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ALGORITMY

36. SNEDECOR

AN ALGORITHM FOR FISHER - SNEDECOR'S F -TEST WITHOUT APPLICATION OF CRITICAL VALUES

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The algorithm suggested in this paper computes the probability that Fisher-Snedecor's test statistic will exceed the value F actually observed, i.e.

$$(1) \quad \alpha_{m,n}(F) = \frac{\left(\frac{m}{n}\right)^{m/2}}{B\left(\frac{m}{2}, \frac{n}{2}\right)} \int_F^\infty y^{m/2-1} \left(1 + \frac{m}{n} y\right)^{-(m+n)/2} dy,$$

where m, n is the pair of numbers of degrees of freedom. This probability may be hence called the significance degree, similarly as in the case of Student's t -statistic treated in our previous paper [1]. Since the latter represents a special case of the present problem (with $m = 1, F = t^2$), the features of the algorithm and further remarks made in [1] apply here, too (except the distinction between one-sided and two-sided tests) and will not be repeated.

Analogously as in [1], the relation

$$(2) \quad \alpha_{m,n}(F) = A_{m,n}(x)$$

with

$$(3) \quad x = \left(1 + \frac{m}{n} F\right)^{-1},$$

$$(4) \quad A_{m,n}(x) = \frac{1}{B\left(\frac{m}{2}, \frac{n}{2}\right)} \int_0^x y^{n/2-1} (1 - y)^{m/2-1} dy$$

holds and the algorithm is based on the following recurrence relations and initial conditions

$$(5) \quad A_{m,n}(x) = A_{m,n-2}(x) - \frac{\Gamma\left(\frac{m+n}{2} - 1\right)}{\Gamma\left(\frac{m}{2}\right) \Gamma\left(\frac{n}{2}\right)} x^{n/2-1} (1-x)^{m/2}$$

$$(m > 0, n > 2),$$

$$(6) \quad A_{m,n}(x) = A_{m-2,n}(x) + \frac{\Gamma\left(\frac{m+n}{2} - 1\right)}{\Gamma\left(\frac{m}{2}\right) \Gamma\left(\frac{n}{2}\right)} x^{n/2} (1-x)^{m/2-1}$$

$$(m > 2, n > 0),$$

$$(7) \quad A_{1,1}(x) = \frac{2}{\pi} \arcsin \sqrt{x}, \quad A_{2,1}(x) = \sqrt{x}, \quad A_{2,2}(x) = x,$$

$$(8) \quad A_{m,n}(x) = 1 - A_{n,m}(1-x).$$

The last relation follows from the definition, the remainder can be proved by differentiation (the relations (5) and (6) being equivalent due to (8)).

If the statistic F has its usual form

$$(9) \quad F = \left(\frac{y}{m}\right) / \left(\frac{z}{n}\right),$$

where y, z are certain sampling characteristics, then the transform (3), which is used instead of the statistic F , may be evaluated directly from y, z in the form

$$(10) \quad x = z / (y + z).$$

```

real procedure SNEDECOR( $x, m, n$ ); value  $x$ ; real  $x$ ; integer  $m, n$ ;
begin
  real  $a, b, c, d, e, f$ ; integer  $i$ ;
  procedure G;
    begin  $c := c \times x$ ;
      for  $f := e$  step 2 until  $i$  do
        begin  $a := a + b$ ;  $d := b \times c$ ;  $b := d/f$ ;  $c := c + 2 \times x$  end  $f$ 
      end G;
  procedure H;
    begin  $x := 1 - x$ ; G;  $b := -d$ ;  $c := i + 1$ ;  $e := 3$ ;  $i := n$ ;
       $x := 1 - x$ ; G
    end H;
  procedure P; begin  $b := \text{sqrt}(x)$ ;  $c := 1$ ; H end;

```

```

procedure Q; begin b := 1; c := n; G; a :=  $a \times (1 - x) \uparrow (n \div 2)$  end;
if  $n > (n \div 2) \times 2$ 
  then
    begin i := m;
      if  $m > (m \div 2) \times 2$ 
        then
          begin a := 0.63661977  $\times$  arcsin(sqrt(x));
            b := 0.63661977  $\times$  sqrt((1 - x)  $\times$  x); e := 2;
            d := b; e := 3; H
          end
        else begin a := 0; e := 2; P end
      end
    else
      begin a := 0; e := 2;
        if  $m > (m \div 2) \times 2$ 
          then
            begin i := n; n := m; m := i; x := 1 - x; P; x := 1 - x;
              n := m; m := i; a := 1 - a
            end
          else
            if  $m > n$ 
              then
                begin i := n; n := m; Q; n := i; a := 1 - a end
              else
                begin i := m; x := 1 - x; Q; x := 1 - x end
              end
            end;
          SNEDECOR := a
        end SNEDECOR
      end
    end
  end

```

The result is obtained with the accuracy of at least about 5 decimal places. We give some check values:

```

SNEDECOR (0.3, 1, 1) = 0.36901
SNEDECOR (0.25, 1, 10) = 0.00027
SNEDECOR (0.75, 1, 19) = 0.02099
SNEDECOR (0.5, 4, 10) = 0.10937
SNEDECOR (0.4, 10, 6) = 0.58010
SNEDECOR (0.7, 3, 8) = 0.38890
SNEDECOR (0.6, 4, 9) = 0.28109
SNEDECOR (0.1, 3, 1) = 0.39582
SNEDECOR (0.2, 5, 11) = 0.00143
SNEDECOR (0.3, 7, 3) = 0.55292
SNEDECOR (0.75, 10, 1) = 0.99973

```

The program has been tested in the symbolic language MOST [3] and implemented in the Biophysical Institute, Faculty of General Medicine, Charles University Prague for the computer ODRA 1013 [4].

- [1] *Režný, Z., Jirkovský, J.:* STUDENT. An algorithm for Student's *t*-test without application of critical values, *Aplikace matematiky* 19 (1974), 133–135.
- [2] *Janko, J.:* Statistical Tables (in Czech), Publ. House of the Czechoslovak Acad. of Sci., Prague 1958.
- [3] *Szczepkiewicz, J.:* Programming in the autocode MOST I (in Polish), Elwro Publication 03-VI-1, Wrocław.
- [4] *Černý, V., Půr, J.:* Programmer's Manual on Automatic Computer ODRA 1013 (in Czech), Kanc. stroje n. p., Hradec Králové 1967.