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ALGORITMY

31. PERMUT

SIMPLE ALGORITHM GENERATING ALL PERMUTATIONS

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The procedure *PERMUT* generates all permutations of the sequence $a_m, a_{m+1}, \dots, a_{m+n-1}$. In every step of its work, it generates a permutation, assigned for the variables $a[m], \dots, a[m+n-1]$. The first permutation is the original one, entering into the algorithm; after generating all the permutations, the procedure *PERMUT* restitutes the original order of values in the array a . Any permutation can be processed by the procedure *P* immediately after it has been generated. Besides the permutations, the number of their inversions is generated and assigned for the variable r^*). If we do not need it, we can simplify the algorithm by omitting all statements where the variable r occurs.

```

procedure PERMUT( $a, m, n, p$ );
value  $m, n$ ;
array  $a$ ; integer  $m, n$ ; procedure  $P$ ;
begin
    integer  $i, k, r$ ; integer array  $j[1 : n - 1]$ ; real  $b$ ;
    for  $i := 1$  step 1 until  $n - 1$  do  $j[i] := i + m; r := 0$ ;
     $G : P; i := n - 1$ ;
     $H : \text{if } i \leq 0 \text{ then go to } D$ ;
    if  $j[i] = m$  then go to  $F$ ;
     $b := a[j[i]]$ ;  $a[j[i]] := a[j[i] - 1]$ ;  $j[i] := j[i] - 1$ ;  $a[j[i]] := b$ ;
     $r := r + 1$ ; go to  $G$ ;
     $F : j[i] := j[i] + i; b := a[m]$ ;
    for  $k := m$  step 1 until  $m + i - 1$  do  $a[k] := a[k + 1]$ ;
     $a[m + i] := b; r := r - i; i := i - 1$ ; go to  $H$ ;
     $D :$ 
end PERMUT;
    
```

*) The number of inversions in a permutation $a_{j_1}, a_{j_2}, \dots, a_{j_n}$ is defined commonly as the number of all pairs of indices $(i, k), 1 \leq i, k \leq n$, where $i < k, j_i > j_k$.

The array a may be of course of the type integer, in which case the variable b is to be declared integer too.

Example. Calling $PERMUT(a, 1, 4, p)$ where $a_i = i$ and P prints the generated permutation, followed by the number of its inversions in parentheses:

1234(0), 1243(1), 1423(2), 4123(3), 1324(1), 1342(2), 1432(3), 4132(4),
3124(2), 3142(3), 3412(4), 4312(5), 2134(1), 2143(2), 2413(3), 4213(4),
2314(2), 2341(3), 2431(4), 4231(5), 3214(3), 3241(4), 3421(5), 4321(6).

An interesting property of the algorithm may be seen which holds not only for the given example, but generally: the sequence of generated permutations has certain symmetrical attributes, namely that the sum of the pair of the numbers of the inversions of the i -th and $(n! - i + 1)$ -th permutations is equal to $\binom{n}{2}$, and that the order of the i -th permutation is exactly inverse to the order of the $(n! - i + 1)$ -th one ($i = 1, \dots, n!$).

The algorithm has been programmed and tested in the small computer ODRA 1013 [1] in the programming language MOST [2].

References

- [1] Černý, V., Pár J.: Programmer's Manual on Automatic Computer ODRA 1013 (in Czech). Kancelářské stroje n. p., Hradec Králové 1967.
- [2] Szczepkiewicz, J.: Programming in the autocode MOST 1 (in Polish). ELWRO Publication 03—VI—1, Wrocław.