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TWIN PRIME PROBLEM IN AN ARITHMETIC WITHOUT INDUCTION

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Abstract: We prove that the twin prime problem is undecidable in a first-order arithmetic without induction, stronger than Robinson's arithmetic.

Key words: First-order arithmetic without induction, twin prime problem, undecidable.

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Introduction. In this paper we prove that the twin prime problem is undecidable in certain first-order arithmetic Ar without induction.

Moreover, our Ar will be stronger than Robinson's arithmetic (but weaker than Peano one). We will present a parametrical construction of a substructure of a fixed non-standard model \mathcal{U} of Peano arithmetic. As parameters we will have a submodel of Ar and a non-standard element of \mathcal{U} . The required models are obtained by an appropriate choice of parameters.

§ 0. Preliminaries

0.0.0. Let L be a first-order language with a binary predicate $<$. Let $\varphi(x)$ be a formula of L . We denote by $(\exists x)\varphi(x)$ the formula $(\forall y)(\exists x)(y < x \ \& \ \varphi(x))$,

where y is not a variable of φ . Let \mathcal{U} and \mathcal{L} be structures for L . By $\mathcal{U} \subset \mathcal{L}$ ($\mathcal{U} < \mathcal{L}$) we mean that \mathcal{U} is a substructure of \mathcal{L} (\mathcal{U} is an elementary substructure of \mathcal{L}). The language obtained from L by adding all the names a of individuals a of \mathcal{U} is denoted by $L(\mathcal{U})$. We expand \mathcal{U} to a structure $\underline{\mathcal{U}}$ for $L(\mathcal{U})$ as follows: if \underline{a} is the name of an individual a of \mathcal{U} then $\underline{\mathcal{U}}$ assigns a to \underline{a} . Let M be a nonempty subset of \mathcal{U} (where $\mathcal{U} = A$ is the universe of \mathcal{U}). If there is a substructure of $\underline{\mathcal{U}}$ with universe M then it is designated by \mathcal{U}/M .

The expression $\mathcal{U} \subset \mathcal{L}$ ($\mathcal{U} \leq \mathcal{L}$) stands for 1) $\mathcal{U} \subset \mathcal{L}$ ($\mathcal{U} < \mathcal{L}$), 2), if $a \in A$ and $b \in B$, then $a \leq b$. (\mathcal{L} is an (elementary) end-extension of \mathcal{U} .) Writing $\mathcal{U} \subset \mathcal{L}$ we mean that $\mathcal{U} \subseteq \mathcal{L}$ and $A \neq B$. (\mathcal{L} is a proper end-extension of \mathcal{U} .) $\mathcal{U} < \mathcal{L}$ is defined analogously.

0.1.0. The language J of Peano arithmetic P is $\langle 0', +, \cdot, < \rangle$. Let \mathcal{N} be the standard model of P . For $n \in \mathbb{N}$ we denote by n the constant term $0' \dots'$, where $'$ is applied n -times.

i, j, k, l, m, n are variables for elements of \mathbb{N} .

Remark. We work in the logic with equality.

0.1.1. Let $s(i)$, $i = 1, \dots, 5$ be symbols such that $s(1)$ is the binary predicate $x | y$, $s(2)$ is the unary predicate $\text{Prm}(x)$, $s(3)$ is the unary predicate $\text{Prm}_2(x)$, $s(4)$ is the binary function $e(x, y)$, and $s(5)$ is the binary function $r(x, y)$.

Let $\mathcal{G}_i, i = 1, 2, 3, 4, 5$ be the following formulas:

\mathcal{G}_1 is the formula $(\exists z)(y = x.z)$, \mathcal{G}_2 is the formula $y | x \rightarrow (y = \bar{1} \vee y = x)$, \mathcal{G}_3 is $\text{Prm}(x) \& \text{Prm}(x + \bar{2})$,
 \mathcal{G}_4 is $(x > 0 \& y > \bar{1} \& y^2 | x \& y^{z+1} \nmid x) \vee ((x = 0 \vee y \leq \bar{1}) \& z = 0)$,
 \mathcal{G}_5 is $(x > 0 \& y > \bar{1} \& (\exists u)(u = e(x, y) \& x = y^u.z)) \vee \vee ((x = 0 \vee y \leq \bar{1}) \& z = 0)$.

Remark. By $x \nmid y$ we mean $\neg(x | y)$.

Let P designate also the theory obtained from P by adding the functions x^y and the symbols $s(i)$ defined by $\mathcal{G}_i, i = 1, \dots, 5$.

0.1.2. Throughout the paper, $\mathcal{M}_0, \mathcal{U}_0, \mathcal{U}_1, \mathcal{U}$ are non-standard models of P such that

$$\mathcal{N} \leq \mathcal{M}_0 \leq \mathcal{U}_0 \leq \mathcal{U}_1 \leq \mathcal{U}$$

and α is a fixed element of $A - A_1$. We use McDowell-Specker's theorem. (See [1].)

If there is no danger of confusion, we write $+, \cdot, <$ etc. instead of $+^{\mathcal{U}}, \cdot^{\mathcal{U}}, <^{\mathcal{U}}$ etc.

Let \mathcal{U}^* be "integers over \mathcal{U} ". \mathcal{U}^* is an ordered domain. If a, b are elements of \mathcal{A}^* , $-a$ designates the inverse element of a . $a - b$ designates $a + (-b)$, and $|a|$ designates absolute value of a . If $b | a$, we denote by $\frac{a}{b}$ the element c with $a = b.c$. For $B \subseteq A$, we put $B^- = \{-a; a \in B\}$ and $B^* = B^- \cup B$. If $\mathcal{L} \subseteq \mathcal{U}$ and $\mathcal{L} \models x < y \rightarrow \rightarrow (\exists z)(z \neq 0 \& x + z = y)$ then $\mathcal{L}^* = \mathcal{U}^*/B$ is a subdomain of \mathcal{U}^* .

§ 1. Arithmetic Ar and some models of it

1.0.0. Ar is a first-order theory with the language

J. The nonlogical axioms of Ar are the following:

(a) $x + 0 = x$	$x \cdot 0 = 0$
$x + y = y + x$	$x \cdot y = y \cdot x$
$x + (y + z) = (x + y) + z$	$x \cdot (y \cdot z) = (x \cdot y) \cdot z$
$x + y' = (x + y)'$	$x \cdot y' = x \cdot y + z$
$x \cdot (y + z) = x \cdot y + x \cdot z$	

- (b) 1) $\neg(x < x)$
 2) $x < y \& y < z \rightarrow x < z$
 3) $x < y \vee x = y \vee y < x$
 4) $x < y' \leftrightarrow x < y \vee x = y$
 5) $0 < x \vee 0 = x$
 6) $0 < x \rightarrow (\exists y)(y' = x)$
 7) $x < y \leftrightarrow (\exists z \neq 0)(x + z = y)$

(c) $x < y \& 0 < u \leq v \rightarrow x + u < y + v \& x \cdot u < y \cdot v$

(d) (schema) $\{ \sigma_n; n \in \mathbb{N} - \{0\} \}$,

where σ_n is the formula $(\forall x)(\exists y < x)(\exists z < \bar{n})(x + y \cdot \bar{n} + z)$.

1.0.1. Proposition. The following sentences are provable in Ar:

- (i) $x \neq 0 \rightarrow (\exists y)(\forall z)(y < x \& z < x \rightarrow z \neq y)$,
- (ii) $x < y \rightarrow x' < y'$,
- (iii) $x' = y' \rightarrow x = y$,
- (iv) $x < y \rightarrow x \neq y$.

1.0.2. Let Ar designate also the theory obtained from Ar by adding the symbols $s(i)$ defined by $\langle g_i, i = 1, 2, 3 \rangle$.

1.1.0. Let \mathcal{M}_1 be a model of Ar such that

$$\mathcal{A}_0 \subseteq \mathcal{M}_1 \subseteq \mathcal{A}_1$$

Let $s \in \mathcal{A}_0$.

We define, for $i = 0, 1$,

$M_{1i}[s] = \{ \alpha^k a_k + \dots + \alpha a_1 + a_0; k \in \mathbb{N} - \{0\}, a_1, \dots$

$\dots, a_k \in M_1^*, a_k > 0, a_0 \in M_1^*,$

there exists an $e \in A_0 - \mathbb{N}$ such that $s^e \mid \mathcal{M}_1^* a_1, \dots$

$\dots, s^e \mid \mathcal{M}_1^* a_k \},$

$M_{1i}(s) = M_{1i}[s] \cup M_i.$

Lemma. Let $a \in M_{1i}$, $i = 0, 1$. Then there is precisely one $k \in \mathbb{N}$ and $a_1, \dots, a_k \in M_1^*$, $a_k > 0$, $a_0 \in M_1^*$ such that

$$a = \alpha^k a_k + \dots + \alpha a_1 + a_0.$$

Proof is obvious.

Notation. For $a \in M_{1i}[s]$, $i = 0, 1$, we denote by $v(a)$ the standard number k and by a_1, \dots, a_k elements of M_1^* , $a_k > 0$, and a_0 element of M_1^* such that $a = \alpha^k a_k + \dots + \alpha a_1 + a_0$.

Lemma. $M_{1i}(s)$ is the universe of a substructure of \mathcal{U} $i = 0, 1$.

Proof. Let $a, b \in M_{1i}[s]$. Obviously $a' \in M_{1i}[s]$. Let $v(a) \leq v(b)$. For $0 \leq i \leq v(a)$ we have $(a + b)_i = a_i + b_i$, for $v(a) < i \leq v(b)$ we have $(a + b)_i = b_i$. There is an $e \in A_0 - \mathbb{N}$ such that $s^e \mid \mathcal{M}_1^* a_i, i = 1, \dots, v(a)$, $s^e \mid \mathcal{M}_1^* b_i, i = 1, \dots, v(b)$. Consequently, $a + b \in M_{1i}[s]$. We also have $(a \cdot b) = \sum_{k+l=i} a_k b_l$; for $i \geq 1$ we have $s^e \mid \mathcal{M}_1^* \sum_{k+l=1} a_k b_l$. Thus, $a \cdot b \in M_{1i}[s]$. Similarly for $a \in M_i$ and $b \in M_{1i}[s]$ etc.

1.1.1. We put $\mathcal{M}_{1i}(s) = \mathcal{U} / M_{1i}(s)$, $i = 0, 1$. We write \mathcal{M}_{1i} for $\mathcal{M}_{1i}(s)$, $i = 0, 1$.

1.1.2. Theorem. Let $n \mid s$ for every $n \in \mathbb{N}$. Then $\mathcal{M}_{1i}(s) \models Ar$, $i = 0, 1$.

Proof. We have $\mathcal{M}_{1i} \subseteq \mathcal{U}$. Only the axioms (b6), (b7) and the schema (d) are not general closures of open formulas and, consequently it suffices to prove that \mathcal{M}_{1i} is a model of these axioms. Obviously $\mathcal{M}_{1i} \models (b6)$. We will prove $\mathcal{M}_{1i} \models (b7)$. Let $a, b \in M_{1i}[s]$ and $a < b$. Thus $v(a) \leq v(b)$. If $v(a) = v(b)$, put $j = \max \{i; a_i \neq b_i\}$. If $b_j - a_j < 0$, then we have $\alpha^j(b_j - a_j) + \dots + (b_0 - a_0) \leq -\alpha^j + \alpha^{j-1}|b_{j-1} - a_{j-1}| + \dots + |b_0 - a_0| \leq -\alpha^j + \alpha^{j-1} \cdot j \cdot \max \{|b_i - a_i|; i = 0, \dots, j-1\} < 0$. Thus $b_j - a_j > 0$. On the other hand, if $v(a) < v(b)$ then obviously $b - a \in M_{1i}[s]$. Thus $\mathcal{M}_{1i} \models (b7)$. It remains to prove the schema (d). Let $n \in \mathbb{N}$, $n > 0$, $a \in M_{1i}[s]$, $k = v(a)$. There are $\tilde{a}_0 \in M_1^*$, $\tilde{a}_0 \in M_1^*$ such that $0 \leq \tilde{a}_0 < n$ and $a_0 = n \cdot \tilde{a}_0 + \tilde{a}_0$.

Put $b = \alpha^k \cdot \frac{a_k}{n} + \dots + \alpha \cdot \frac{a_1}{n} + \tilde{a}_0$. There exists an $e \in A_0 - \mathbb{N}$ such that $s^e \mid \mathcal{M}_1^* a_i, \frac{a_i}{n} \in M_1^*$ and $s^{e-1} \mid \mathcal{M}_1^* \frac{a_i}{n}, i = 1, \dots, k$. Consequently, $b \in M_{1i}[s]$. Evidently $a = n \cdot b + \tilde{a}_0$. Hence $\mathcal{M}_{1i} \models \omega_n$.

1.2.0. Let $M \subseteq |\mathcal{U}|$, $a \in M$. We say that a is decomposable in M if there are $b, c \in M$ such that $a = b \cdot c$.

1.2.1. Lemma. Let $a \in M_{1i}[s]$, $a_0 \in \{-1, 1\}$, $v(a) \geq 2$. Then a is decomposable in $M_{1i}[s]$, $i = 0, 1$.

Proof. $a_0 = 1$. Let $d, e \in A_0 - \mathbb{N}$, $e < d$, $\hat{a}_i \in M_1^*$, $a_i = \hat{a}_i \cdot s^{d+e}$, $i = 1, \dots, k$, $k = v(a)$. Let $x_0 = y_0 = 1$, $x_1 = s^e$ and $y_{i+1} = a_{i+1} - y_i \cdot s^e$ if $0 \leq i < k-1$ and $y_{k-1} = \hat{a}_k \cdot s^d$.

Obviously, $\frac{y_i}{s^e} \in M_1^*$, $i = 1, \dots, k-1$. Thus, $y = \alpha^{k-1} \cdot y_{k-1} + \dots + 1 \in M_{11}[s]$, $x = \alpha \cdot s^e + 1 \in M_{11}[s]$. We have $(x \cdot y)_0 = 1$, $(x \cdot y)_i = y_i + s^e y_{i-1} = a_1 - y_{i-1} \cdot s^e + y_{i-1} \cdot s^e = a_1$ for $i = 1, \dots, k-1$ and $(x \cdot y)_k = s^e y_{k-1} = a_k$. Consequently, $a = x \cdot y$. Analogously for $a_0 = -1$.

1.2.2. Lemma. Let $a \in M_{11}[s]$, $b \in M_i$, $i = 0, 1$.

(i) If $\mathcal{M}_{11} \models b \mid a$ then $\mathcal{M}_1^* \models b \mid a_j$, $j = 0, \dots, v(a)$.

(ii) If $b \mid s$ and $\mathcal{M}_1^* \models b \mid a_0$ then $\mathcal{M}_{11} \models b \mid a$.

Proof. (i) If $a = b \cdot c$ and $c \in M_{11}[s]$, then $a_i = b \cdot c_i$, $i = 0, 1, \dots, v(a)$.

(ii) We have $\frac{b}{s} \in A_0$, and hence $\frac{a_i}{s} \in M_1^*$, $i = 1, \dots, v(a)$. Since $\frac{a_0}{s} \in M_1^*$, the statement follows.

§ 2. The consistency of Ar with $\neg (\exists x) \text{Prm}(x)$ and with $(\exists x) \text{Prm}(x)$ & $\neg (\exists x) \text{Prm}_2(x)$

The models in question are $\mathcal{M}_{10}(s)$ with $\mathcal{M}_1 = \mathcal{U}_1$.

2.0.0. Theorem. $\text{Ar} \cup \{ \neg (\exists x) \text{Prm}(x) \}$ is consistent.

Proof. Let $L \in A_0 - M_0$, $s = Ll$. We prove that $\mathcal{M}_{10} = \mathcal{M}_{10}(s)$ (with $\mathcal{M}_1 = \mathcal{U}_1$) is the required model. First, $s \in A_0$ and for every standard n we have $n \mid s$. Thus,

$\mathcal{M}_{10}(s) \models \text{Ar}$ follows by 1.1.2.

Let $a \in M_{10}[s]$, $v(a) \geq 2$. If $a_0 = \pm 1$, then

$\mathcal{M}_{10} \models \neg \text{Prm}(a)$ follows from 1.2.1. If $a_0 = 0$ then evidently $\mathcal{M}_{10} \models \neg \text{Prm}(a)$. If $a_0 \notin \{0, +1, -1\}$, then $|a_0| \in M_0$ and $|a_0| \mid \mathcal{M}_{10} a$ (this follows from $|a_0| \mid s$ and (ii) of 1.2.2). Consequently, $a \in M_{10}[s]$ and $v(a) \geq 2$ implies

$\mathcal{M}_{10} \models a \mid x \rightarrow \neg \text{Prm}(x)$.

Now, we will prove the consistency of Ar with

$$(\exists x)\text{Prm}(x) \ \& \ \neg (\exists x)\text{Prm}_2(x).$$

2.1.0. As it is well known,

- (i) $P \vdash \text{Prm}(p) \ \& \ p \mid x.y \rightarrow p \mid x \vee p \mid y$,
- (ii) $P \vdash \text{Prm}(p) \ \& \ p \nmid z \ \& \ z \mid p^x.y \rightarrow z \mid y$.

2.1.1. Let $p \in M_0 - N$ be prime, $L \in A_0 - M_0$ and

$$s = r(L, p).$$

(For the definition of r see 0.1.1.)

Lemma. If $d \in M_0$ and $d > 1$, then $r(d, p) \mid s$.

Proof. We first prove that $c \in M_0$ and $p \nmid c$ implies $c \mid s$. This follows from (ii) of 2.1.0 using $c \mid Ll$ and $Ll = p^{e(L, p)} \cdot s$.

We have $r(d, p) < d$, hence $r(d, p) \in M_0$ and $p \nmid r(d, p)$.

Consequently, $r(d, p) \mid s$.

As a consequence we obtain immediately!

Corollary. For every standard n , $n \mid s$.

2.1.2. Let $\mathcal{M}_1 = \mathcal{U}_1$.

$\mathcal{M}_{10}(s) \models \text{Ar}$ follows from 1.1.2 by Corollary from 2.1.1.

Theorem. (1) $\mathcal{M}_{10}(s) \models (\exists x)\text{Prm}(x)$,

(2) $\mathcal{M}_{10}(s) \models \neg (\exists x)\text{Prm}_2(x)$.

Proof. (1) (a) Let $a = \alpha^k a_k + a_0 \in M_{10}[s]$, $a_k \in M_1$, $a_0 \in M_0$, $\text{Prm}(a_0)$ and $a_0 \nmid a_k$. We prove that a is not decomposable in $M_{10}[s]$. If $a = x.y$ and $x, y \in M_{10}[s]$, then $k \geq 2$, $v(x) + v(y) = k$ and $x_0.y_0 = a_0$. Let $|x_0| = 1$, $|y_0| = a_0$. If $j < v(y)$ and $a_0 \nmid y_j$, $i = 0, \dots, j$, then $a_0 \mid y_{j+1}$ follows

from $0 = a_{j+1} = \sum_{m+n=j+1} x_m \cdot y_n$. Thus $a_0 \mid a_k$ follows from $a_k = x_{v(x)} \cdot y_{v(y)}$, which is a contradiction.

(b) If $e \in A_0 - N$, then we have $\text{Prm}^{\mathcal{M}_{10}}(\alpha^k s^e + p)$.

Proof. $\alpha^k s^e + p$ is not decomposable in $M_{10}[s]$ by

(a). Let $1 < b$, $b \in M_0$ and $b \mid \alpha^k s^e + p$. Thus $b \mid s^e$ and $b \mid p$ and, consequently, $b = p$. Finally, $p \mid s$ follows from $p \mid s^e$, which is a contradiction.

Clearly, $a \in M_{10}[s]$ implies $\alpha^{v(a)+1} s^e + p > a$, which finished the proof of (1).

We will prove (2). Let $a \in M_{10}[s]$, $v(a) \geq 2$.

(a) If $a_0 = 0$, then $\neg \text{Prm}^{\mathcal{M}_{10}}(a)$ follows from $s^e \mid \mathcal{M}_{10} a$ for some $e \in A_0 - N$.

(b) If $|a_0| = 1$, then $\neg \text{Prm}^{\mathcal{M}_{10}}(a)$ follows by 1.2.1.

(c) If $|a_0| > 1$, and $r(|a_0|, p) \neq 1$, then $\neg \text{Prm}^{\mathcal{M}_{10}}(a)$.

Proof. $r(|a_0|, p) \mid s$ follows from $r(|a_0|, p) \in M_0$ by using lemma in 2.1.1. Thus $r(|a_0|, p) \mid \mathcal{M}_{10} a$ follows from (ii) of 1.2.2.

(d) Let $|a_0| > 1$, $r(|a_0|, p) = 1$. Let t be such that $|a_0| = p^t$.

(d1) If $a_0 > 1$, then $r(|a_0|, p) \neq 1$ and $\neg \text{Prm}^{\mathcal{M}_{10}}(a + 2)$ follows from (c).

(d2) If $a_0 = -2$, then $(a + 2)_0 = 0$ and $\neg \text{Prm}^{\mathcal{M}_{10}}(a + 2)$ follows from (a).

(d3) If $a_0 = -3$, then $|(a + 2)_0| = 1$ and $\neg \text{Prm}^{\mathcal{M}_{10}}(a + 2)$ follows from (b).

(d4) If $a_0 < -3$, then $|(a + 2)_0| > 1$. Let $r(|a_0 + 2|, p) = 1$. Then there exists a \tilde{t} with $|a_0 + 2| = p^{\tilde{t}}$. Thus $|a_0| - |a_0 + 2| = 2 = p^{\tilde{t}} \cdot (p^{t-\tilde{t}} - 1)$, which is a contradiction.

Thus $r(a_0 + 2, p) \neq 1$ and $\neg \text{Prm}_{\mathcal{M}_{10}}(a + 2)$ follows from (c).

Consequently, $\neg \text{Prm}_2^{\mathcal{M}_{10}}(a)$ follows from (a), (b), (c), (d).

Let $a \in M_{10} \setminus s$, $v(a) \geq 2$. Since $\mathcal{M}_{10} \models a < x \rightarrow \neg \text{Prm}_2(x)$, the proof is completed.

§ 3. The consistency of Ar with $(\exists x)\text{Prm}_2(x)$

3.0.0. At first we are going to construct a model

\mathcal{M}_1 . Let $\beta \in A_1 - A_0$ be prime, $L \in A_0 - N$ and $s = L$. Put $M' = \{ \beta \cdot a_1 + a_0; a_1 > 0, a_1 \in A_1, a_0 \in A_0^* \}$ and there is an $e \in A_1 - N$ with $s^e \mid a_1$,

and

$$M_1 = M' \cup A_0.$$

Lemma. If $a \in M'$, then there is exactly one $a_1 \in A_1$ and $a_0 \in A_0^*$ such that $a = \beta \cdot a_1 + a_0$ and $a_1 > 0$.

Proof is obvious.

Notation. For $a \in M'$, we denote a_0, a_1 the elements of A_1^* such that $a_1 > 0, a_0 \in A_0^*$ and $a = \beta \cdot a_1 + a_0$.

Lemma. M_1 is the universe of a substructure of \mathcal{U}_1 .

3.0.1. Put $\mathcal{M}_1 = \mathcal{U}_1 / M_1$.

Lemma. (0) $\mathcal{U}_0 \subseteq \mathcal{M}_1 \subseteq \mathcal{U}_1$,

(1) $\mathcal{M}_1 \models \text{Ar}$,

(2) there is a $c \in M'$ such that $\mathcal{M}_1 \models \text{Prm}_2(c)$.

Proof: (0) obvious. (1) can be proved similarly as Theorem 1.1.2. (2): First, we shall prove the following statements:

(a) $a \in M'$ and $n \in N$ imply $n \mid a_1$ and $\frac{a}{n} \notin N$. (Obvious.)

(b) If $a \in M'$, $b \in A_0$, then $b \mid a_1$ and $b \mid a_0$ follows from $b \mid a$.

(c) If $a, b \in M'$, $a \cdot b = \beta^2 \cdot u + v$ and $v \in A_0^*$, $a_1, b_1 \in A_0$, then $a_1 b_0 + b_1 a_0 = 0$. (Indeed, we have $\beta \cdot a_1 b_1 + a_1 b_0 + b_1 a_0 = \beta \cdot u$. Thus $\beta \mid a_1 b_0 + b_1 a_0$ and $a_1 b_0 + b_1 a_0 = 0$ follows from $a_1 \cdot |b_0| + b_1 \cdot |a_0| < \beta$.)

(d) If $a = \beta^2 \cdot u + v$, $a \in M'$, $u, v > 0$ and $u, v \in A_0$, then a is not decomposable in M' . (Let $x, y \in M'$ and $x \cdot y = a$. Hence $v = x_0 y_0$ and, consequently $\text{sign}(x_0) = \text{sign}(y_0)$.)

If $x_1, y_1 \in A_0$, then $x_1 y_0 + y_1 x_0 = 0$ follows from (c). Thus $x_1, y_1 \in A_0$ implies $\text{sign}(x_0) \neq \text{sign}(y_0)$, a contradiction.

We have $\beta \cdot u = \beta \cdot x_1 y_1 + x_1 y_0 + y_1 x_0$. If $x_1 \notin A_0$ and $\text{sign}(x_0) = 1$, then, obviously, $u \notin A_0$, a contradiction. We shall prove that $u \notin A_0$ follows from $x_1 \notin A_0$ and $\text{sign}(x_0) = -1$. We have $x_1 \cdot |y_0| < x_1 \cdot \beta$, $y_1 \cdot |x_0| < y_1 \cdot \beta$. Thus $\beta \cdot (x_1 + y_1) > x_1 \cdot |y_0| + y_1 \cdot |x_0|$, and consequently

$$u > x_1 y_1 - (x_1 + y_1) = (x_1 \cdot \frac{y_1}{2} - x_1) + (y_1 \cdot \frac{x_1}{2} - y_1) > x_1 + y_1 \notin A_0. \quad (2 \mid y_1, 2 \mid x_1 \text{ and } \frac{x_1}{2} > 2, \frac{y_1}{2} > 2 \text{ follows}$$

from (a).) The statement (d) is proved.

Let $e \in A_0 - N$, $u = \beta^2 s^e + s^e - 1$. We prove $\text{Prm}_2 \text{Prm}_1(u)$. Note that u is not decomposable in M (this follows from (d) and $s^e \in A_0$). If $a > 1$, $a \in A_0$ and

$\text{Prm}_1 \mid a \mid u$, then $a \mid \beta \cdot s^e$ and $a \mid s^e - 1$. β is prime, thus $a \mid s^e$ follows by using (ii) of 2.1.0, a contradiction. We have $\text{Prm}_2 \text{Prm}_1(u)$. Case $u + 2$ can be proved like the case u . Clearly, $u \in A_0$ and u is the required element c .

3.1.0. Let \mathcal{M}_1 , s be as in 3.0.0. We have
 $\mathcal{M}_1(s) \models \text{Ar}$.

Theorem. $\mathcal{M}_1(s) \models (\exists x) \text{Prm}_2(x)$.

Proof. (a) Let $a \in M_1[s]$, $v(a) = k$, $a_{k-1} = a_{k-2} = \dots = a_1 = 0$, $\text{Prm}_{\mathcal{M}_1}(a_0)$ and $a_0 \not\vdash_{\mathcal{M}_1} a_k$. Then
 $\text{Prm}_{\mathcal{M}_1}(a)$.

We shall first prove that a is not decomposable in $M_1[s]$.

Contrarywise, assume that $a = x \cdot y$ and $x, y \in M_1[s]$. Then $x_0 \cdot y_0 = a_0$ and $v(x) + v(y) = k$. Let $|x_0| = 1$, $|y_0| = a_0$. Thus $a_0 \mid_{\mathcal{M}_1} y_0$. Let $j < v(y)$ and $a_0 \mid_{\mathcal{M}_1} y_i$, $i = 0, 1, \dots, \dots, j$. $|y_{j+1}| = |\sum_{m+n=j} x_{m+1} y_n|$ follows from $0 = \sum_{m+n=j+1} x_m y_n$, and consequently $a_0 \mid_{\mathcal{M}_1} y_{j+1}$. Thus $a_0 \mid_{\mathcal{M}_1} y_i$, $i = 0, \dots, v(y)$. We have $a_k = x_{v(x)} \cdot y_{v(y)}$. Consequently, $a_0 \mid_{\mathcal{M}_1} a_k$, a contradiction.

Let $b \in M_1$, $b > 1$ and $b \mid_{\mathcal{M}_1} a$. Then $b \mid_{\mathcal{M}_1} a_k$ and $b \mid_{\mathcal{M}_1} a_0$. Thus $b = a_0$, a contradiction.

(b) Let $e \in A_0 - N$, $p \in M_1 - A_0$ with $\text{Prm}_2^{\mathcal{M}_1}(p)$ (by using (2) of 3.0.1). $p \not\vdash_{\mathcal{M}_1} s^e$ and $p + 2 \not\vdash_{\mathcal{M}_1} s^e$ follows from $s^e \in A_0$. Let $c(k) = \alpha^k s^e + p$, $k \in N$ and $k \geq 1$. $\text{Prm}_2^{\mathcal{M}_1}(c(k))$ follows from a. Clearly, if $a \in M_1[s]$, then $a < c(v(a) + 1)$, and hence the proof is completed.

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