K$  and in which two vertices are adjacent if and only if their union is  $K$ . By Lemma 1 the vertices corresponding to non-empty proper subsets of  $K$  form a subgraph of this graph isomorphic to the complement of  $G(k)$ . The vertex corresponding to the empty set is isolated.

We see that every further information about  $T$  from  $J(T)$  can be obtained only from the numbers of vertices of the classes of the adjacency equivalence.

**Lemma 2.** *Let the graph  $IG(n)$  be given for some  $n$ . Then for each vertex  $x$  of  $IG(n)$ , we can determine the cardinality of the set to which  $x$  corresponds.*

*Proof.* The vertices corresponding to one-element subsets of the set  $M$  (from the definition of  $IG(n)$ ) form an independent set in  $IG(n)$  of the cardinality  $n$ . No other family of  $n$  pairwise disjoint non-empty subsets of  $M$  can exist. Therefore a vertex  $x$  of  $IG(n)$  corresponds to a one-element subset of  $M$  if and only if it belongs to the

(unique) independent set  $M_0$  of the greatest cardinality. If  $2 \leq m \leq n - 1$ , a vertex  $y \notin M_0$  corresponds to a set of the cardinality  $m$  if and only if it is adjacent to exactly  $m$  vertices of  $M_0$  (i.e., the corresponding set has non-empty intersections with exactly  $m$  one-element sets).

**Theorem 2.** *Let  $T$  be a finite tree, let  $J(T)$  be its join graph. Given  $J(T)$ , we can reconstruct  $T$  up to isomorphism.*

*Proof.* We construct the adjacency reduct of  $\bar{J}(T)$  and its complement; by Theorem 1 it is isomorphic to  $IG(k)$ , therefore we can determine the number  $k$  of terminal vertices of  $T$ . By Theorem 1 and Lemma 2 we can determine the classes of the adjacency equivalence in  $J(T)$  which correspond to families  $\mathcal{F}(L)$  for  $L$  of the cardinality 1; to these families we assign vertices  $t_1, \dots, t_k$  in a one-to-one way. These vertices  $t_1, \dots, t_k$  are the terminal vertices of  $T$ ; the cardinality of the class corresponding to  $t_i$  for  $i = 1, \dots, k$  is the number of the subtrees of  $T$  which contain  $t_i$  and no other terminal vertex of  $T$ . Thus we take  $K = \{t_1, \dots, t_k\}$ . If  $L$  is a subset of  $K$  with  $l \geq 2$  vertices, then  $\mathcal{F}(L)$  is the class of all vertices of  $\bar{J}(T)$  which are adjacent to all vertices from the classes corresponding to vertices of  $L$  and not adjacent to vertices from the classes corresponding to vertices of  $K - L$ . Therefore for each proper subset  $L$  of  $K$  we can determine the number  $\mu(L)$  of subtrees of  $T$  which contain all vertices of  $L$  and no vertex of  $K - L$ .

If  $T$  is a snake (a tree consisting of one simple path), we have  $|K| = 2$ , therefore  $K = \{t_1, t_2\}$  and there are three proper subsets of  $K$ , namely  $\{t_1\}$ ,  $\{t_2\}$  and  $\emptyset$ . Then  $\bar{J}(T)$  consists of a complete bipartite graph with some isolated vertices added (by Theorem 1). Therefore we can conclude that  $T$  is a snake; otherwise  $\bar{J}(T)$  would contain a triangle. We determine  $\mu(\{t_1\})$  and  $\mu(\{t_2\})$ ; evidently  $\mu(\{t_1\}) = \mu(\{t_2\}) = n - 1$ , where  $n$  is the number of vertices of  $T$ . By its number of vertices a snake is determined up to isomorphism.

If  $T$  is not a snake, we proceed by induction with respect to the number  $n$  of vertices of  $T$ . For  $n = 2$  and  $n = 3$  the tree is always a snake and for this case the assertion was proved. Let  $n_0 \geq 4$ . Suppose that the assertion is true for  $n = n_0 - 1$  and prove it for  $n = n_0$ .

Let  $T$  be a tree with  $n_0$  vertices, let its join graph  $J(T)$  be given. For  $T$  let  $K = \{t_1, \dots, t_k\}$ . Let  $T'$  be the tree obtained from  $T$  by deleting  $t_k$ ; this is a tree with  $n_0 - 1$  vertices. Let  $t'$  be the vertex adjacent to  $t_k$  in  $T$ . Distinguish two cases:

- (i) The vertex  $t'$  is a terminal vertex of  $T'$ .
- (ii) The vertex  $t'$  is not a terminal vertex of  $T'$ .

In the case (ii) there exists a proper subset  $L$  of  $K - \{t_k\}$  such that the least tree from  $\mathcal{F}(L)$  contains  $t'$  and thus all trees from  $\mathcal{F}(L)$  contain  $t'$ . Then there is a one-to-one correspondence between  $\mathcal{F}(L)$  and  $\mathcal{F}(L \cup \{t_k\})$ ; each tree from  $\mathcal{F}(L)$  is obtained from a tree from  $\mathcal{F}(L \cup \{t_k\})$  by deleting the vertex  $t_k$  and the edge  $t't_k$  and each tree from  $\mathcal{F}(L \cup \{t_k\})$  is obtained from a tree from  $\mathcal{F}(L)$  by adding this vertex and this edge. Therefore  $\mu(L) = \mu(L \cup \{t_k\})$ . In the case (i), to each proper

subset  $L$  of  $K - \{t_k\}$  there exist trees from  $\mathcal{T}(L)$  which contain  $t'$  and their number is equal to  $\mu(L \cup \{t_k\})$ , but there are also trees from  $\mathcal{T}(L)$  which do not contain  $t'$ ; therefore  $\mu(L) > \mu(L \cup \{t_k\})$  for each  $L \subset K - \{t_k\}$ . We see that we are able to recognize whether (i) or (ii) occurs.

Consider the case (i). In the tree  $T'$  we determine the classes  $\mathcal{T}'(L)$  and numbers  $\mu'(L)$  analogous to  $\mathcal{T}(L)$  and  $\mu(L)$  for all proper subsets  $L$  of  $K' = (K \cup \{t'\}) - \{t_k\}$ . If  $L$  is a proper subset of  $K'$  and  $t' \notin L$ , then a tree from  $\mathcal{T}(L)$  belongs to  $\mathcal{T}'(L)$  if and only if it does not contain  $t'$ . As was shown above, the number of trees from  $\mathcal{T}(L)$  containing  $t'$  is equal to  $\mu(L \cup \{t_k\})$ . Therefore  $\mu'(L) = \mu(L) - \mu(L \cup \{t_k\})$ . If  $t' \in L$ , then we can prove analogously as above that there is a one-to-one correspondence between  $\mathcal{T}'(L)$  and  $\mathcal{T}((L - \{t'\}) \cup \{t_k\})$  for  $L \neq \{t'\}$ , therefore in this case  $\mu'(L) = \mu((L - \{t'\}) \cup \{t_k\})$ . For  $L = \{t'\}$  there is such a correspondence between  $\mathcal{T}'(L)$  and the set obtained from  $\mathcal{T}((L - \{t'\}) \cup \{t_k\})$  by deleting the one-vertex tree consisting of  $t_k$ ; thus  $\mu'(\{t'\}) = \mu(\{t_k\}) - 1$ . Hence we can transform  $J(T)$  to  $J(T')$ : By the induction assumption the tree  $T'$  can be reconstructed from  $J(T')$ . In the reconstructed tree  $T'$  we find the vertex  $t'$  and add the vertex  $t_k$  and the edge  $t't_k$  to  $T'$  to obtain  $T$ .

Now consider the case (ii). We put  $K' = K - \{t_k\}$ . If  $L$  is a proper subset of  $K'$ , then each tree from  $\mathcal{T}(L)$  is also in  $T'$  and thus  $\mu'(L) = \mu(L)$  for each such  $L$ . We reconstruct the tree  $T'$ . Now it is more difficult to find  $t'$  in  $T'$ , because  $t'$  is not a terminal vertex of  $T'$ . Let  $u_1, u_2$  be two elements of  $K'$ . If  $t'$  lies between  $u_1$  and  $u_2$ , then each tree from  $\mathcal{T}(\{u_1, u_2\})$  contains  $t'$  and analogously to the above considerations we can prove  $\mu(\{u_1, u_2\}) = \mu(\{u_1, u_2, t_k\})$ . In the opposite case  $\mu(\{u_1, u_2\}) > \mu(\{u_1, u_2, t_k\})$ . Therefore we can determine for any two vertices of  $K'$  whether  $t'$  lies between them or not. If there exist three vertices  $u_1, u_2, u_3$  of  $K'$  such that  $t'$  lies between any two of them, then  $t'$  is uniquely determined [1]. If not, then there is an equivalence on  $K'$  such that two vertices are in this equivalence if and only if  $t'$  does not lie between them and this equivalence has exactly two classes  $L_1, L_2$ . Let  $T_1$  (or  $T_2$ ) be the least tree from  $\mathcal{T}(L_1)$  (or  $\mathcal{T}(L_2)$ , respectively). The trees  $T_1, T_2$  are disjoint and there exists a path  $P$  in  $T'$  connecting a vertex  $v_1$  of  $T_1$  with a vertex  $v_2$  of  $T_2$  whose inner vertices belong neither to  $T_1$  nor to  $T_2$ . One of those inner vertices is  $t'$ . If  $d$  is the distance between  $v_1$  and  $t'$ , then  $\mathcal{T}(L_1)$  contains exactly  $d$  trees not containing  $t'$  and  $\mu(L_1 \cup \{t_k\})$  trees containing  $t'$ . This yields  $d = \mu(L_1) - \mu(L_1 \cup \{t_k\})$  and  $t'$  is determined. We add the vertex  $t_k$  and the edge  $t't_k$  to  $T'$  and  $T$  is reconstructed.

#### References

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