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## Correspondence between interval $\pi$ -equivalences and Sd-functions

#### JIŘÍ WITZANY

Abstract. In this paper we study interval  $\pi$ -equivalences, that is we want to study Sdfunctions from the class of rational numbers Q to Q by means of these  $\pi$ -equivalences. A theorem is proved which says that to each interval  $\pi$ -equivalence there exists an Sd<sup>\*</sup>function to which the  $\pi$ -equivalence corresponds.

Keywords: Alternative Set Theory, interval  $\pi$ -equivalence, function.

Classification: 03E70, 54C30

#### Introduction.

A classical real function  $\mathcal{F}$  (i.e. a closed figure in  $Q^2$ ) can be represented by an Sd-function  $F: Q \to Q$  such that  $\mathcal{F} = \operatorname{Fig}(F)$ . We want to study  $\mathcal{F}$  by means of that Sd-function F and the Sd-function by means of an interval  $\pi$ -equivalence  $R_F$  on the class of all rational numbers Q which is in a canonical way assigned to F.

Throughout the paper we use usual notations and principles of the Alternative Set Theory (see [V]). In the first section, basic propositions concerning interval  $\pi$ symmetries are proved, discrete basis theorem is also proved. Then the structure of Q and the  $\pi$ -symmetries are studied in a connection with automorphisms. Finally there is proved an important theorem stating that to each interval  $\pi$ -equivalence Rthere exists an  $Sd^*$ -function F such that  $R = R_F$ .

First section, basic notions and motivations of this paper are due to P.Vopěnka. I also thank K.Čuda for many valuable remarks to the studied matter.

#### 1. Interval $\pi$ -symmetries (equivalences).

Let the letters x, y, z (event. with indices) be variables for rational numbers from Q.

Definition. A symmetry R is called to be an interval if

$$(\forall x, y, z)(\langle x, z \rangle \in R \& x \leq y \leq z \rightarrow \langle x, y \rangle \in R \& \langle y, z \rangle \in R).$$

Obviously if  $\mathcal{M} \neq 0$  is a class of interval symmetries (equivalences) then  $\cap \mathcal{M}$  is an interval symmetry (equivalence). If R is a symmetry then we denote

$$\overline{R} = \{ \langle x, y \rangle; (\exists x_1, y_1) (\langle x_1, y_1 \rangle \in R \& x_1 \leq x, y \leq y_1) \}.$$

Obviously  $\overline{R}$  is an interval symmetry. If R is an equivalence then  $\overline{R}$  is an interval equivalence. If R is a  $\pi$ -class then  $\overline{R}$  is also a  $\pi$ -class.

**Definition.** Let R be an interval symmetry. We say that X is an R-cut if

- 1)  $X \subseteq Q \& \emptyset \neq X \neq Q$ ,
- 2)  $(\forall x, y)(x \in X \& y \leq x \rightarrow y \in X),$
- 3) R''X = X.

We say that x is its inner or outer R-head if  $X = \{y; y \le x\} \cup R''\{x\}$  or  $Q - X = \{y; x \le y\} \cup R''\{x\}$  respectively.

**Proposition 1.** Let  $S \subseteq R$  be two interval symmetries. Let X be an R-cut. Then X is an S-cut. If x is moreover an inner (outer) S-head of the cut X, then x is an inner (outer) R-head of the cut X.

**PROOF:**  $X \subseteq S''X \subseteq R''X = X$ . Let  $X = \{y; y \le x\} \cup S''\{x\}$  then  $X \subseteq \{y; y \le x\} \cup R''\{x\} \subseteq R''X = X$ .

We say that a property  $\varphi(n)$  holds for almost all  $n \in FN$  if there exists an  $m \in FN$  such that  $\varphi(n)$  holds for all  $n \ge m$ .

**Proposition 2.** Let  $\{R_n; n \in FN\}$  be a sequence of interval  $\pi$ -symmetric such that  $R_{n+1} \subseteq R_n$  for all n. Put  $R = \cap \{R_n; n \in FN\}$ . Let  $X \subseteq Q$  be such that the classes X, Q - X are revealed. The following holds:

- (a) R''X = X iff  $R''_nX = X$  for almost all  $n \in FN$ .
- (b) X is an R-cut iff X is an R<sub>n</sub>-cut for almost all n. Moreover X has an inner (outer) R-head iff X has an inner (outer) R<sub>n</sub>-head for almost all n.

**PROOF:** (a) The case of  $X = \emptyset$  or X = Q is trivial, hence let  $\emptyset \neq X \neq Q$ . Let  $R''_m X = X$ , then  $X \subseteq R'' X \subseteq R''_m X = X$ . On the other hand let X = R'' X. Let us suppose that  $X \neq R''_n X$  for all  $n \in FN$ , hence  $R_n \cap (X \times (Q - X)) \neq \emptyset$  for all n. Then  $R \cap (X \times (Q - X)) \neq \emptyset$ , thus  $X \neq R'' X$  - a contradiction. By that we have proved that there exists an m such that  $X = R''_m X$ . Let  $n \ge m$  then  $X \subseteq R''_m X \subseteq R''_m X = X$ .

(b) From (a) it follows that X is an R-cut iff there exists an m such that X is an  $R_n$ -cut for all  $n \ge m$ . Let x be an inner (outer) R-head of the cut X. Then (by the proposition 1) x is an inner (outer)  $R_n$ -head of X for almost all n. Let conversely  $x_n$  be an inner  $R_n$ -head of the cut X for all  $n \ge m$ . Let  $x \in X$  be such that  $x_n \le x$  for all  $n \ge m$ . We prove that x is an inner R-head of the cut X. Let  $x \le y$ . If  $y \in X$  then  $\langle x_n, y \rangle \in R_n$  for all  $n \ge m$ , thus  $y \in R''\{x\}$ . Hence  $X \subseteq \{y; y \le x\} \cup R''\{x\} \subseteq R''X = X$ , which means that x is an inner R-head of X. The case of the outer head is similar.

**Proposition 3.** Let R be an interval  $\pi$ -symmetry. Let X be an Sd-class such that  $X \subseteq Q, \emptyset \neq X \neq Q, R''X = X$ . Then there exists a set-theoretically definable R-cut Y.

**PROOF:** Obviously Q - X is an Sd-class and R''(Q - X) = Q - X. Let us assume that Q - X is not an R-cut. Then there exist  $x_0 \in X, y_0 \in (Q - X), x_0 < y_0$ , thus  $x_0 \in X, y_0 \notin X$ . Put  $Y = \{x; (\exists y \in X) (x \le y < y_0)\}$ . Obviously Y is an Sd-class which satisfies the first two conditions from the definition of R-cut. Let us prove that it satisfies the third condition. By contradiction let us assume that there exist

 $x_1 \in Y, z \notin Y$  such that  $\langle x_1, z \rangle \in R$ . Let  $x_2 \in X$  be from the definition of Y such that  $x_1 \leq x_2 < y_0$ . Obviously  $x_2 < z$  because otherwise it would be  $z \in Y$ . It implies  $\langle x_2, z \rangle \in R$ , thus  $z \in R''X = X$ . If  $y_0 \leq z$  then we would have  $\langle x_2, y_0 \rangle \in R$ , thus  $y_0 \notin R$ . Consequently  $z < y_0$ . Since  $z \in X, z \in Y$ , and this is the desired contradiction.

**Proposition 4.** Let R be an interval  $\pi$ -symmetry. Then there exists its generating sequence  $\{R_n; n \in FN\}$  such that  $R_n$  is an interval Sd-symmetry for all n. Moreover if R is an equivalence then  $R_{n+1} \circ R_{n+1} \subseteq R_n$  can be assumed for all n.

PROOF: Let  $\{S_n; n \in FN\}$  be a generating sequence of the  $\pi$ -symmetry R. Obviously  $\overline{S}_n$  is an interval Sd-symmetry,  $\overline{S}_{n+1} \subseteq \overline{S}_n$ ,  $S_n \subseteq \overline{S}_n$  for all  $n \in FN$ . From this  $R = \cap\{S_n; n \in FN\} \subseteq \cap\{\overline{S}_n; n \in FN\}$ . Let  $\langle x, y \rangle \in \cap\{\overline{S}_n; n \in FN\}$ . We want to prove  $\langle x, y \rangle \in R$ . There exists a sequence  $\{\langle x_n, y_n \rangle; n \in FN\}$  such that  $\langle x_n, y_n \rangle \in S_n, x_n \leq x, y \leq y_n$  for all n. Let  $\{\langle x_\alpha, y_\alpha \rangle; \alpha \in \gamma\}$  be a prolongation of this sequence such that  $x_\alpha \leq x, y \leq y_\alpha$  for  $\alpha \in \gamma$  and  $\langle x_\alpha y_\alpha \rangle \in S_n$  for  $n \in FN$ ,  $\alpha \geq n$ . Hence  $\langle x_\alpha, y_\alpha \rangle \in R$  if  $\alpha \in \gamma - FN$  and since R is an interval symmetry, we see  $\langle x, y \rangle \in R$ . We have proved that  $\{\overline{S}_n; n \in FN\}$  is a generating system of the  $\pi$ -symmetry R with the desired properties. If R is moreover an equivalence then by the theorem III.1.1[V] it is possible to select from this sequence a generating subsequence  $\{R_n; n \in FN\}$  such that  $R_{n+1} \circ R_{n+1} \subseteq R_n$  for all n.

**Definition.** We say that a class  $D \subseteq Q$  is a discrete basis of a  $\pi$ -symmetry R if

1)  $(\forall x)(\exists y \in D)(\langle x, y \rangle \in R),$ 

2) 
$$(\forall \gamma \in N)$$
 Set  $\{x; x \in D \& -\gamma \leq x \leq \gamma\}$ .

We say that  $x, y \in D$  are neighbouring if  $x \neq y$  and

$$(\forall z)(\min\{x,y\} < z < \max\{x,y\} \to z \notin D).$$

**Theorem.** Let R be an interval  $\pi$ -symmetry. Then the following conditions are equivalent:

- (a) There exists a discrete basis of the  $\pi$ -symmetry R.
- (b) Each set-theoretically definable R-cut has an inner and an outer R-head.

**PROOF:** (a) $\rightarrow$ (b). Let D be a discrete basis of the  $\pi$ -symmetry R. Let X be a set-theoretically definable R-cut. Let  $x_0 \in X, y_0 \notin X$  and  $x_1, y_1 \in D$  be such that  $\langle x_0, x_1 \rangle \in R, \langle y_0, y_1 \rangle \in R$ . Obviously  $x_1 \in X, y_1 \notin X$ . Let  $\gamma \in N$  be such that  $-\gamma \leq x_1 < y_1 \leq \gamma$ . Put  $u = X \cap \{x; x \in D \ \& \ -\gamma \leq x \leq \gamma\}$ . We see  $x_1 \in u, y_1 \notin u$ . Let  $x_2$  be the greatest element in the set u in the natural ordering of  $Q, y_2 \in D, y_2 > x_2$  its neighbour in D. Obviously  $y_2 \notin X$ . If  $x_2 \leq z \leq y_2$  then either  $\langle x_2, z \rangle \in R$  or  $\langle z, y_2 \rangle \in R$  and these two cases exclude one another because  $\langle x_2, z \rangle \in R$  implies  $z \in R''X = X$  and  $\langle z, y_2 \rangle \in R$  implies  $z \in R''(Q - X) = Q - X$ . From this it follows  $X = \{y; y \leq x_2\} \cup R''\{x_2\}, Q - X = \{y; y \geq y_2\} \cup R''\{y_2\}$ . Consequently  $x_2$  is an inner and  $y_2$  and  $y_2$  an outer R-head of the R-cut X.

(b) $\rightarrow$ (a). Let  $\{R_n; n \in FN\}$  be a generating sequence of the  $\pi$ -symmetry R such that  $R_n$  is an interval Sd-symmetry for all n (see Proposition 4). There exist (by

the theorem III.1.3[V]) Sd-classes  $D_n$  such that  $D_n$  is a maximal  $R_n$ -net. If  $m \leq n$  then  $R_n \subseteq R_m$  and hence

(1')  $(\forall n)(m \leq n \rightarrow (\forall x)(\exists y \in D_n)(\langle x, y \rangle \in R_m)).$ 

We will prove that the following holds:

(2')  $(\forall n)(\forall \gamma \in N)$  Set  $\{x, x \in D_n \& -\gamma \leq x \leq \gamma\}$ .

Choose an  $n \in FN$  and let  $\gamma \in N$  be such that (2') does not hold. Put  $Y = \{y; \operatorname{Set}\{x; x \in D_n \& -\gamma \leq x \leq \gamma \& x \leq y\}\}$ . Obviously Y is an Sdclass satisfying the first two conditions from the definition of the R-cut. We prove  $Y = R_n''Y$ . Let  $x_0 \in Y, y_0 \notin Y, \langle x_0, y_0 \rangle \in R_n$ , obviously  $x_0 < y_0$ . If  $z_1, z_2 \in D_n$ would be such that  $x_0 \leq z_1 < z_2 \leq y_0$  then  $\langle z_1, z_2 \rangle \in R_n$ , which is impossible because  $D_n$  is an  $R_n$ -net. So between  $x_0, y_0$  there lies at most one element of the class  $D_n$  and so  $y_0 \in Y$  - a contradiction. By that we have proved that Y is an  $R_n$ -cut and thus also an R-cut. Let  $x_1$  be its inner and  $y_1$  outer Rhead and thus also  $R_n$ -head (see Proposition 1). Let  $z_1, z_2, z_3 \in D_n$  be such that  $x_1 \leq z_1 < z_2 < z_3 \leq y_1$ . Then either  $\langle x, z_2 \rangle \in R_n$  and so  $\langle z_1, z_2 \rangle \in R_n$  or  $\langle z_2, y_1 \rangle \in R_n$  and so  $\langle z_2, z_3 \rangle \in R_N$ . Between  $x_1, y_1$  thus there can lie at most two elements of the class  $D_n$ . This implies  $y_1 \in Y$  - a contradiction.

Let  $\{D_{\alpha}; \alpha \in \delta\}$  be an  $Sd^*$ -prolongation of the sequence  $\{D_n; n \in FN\}$  such that for all  $\alpha \in \delta$  the following holds:

$$(\forall \gamma \in N) \operatorname{Set} \{x; x \in D_{\alpha} \& -\gamma \leq x \leq \gamma\}.$$

Let  $\delta_m \in (\delta - FN)$  be such for all  $\alpha \in N, m \leq \alpha \leq \delta_m$ , it holds

$$(\forall x)(\exists y \in D_{\alpha})(\langle x, y \rangle \in R_m).$$

Choose  $\alpha \in \delta - FN$  so that  $\alpha \leq \delta_m$  for all m. Then

 $(\forall m)(\forall x)(\exists y \in D_{\alpha})(\langle x, y \rangle \in R_m),$ 

thus  $R''_m\{x\} \cap D_\alpha \neq \emptyset$  for every m, x. Since  $D_\alpha$  is an  $Sd^*$ -class,  $D_\alpha$  is revealed and so  $R''\{x\} \cap D_\alpha = \cap\{R''_m\{x\}; m \in FN\} \neq \emptyset$ . But it means that

$$(\forall x)(\exists y \in D_{\alpha})(\langle x, y \rangle \in R).$$

We have proved that  $D_{\alpha}$  is a discrete basis of the  $\pi$ -symmetry R.

**Proposition 5.** Let R be an interval  $\pi$ -symmetry which has a discrete basis D. Let X be an R-cut, X a sharp class, i.e.  $(\forall u) \operatorname{Set}(X \cap u)$ . Then X is an Sd-class.

**PROOF:** Let  $\gamma \in N$  be such that  $-\gamma \in X, \gamma \notin X$ . Put  $d = \{y; -\gamma \le y \le \gamma \& y \in D\}$ . Then  $X = \{y; y \le \gamma\} \cup R''(d \cap X), Q - X = \{y; \gamma \le y\} \cup R''(d - X)$ , thus X and Q - X are  $\pi$ -classes and so Sd-classes.

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**Proposition 6.** Let R be an interval compact  $\pi$ -symmetry. Then R has a discrete basis which is a set and the class  $\mathcal{M}$  of all set-theoretically definable R-cuts is at most countable.

PROOF: Let  $\{R_n; n \in FN\}$  be a generating sequence of the  $\pi$ -symmetry R. Let d be a set such that  $(\forall x)(\exists y \in d)(\langle x, y \rangle \in R)$  (see Theorem III.1.5[V]). Obviously d is a discrete basis of the  $\pi$ -symmetry R. If  $x \in d$  then  $\overline{x}$  denote its right neighbour in the set d. For  $X \in \mathcal{M}$  let  $c_X$  denote the greatest element of the set  $X \cap d$ . Obviously  $\langle c_X, \overline{c}_X \rangle \notin R$  and  $c_X \neq c_Y$  for  $X, Y \in \mathcal{M}, X \neq Y$ . Put  $A_n = \{c_X; X \in \mathcal{M} \& \langle c_X, \overline{c}_X \rangle \notin R_n\}$ . Obviously

$$\{c_X; X \in \mathcal{M}\} = \cup \{A_n; n \in FN\}.$$

Thus it suffices to prove that each  $A_n$  is a finite class. If  $x, y \in A_n, x < y$  then  $\overline{x} \leq y$  from the definition of  $\overline{x}$ , hence  $\langle x, y \rangle \notin R_n$ . But this means that  $A_n$  is an  $R_n$ -net and so by the theorem III.1.3[V]  $A_n$  is finite.

More generally as a consequence of some deeper results of  $[\check{\mathbf{C}}]$  it can be proved that the class of all clopen figures in a compact  $\pi$ -symmetry is countable.

#### 2. Interval $\pi$ -symmetries and automorphisms.

Let  $\stackrel{\circ}{=}$  mean the basic equivalence on the universe V (see [ $\check{\mathbf{C}}\mathbf{K}$ ] or the section V.1[ $\mathbf{V}$ ]).

**Proposition 1.** Let  $X \subseteq Q$  be an Sd-cut. Then Fig.(X) is also a cut.

**PROOF:** If  $F: V \to V$  is an automorphism then obviously F''X is also an Sd-cut. Since  $\operatorname{Fig}_{\bullet}(X) = \bigcup \{F''X; F \text{ is an automorphism}\}, \operatorname{Fig}_{\bullet}(X)$  is a cut.

**Proposition 2.** Let F be an automorphism,  $X \subseteq Q$  an Sd-cut which is not Sdg. Then  $F''X \neq \operatorname{Fig}_{\bullet}(X)$ .

**PROOF:** Let us suppose that  $F''X = \operatorname{Fig}_{\underline{\bullet}}(X)$ . It implies that  $\operatorname{Fig}_{\underline{\bullet}}(X)$  is an Sd-class. It is also a  $\stackrel{\circ}{=}$ -figure, it is proved in the section V.1[V] that then it is an  $Sd_{\emptyset}$ -class. Consequently X is an  $Sd_{\emptyset}$ -class – a contradiction.

Sd-cuts represent classical real numbers in the sense of Dedenkind's cuts. The following proposition says that these Sd-cuts are being moved by automorphisms in the limits given by Sdg-cuts which are firm with respect to the automorphisms.

Define an interval  $\pi$ -equivalence

$$R_0 = \cap \{Z^2 \cup (Q-Z)^2; Z \text{ is an } Sd_{\emptyset} \text{-cut}\}$$

Proposition 3. Let X be a cut, then  $\operatorname{Fig}_{\bullet}(X) = \operatorname{Fig}_{R_0}(X)$ .

**PROOF:** Obviously  $\operatorname{Fig}_{\underline{\bullet}}(X) \subseteq \operatorname{Fig}_{R_0}(X)$ . From the definition of  $\overset{\circ}{=}$ 

 $\operatorname{Fig}_{\bullet}(X) = \cap \{A; A \text{ is } Sd_{\bullet} \& X \subseteq A\}.$ 

For an  $Sd_{\emptyset}$ -class  $A \supseteq X$  put  $Z_A = \{y \in A; (\forall z)(z \le y \to z \in A)\}$ , it is an  $Sd_{\emptyset}$ -cut. Since  $\operatorname{Fig}_{\bullet}(X) = \cap\{Z_A; A \text{ is } Sd_{\emptyset} \& X \subseteq A\}$ ,  $\operatorname{Fig}_{R_0}(X) \subseteq \operatorname{Fig}_{\bullet}(X)$ .  $Sd_{\emptyset}$ -cuts occupy a special place among all rational cuts. Thus let us define a class of concrete real numbers:

 $CR = \{X; X \text{ is an } Sd_{\emptyset} \text{-cut } \& X \text{ has not a last element} \}.$ 

All finite rational and algebraic numbers,  $\pi$ , e etc. belong in the classical sense to CR. This class is countable and is closed under algebraic operations and under the operation of supremum over  $Sd_{\theta}$ -subclasses. The nonexistence of an infinitesimally small concrete real number is equivalent to the axiom of elementary equivalence.

Now let R be an interval  $\pi$ -symmetry and X an Sd-R-cut. We say that X is limit if X has not its inner or outer head. We will give a sufficient condition on R to have a limit Sd-cut.

In the rest of this section we suppose that the axiom of elementary equivalence holds (i.e. Def = FV).

**Proposition 4.** Let S be an interval  $Sd_{\emptyset}$ -symmetry. If there exists an Sd-cut X of S such that  $X \cap BQ \neq \emptyset$ ,  $BQ - X \neq \emptyset$  and  $X \notin Sd_{\emptyset}$  then S has a limit Sd-cut.

**PROOF:** Let X be an Sd-cut,  $X \notin Sd_{\emptyset}, S''X = X$ . Let us suppose that  $0 \in X$ ,  $1 \notin X$ . Let A be a maximal  $Sd_{\emptyset}$ -S-net on [0,1]. If A would be a finite class then X could not be limit. If Set(A) then  $card(A) \in Def$  but  $card(A) \notin FN$ . Thus A is a proper uncountable Sd-class. From the theorem of the preceding section it follows that there has to exist an Sd-cut Y of S which is limit.

**Corollary.** Let R be an interval  $\pi_{\emptyset}$ -equivalence. If R has an Sd-cut X such that  $X \cap BQ \neq \emptyset$ ,  $BQ - X \neq \emptyset$  and  $X \notin Sd_{\emptyset}$  then R has a limit Sd-cut.

PROOF: See Proposition 1.2.

The converse implication does not hold – the interval  $\pi_{0}$ -equivalence

$$R_{+} = \{ \langle x, y \rangle; x = y = 0 \text{ or } x \neq 0 \& y \neq 0 \& (\forall n) (|\frac{x}{y} - 1| \le \frac{1}{n}) \}$$

has just two Sd-cuts  $\{x; x < 0\}$  and  $\{x; x \le 0\}$  which are both Sd<sub>0</sub> and limit.

3. Correspondence between interval  $\pi$ -equivalences and rational Sd-functions.

**Definition.** Let  $F: Q \to Q$  be a function. We define the relation

$$R_F = \{\langle x, y \rangle; (\forall z)(\min\{x, y\} \le z \le \max\{x, y\} \to F(x) \doteq F(z) \doteq F(y))\},\$$

where  $\doteq$  is the standard compact indiscernible equivalence on Q.

**Proposition 1.** If  $F : Q \to Q$  is an Sd-function then  $R_F$  is an interval  $\pi$ -equivalence.

**PROOF:** Obviously  $R_F$  is an interval equivalence. Let  $\{S_n; n \in FN\}$  be a generating system of  $\doteq$ . Put

$$R_{F,n} = \{ \langle x, y \rangle; (\forall z)(\min\{x, y\} \le z \le \max\{x, y\} \rightarrow \langle F(x), F(z) \rangle \in S_n \& \langle F(z), F(y) \rangle \in S_n \} \},$$

then  $R_{F,n}$  is an Sd-class and  $R_F = \cap \{R_{F,n}; n \in FN\}$ .

**Examples.** (i) Let  $x \in Q$ , then  $\alpha \le x < \alpha + 1$  for an  $\alpha \in N \subseteq Q$ . Define  $F(x) = (x - \alpha)(-1)^{\alpha} + (\alpha + 1 - x)(-1)^{\alpha+1}$  (see fig. 1). Then  $R_F = \{\langle x, y \rangle; |x - y| \doteq Q\}$  is a noncompact  $\pi$ -equivalence on Q.



(ii) Let  $x \in Q$ , then  $x = \alpha/\beta$  where  $\alpha, \beta \in N$  are relatively prime. Put  $F(x) = (-1)^{\alpha}$ . Then  $R_F = \{\langle x, y \rangle; x = y\}$  is a discrete equivalence on Q.

We want to investigate Sd-functions from Q to Q through interval  $\pi$ -equivalences  $R_F$ . Results on interval  $\pi$ -equivalences can be applied on Sd-functions.

In this section we will show that to each interval  $\pi$ -equivalence R there exists an  $Sd^*$ -function F such that  $R = R_F$ .

**Definition.** Let R be a symmetry on Q, we define the relation of connectedness of R as usually

$$Cntd_R(u) \equiv (\forall v \subseteq u) (\emptyset \neq v \neq u \to (\exists z_1 \in v) (\exists z_2 \in u - v) (\langle z_1, z_2 \rangle \in R)),$$
  
$$S = \{\langle x, y \rangle; (\exists u) (x, y \in u \& Cntd_R(u))\}.$$

**Proposition 2.** Let R be an interval symmetry, S the relation of connectedness of R. Then

- (a)  $R \subseteq S$  and S is an interval equivalence.
- (b) An Sd-class X is R-cut iff it is S-cut.
- (c) If R is an interval  $\pi$ -symmetry then S is an interval  $\pi$ -equivalence.

**PROOF:** (a) It is obvious from the definition that  $R \subseteq S$ . Since  $u_1 \cap u_2 \neq \emptyset$  &  $Cntd_R(u_1)$  &  $Cntd_R(u_2)$  implies  $Cntd_R(u_1 \cup u_2)$ , we see that S is an equivalence. Finally let x < z < y and  $\langle x, y \rangle \in S$ , then there is a  $u \subseteq Q$  such that  $Cntd_R(u)$  and  $x, y \in u$ . Put  $v = \{x_1 \in u; x_1 \leq z\}$ ,  $z_1 = \max(v), z_2 = \min(u - v)$ . Then necessarily  $\langle z_1, z_2 \rangle \in R$  and so  $\langle z_1, z_2 \rangle \in R$ ,  $\langle z, z_2 \rangle \in R$ . Consequently  $Cntd_R(v \cup \{z\})$ ,  $Cntd_R((u - v) \cup \{z\})$  and  $\langle x, z \rangle \in S$ ,  $\langle z, y \rangle \in S$ . We have proved that S is an interval equivalence.

(b) If an Sd-class X is an S-cut, it is also an R-cut because  $R \subseteq S$ . Let an Sd-class X be an R-cut and  $x \in X, y \notin X$  be such that  $\langle x, y \rangle \in S$ . It means that there is a  $u \subseteq Q$  such that  $x, y \in u$  and  $Cntd_R(u)$ . Put  $v = u \cap X$ , there has to exist  $z_1 \in v, z_2 \in u - v$  such that  $\langle z_1, z_2 \rangle \in R$ , but  $z_1 \in X, z_2 \notin X$  - a contradiction.

(c) If R is an Sd-class then it is obvious from the definition that S is also Sd. Let R be an interval  $\pi$ -symmetry,  $R = \cap \{R_n; n \in FN\}$  where  $R_n$  are interval

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Sd-symmetries (see Proposition 1.4). Let  $S_n$  be the relations of connectedness of  $R_n$ , hence  $S_n$  are Sd-classes. It holds that (see theorem III.3.1[V])

(1) 
$$Cntd_R(u) \Leftrightarrow (\forall n)(Cntd_{R_n}(u)).$$

It is obvious  $S \subseteq \cap \{S_n; n \in FN\}$ . Let  $\langle x, y \rangle \in \cap \{S_n; n \in FN\}$ . Then there are  $u_n \subseteq Q$  such that  $x, y \in u_n$  &  $Cntd_{R_n}(u_n)$ . Take a prolongation  $\{u_{\delta}; \delta \in \gamma\}$  of the sequence  $\{u_n; n \in FN\}$  such that  $\delta \in \gamma - FN$  and  $n \in FN$  implies  $x, y \in u_{\delta}$  and  $Cntd_{R_n}(u_{\delta})$ . Take a  $\delta \in \gamma - FN$ , then  $x, y \in u_{\delta}$  and  $Cntd_R(u_{\delta})$  (see (1)), thus  $\langle x, y \rangle \in S$ .

**Proposition 3.** Let S be an interval Sd-equivalence. Then there exists an Sd-function  $G: Q \to \{-1, 1\}$  such that  $S = R_G$ .

**PROOF:** Put  $S_0 = \{\langle x, y \rangle; (\forall z)(\min\{x, y\} \le z \le \max\{x, y\} \to \langle x, z \rangle \in S \text{ or } \langle z, y \rangle \in S \}$ , obviously  $S_0$  is an interval Sd-symmetry. Let  $S_1$  be the equivalence of connectedness of  $S_0$ . Obviously  $S \subseteq S_1$ . Let  $A_1$  be a maximal set-theoretically definable  $S_1$ -net and  $A \supseteq A_1$  a maximal set-theoretically definable S-net. By induction we construct a function G on  $A_1$ . Let  $P: N \to A_1$  be an Sd-numbering of the class  $A_1$  (if  $A_1$  is a set, we consider  $P: \alpha \to A_1$ ).

- I. G(P(0)) = 1.
- II.  $G(P(\alpha)) = -G(x_0)$  where  $x_0 = \max\{x \in P''\alpha; x < P(\alpha)\}$  if  $\{x \in P''\alpha; x < P(\alpha)\} \neq 0$ .  $G(P(\alpha)) = 1$  otherwise.

By this an Sd-function G on  $A_1$  is defined. Now let  $x \in A$ . Then there exists just one  $x_0 \in A_1$  such that  $\langle x_0, x \rangle \in S_1$ . Thus there is a  $u \subseteq Q$  such that  $x, x_0 \in u$ and  $Cntd_{S_0}(u)$ . Put  $Z = \{z \in A; \min\{x_0, x\} \le z \le \max\{x_0, x\}\}$ . Let  $z_1 \in Z$ , then there exists a  $z_2 \in u$  such that  $\langle z_1, z_2 \rangle \in S$ . Indeed, let  $z_1 \notin u$ , otherwise it should hold with  $z_2 = z_1$ . Put  $v = \{z \in u; z < z_1\}$ , there are  $z_2 \in v, z_3 \in (u - v)$  such that  $\langle z_2, z_3 \rangle \in S$ . Since  $z_2 \le z_1 \le z_3$ , also  $\langle z_1, z_2 \rangle \in S$ . Since Z is an S-net and S is an Sd-equivalence, there is a one-one Sd-function from Z into u. Thus Set(Z)and we can put  $\alpha = card(Z)$  and  $G(x) = (-1)^{\alpha-1}G(x_0)$ .

Finally let  $x \in Q$ . Then put  $G(x) = G(x_0)$  where  $x_0 \in A$  is such that  $\langle x_0, x \rangle \in S$ . We have defined an Sd-function  $G: Q \to \{-1, 1\}$ . It remains to prove that  $S = R_G$ . It is obvious that  $S \subseteq R_G$ . Let  $\langle x, y \rangle \notin S, x < y$ , we can suppose  $x, y \in A$ . We shall use the common notation  $[x, y] = \{z, x \le z \le y\}$  and  $(x, y) = \{z; x < z < y\}$ . If there exist  $z_1, z_2 \in A \cap [x, y], z_1 < z_2$  such that  $\langle z_1, z_2 \rangle \in S_1$ , then from the definition of G it is obvious that  $\langle x, y \rangle \notin R_G$ . Let us suppose the contrary. Then  $(A - A_1) \cap (x, y) = 0$  and  $\langle x, y \rangle \notin S_1$ . If  $Set(A \cap [x, y])$  then  $u = A \cap [x, y]$  would be  $S_0$  connected -a contradiction. Thus  $A \cap [x, y]$  and also  $A_1 \cap [x, y]$  is an uncountable proper Sd-class. Let  $x = P(\alpha)$ , necessarily there exist a  $\beta > \alpha$  such that  $P(\beta) \in A_1 \cap [x, y]$ . Let  $\beta_0$  be the first such  $\beta$ . Then  $G(P(\alpha)) \neq G(P(\beta_0))$ , hence again  $\langle x, y \rangle \notin R_G$ .

**Theorem.** Let R be an interval  $\pi$ -equivalence. Then there exists an Sd<sup>\*</sup>-function F such that  $R = R_F$ .

**PROOF:** Let  $R = \cap \{R_n; n \in FN\}$  where  $\{R_n; n \in FN\}$  is a generating system consisting of interval Sd-symmetries such that  $R_{n+1} \circ R_{n+1} \subseteq R_n, \{R_\alpha; \alpha \in \gamma\}$  be its Sd\*-prolongation and  $S_\alpha$  the relations of connectedness of  $R_\alpha$ . We will construct a sequence  $\{\langle F_n, A_n \rangle; n \in FN\}$  of Sd-functions  $F_n : Q \to Q$  and set-theoretically definable maximal  $R_n$ -nets  $A_n$ .

We say that an  $x \in Q$  lies between connected neighbours  $x_1, x_2 \in A_n$  if  $x_1 \leq x \leq x_2, x_1$  and  $x_2$  are neighbours in  $A_n$  and  $(x_1, x_2) \in S_n$ . We say that x lies on the edge  $x_0 \in A_n$  if  $x \geq x_0, x \in R''_n\{x_0\}$  and  $x_0$  is maximal in  $S''_n\{x_0\} \cap A_n$  or  $x \leq x_0, x \in R_n\{x_0\}$  and  $x_0$  is minimal in  $S''_n\{x_0\} \cap A_n$ . We want to satisfy the following conditions (for  $\alpha \in FN$ ):

 $\begin{array}{l} (\mathbf{a}_{\alpha}) \ \text{ If } m < \alpha \ \text{then } A_{\alpha} \supseteq A_m \ \text{and} \ (\forall z \in A_m)(F_{\alpha}(z) = F_m(z)). \\ (\mathbf{b}_{\alpha}) \ \text{ If } x < y \in A_{\alpha} \ \text{then} \end{array}$ 

$$(\exists z_1, z_2 \in [x, y] \cap A_\alpha)(|F_\alpha(z_1) - F_\alpha(z_2)| \ge 1/4^\alpha).$$

If moreover  $x, y \in A_{\alpha}$  are connected neighbours then

$$|F_{\alpha}(x) - F_{\alpha}(y)| \leq 1/2^{\alpha}.$$

 $\begin{array}{l} (\mathbf{c}_{\alpha}) \quad \text{If } m \leq \alpha \text{ and } x \text{ lies between connected neighbours } x_1, x_2 \in A_m, \text{ then} \\ (1) \quad |F_{\alpha}(x) - F_{\alpha}(x_i)| \leq 1/2^m + (1/4^{m+1} + \dots + 1/4^{\alpha}) \leq 1/2^m + 1/(3.4^m) \\ (i = 1, 2). \\ \text{If } x \text{ lies on the edge } x_0 \in A_m, \text{ then} \\ (2) \quad |F_{\alpha}(x) - F_{\alpha}(x_0)| \leq 1/4^{m+1} + \dots + 1/4^{\alpha} \leq 1/(3.4^m). \end{array}$ 

**Lemma.** Let  $\{\langle F_k, A_k \rangle; k \leq n\}$  satisfy the conditions  $(a_k), (b_k)$  and  $(c_k)$   $(k = 0, \ldots, n)$ . Then there exists an Sd-function  $F_{n+1} : Q \to Q$  and a set-theoretically defined maximal  $R_{n+1}$ -net  $A_{n+1}$  satisfying again the conditions  $(a_{n+1}), (b_{n+1})$  and  $(c_{n+1})$ .

**PROOF:** Let  $A_{n+1} \supseteq A_n$  be a maximal set-theoretically defined  $R_{n+1}$ -net. Proposition 3 says that there is an Sd-function G such that  $S_{n+1} = R_G$ . Let us define  $F_{n+1}$  firstly in the points of  $A_{n+1}$ . For  $z \in A_n$  put  $F_{n+1}(z) = F_n(z)$ . For  $z \in A_{n+1} - A_n$  we distinguish two cases:

A. z lies between two connected neighbours  $x, y \in A_n$ .

If  $\langle x, y \rangle \in S_{n+1}$ , then in all points  $z \in A_{n+1} \cap (x, y)$  define  $F_{n+1}(z)$  so that

(a)  $F_{n+1}(z)$  lies between the values  $F_n(x)$  and  $F_n(y)$ ,

(b) if  $z_1, z_2$  are neighbouring in  $A_{n+1}$ , then

$$|F_{n+1}(z_1) - F_{n+1}(z_2)| \in [1/4^{n+1}, 1/2^{n+1}].$$

There is a  $z \in (A_{n+1} \cap (x, y))$  because  $R_{n+1} \circ R_{n+1} \subseteq R_n$ . Let us suppose the contrary, it means  $[x, y] \subseteq R''_{n+1}\{x, y\}$ . Since  $\langle x, y \rangle \in S_{n+1}$  and  $\langle x, y \rangle \notin R_{n+1}$ there are  $z_1 \in R''_{n+1}\{x\}, z_2 \in R''_{n+1}\{y\}, x \leq z_1 \leq z_2 \leq y$  such that  $\langle z_1, z_2 \rangle \in R_{n+1}$ . It implies  $\langle x, y \rangle \in R_n$  – a contradiction. Moreover by the induction hypothesis

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 $|F_{nx} - F_{ny}| \in [1/4^{n}, 1/2^{n}]$ . It is thus possible to satisfy these two conditions (see fig.2).



Fig. 2

Let  $\langle x, y \rangle \notin S_{n+1}$ . Put  $d = F_n(y) - F_n(x)$ . In all points  $z \in A_{n+1} \cap (x, y)$  define  $F_{n+1}(z)$  so that

- (a)  $F_{n+1}(z) \in \{F_n(x), F_n(x) + d/4\}$  if G(x) = G(z),
- (b)  $F_{n+1}(z) \in \{F_n(y) d/4, F_n(y)\}$  if  $G(x) \neq G(z)$  or  $z \in S''_{n+1}\{y\}$ , (c) if  $z_1, z_2 \in A_{n+1} \cap [x, y]$  are neighbouring, then

$$|F_{n+1}(z_1) - F_{n+1}(z_2)| \ge d/4.$$

It is obvious (see fig.3) that these three conditions can be satisfied.



Fig. 3

B. z lies on the edge  $x \in A_n$ . Let  $z \in S_n\{x\} \cap [x, \infty)$  where  $x = \max(S_n''\{x\} \cap A_n)$ , the second case is similar. Define  $F_{n+1}$  in all points  $z \in S''_n\{x\} \cap [x, \infty) \cap A_{n+1}$  so that

(a)  $F_{n+1}(z) \in \{F_n(x) - 1/4^{n+1}, F_n(x)\}$  if G(x) = G(z), (b)  $F_{n+1}(z) \in \{F_n(x), F_n(x) + 1/4^{n+1}\}$  if  $G(x) \neq G(z)$ (c) if  $z_1, z_2$  are neighbouring in  $A_{n+1}$ , then

$$|F_{n+1}(z-1) - F_{n+1}(z-2)| \ge 1/4^{n+1}.$$

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Again it is possible to satisfy these conditions (see fig. 4).



On the rest of Q define  $F_{n+1}$  so that it is linear on [x, y] where  $x, y \in A_{n+1}$  are connected neighbours or constant on  $S''_{n+1}\{x\} \cap [x, \infty)$  or  $S''_{n+1}\{x\} \cap \langle (-\infty, x]$  where x is maximal or minimal in  $S''_{n+1}\{x\} \cap A_{n+1}$  respectively.

The conditions  $(a_{n+1})$  and  $(b_{n+1})$  are obvious from the construction of  $F_{n+1}$ . Let us prove  $(c_{n+1})$ .

Firstly let m = n + 1. If x lies between connected neighbours  $x_1, x_2 \in A_m$ , then  $|F_m(x) - F_m(x_i)| \le 1/2^m$  (i = 1, 2). If x lies on the edge  $x_0 \in A_m$ , then  $|F_m(x) - F_m(x_0)| = 0$ .

Secondly, let m < n + 1, then the induction hypothesis states that  $(c_n)$  with this *m* holds. Let *x* lie between connected neighbours  $z_1, z_2 \in A_n$ . Then  $F_{n+1}(x)$ lies between  $F_n(z_1), F_n(z_2)$  and since in the first case of  $(c_{n+1})z_1, z_2$  lie between  $x_1, x_2 \in A_m$ , in the second case  $z_1, z_2$  lie on the edge  $z_0 \in A_m$ , (1) or (2) of  $(c_{n+1})$ holds. Let *x* lie on the edge  $z_0 \in A_n$ , then  $|F_{n+1}(x) - F_n(z_0)| \le 1/4^{n+1}$  as it follows from the construction and since in the first case  $z_0$  lies between  $x_1, x_2 \in A_m$ , in the second one on the edge  $x_0 \in A_m$ , we see that (1) or (2) of  $(c_{n+1})$  again holds.

We can suppose that  $R_0 = Q^2$ , then put  $F_0 = 0, A_0 = \{0\}$ . From the lemma it follows that there exists a sequence  $\{(F_n, A_n); n \in FN\}$  with the desired properties. Let  $\{(F_\alpha, A_\alpha); \alpha \in \gamma\}$  be an  $Sd^*$ -prolongation consisting of  $Sd^*$ -functions  $F_\alpha : Q \to Q$  and  $Sd^*$ -maximal  $R_\alpha$ -nets  $A_\alpha$  such that  $(a_\alpha), (b_\alpha), (c_\alpha)$  for  $\alpha \in \gamma$  hold.

Take an  $\alpha \in \gamma - FN$  and put  $F \equiv F_{\alpha}$ . It remains to prove that  $R = R_F$ . Firstly observe that F is bounded, indeed  $|F(x)| \leq 1/3$  for  $x \in Q$  as follows from  $(c_{\alpha})$  with m = 0.

Let  $\langle x, y \rangle \in R$  and  $x_0, y_0 \in A_n$  be such that  $x \in R''_n\{x_0\}, y \in R''_n\{y_0\}$  and  $x_0 = y_0$ or  $x_0, y_0$  are connected neighbours in  $A_n$ . Necessarily there are such  $x_0, y_0$ . From  $(a_{\alpha}), (b_n)$  and  $(c_{\alpha})$  it follows

$$|F(x) - F(y)| \le |F_{\alpha}(x) - F_{\alpha}(x_0)| + |F_n(x_0) - F_n(y_0)| + |F_{\alpha}(y) - F_{\alpha}(y_0)| \le \frac{1}{2^n} + \frac{1}{2(1/2^n + 1/(3 \cdot 4^n))}.$$

Since it holds for each  $n \in FN$ ,  $F(x) \doteq F(y)$ . We have proved generally  $(\forall x, y)$  $(\langle x, y \rangle \in R \to F(x) \doteq F(y))$ . Thus  $\langle x, y \rangle \in R$  implies  $\langle x, y \rangle \in R_F$ . On the other hand let  $\langle x, y \rangle \notin R, x < y$ . If there are  $n \in FN$  and  $x_0, y_0 \in A_n$ such that  $x \leq x_0 < y_0 \leq y$ , then  $(a_\alpha)$  and  $(b_\alpha)$  imply that  $\langle x, y \rangle \notin R_F$ . Let  $\operatorname{card}(A_n \cap [x, y]) \leq 1$  for all n. Firstly let us suppose that  $A_n \cap (x, y) = 0$  for all n. Necessarily  $\langle x, y \rangle \notin S$ . Let n be such that  $\langle x, y \rangle \notin S_n$  and  $x_0, y_0 \in A_n$  such that  $x \in R''_n\{x_0\}, y \in R''_n\{y_0\}$ . Then  $x_0 \leq x < y \leq y_0$  and x lies on the edge  $x_0 \in A_n, y$ lies on the edge  $y_0 \in A_n$ . Since  $x_0, y_0 \in A_n$  are neighbouring,  $(b_n)$  and  $(a_\alpha)$  imply

$$|F(x_0) - F(y_0)| \ge 1/4^n$$

Finally from (2) of  $(c_{\alpha})$  it follows

$$|F(x) - F(y)| \ge |F(x_0) - F(y_0)| - |F(x_0) - F(x)| - |F(y) - F(y_0)| \ge \\ \ge 1/4^n - 2/(3 \cdot 4^n) = 1/(3 \cdot 4^n).$$

Thus  $\langle x, y \rangle \notin R_F$ . If  $A_n \cap \langle x, y \rangle = \{x_0\}$  for an  $n \in FN$  then  $A_m \cap \langle x, y \rangle = \{x_0\}$  for all  $m \ge n$ . Obviously  $\langle x, x_0 \rangle \notin R$  or  $\langle x_0, y \rangle \notin R$ . Since  $\operatorname{card}(A_n \cap \langle x, x_0 \rangle) = \operatorname{card}(A_n \cap \langle x, y \rangle) = 0$  for all  $n \in FN$ , it holds  $\langle x, x_0 \rangle \notin R_F$  or  $\langle x_0, y \rangle \notin R_F$ . This implies  $\langle x, y \rangle \notin R_F$ .

**Corollary.** Let R be an interval  $\pi$ -equivalence. Then there exists a nondecreasing Sd<sup>\*</sup>-function F such that  $R = R_F$  iff R is compact.

**PROOF:** It is obvious that if F is nondecreasing, then  $R_F$  is compact. Let R be compact. By the preceding theorem there exists an  $Sd^*$ -function G such that  $R = R_G$ . It would suffice to construct a "variation" of the function G. But we know that even a classically continuous function has not to have a variation. Nevertheless, in this case it suffices to prove the following

**Lemma.** Let G be a compact rational Sd-function (it means that  $R_G$  is compact). Then there exists its generalized variation, i.e. a nondecreasing Sd-function F such that  $R_G = R_F$ .

**PROOF:** Put 
$$\doteq_{\alpha} = \{ \langle x, y \rangle; |x - y| < 1/\alpha \text{ or } x, y > \alpha \text{ or } x, y < -\alpha \},\$$

$$R_{\alpha} = \{ \langle x, y \rangle; (\forall z \text{ between } x, y) (\langle G(x), G(z) \rangle \in \doteq_{\alpha} \& \langle G(z), G(y) \rangle \in \doteq_{\alpha}) \}.$$

Let  $\gamma > FN$  be such that for each  $\alpha \in \gamma, \alpha \ge 1$  there exists a maximal  $R_{\alpha}$ -net  $u_{\alpha}$ . If  $G(x) \doteq c$  for all  $x \in Q$  then put  $F(x) \equiv c$ . If there are x, y such that  $G(x) \neq G(y)$  then we can suppose that there are x, y such that G(x) = 0, G(y) = 1. Thus we can suppose that there are  $x_0 < y_0$  such that

$$(\forall \alpha \in \gamma)(x_0 = \min(u_\alpha \& y_0 = \max(u_\alpha)) \text{ and } u_\alpha = \{x_0 = x_0^\alpha < \cdots < x_{\omega_\alpha}^\alpha = y_0\}.$$

Put

$$F_{oldsymbol{lpha}}(x^{oldsymbol{lpha}}_{oldsymbol{eta}}) = \sum_{oldsymbol{\delta}=1}^{oldsymbol{eta}} |G(x^{oldsymbol{lpha}}_{oldsymbol{\delta}}) - G(x^{oldsymbol{lpha}}_{oldsymbol{\delta}-1})|.$$

Correspondence between interval  $\pi$ -equivalences and Sd-functions

If  $\langle x_{\beta}^{\alpha}, x_{\beta+1}^{\alpha} \rangle \in S_{\alpha}$ , then let  $F_{\alpha}$  be linear on  $[x_{\beta}^{\alpha}, x_{\beta+1}^{\alpha}]$ . If  $x_{\beta}^{\alpha}$  is an edge of  $S_{\alpha}$ , let  $F_{\alpha}$  be constant on  $S_{\alpha}^{\prime\prime}\{x_{\beta}^{\alpha}\} \cap [x_{\beta}^{\alpha}, \infty)$  or  $S_{\alpha}^{\prime\prime}\{x_{\beta}^{\alpha}\} \cap (-\infty, x_{\beta}^{\alpha}]$ . Finally put

$$F(x) = \sum_{lpha \in \gamma, lpha \geq 1} F_{lpha}(x) / (2^{lpha} F_{lpha}(y_0)) ext{ for } x \in Q$$

F is rational nondecreasing Sd-function such that  $0 \le F \le 2$ . We want to prove  $R_G = R_F$ .

If  $\langle x, y \rangle \in R_G$ , then  $F_{\alpha}(x) \doteq F_{\alpha}(y)$  for all  $\alpha$  and so  $F(x) \doteq F(y)$ , i.e.  $\langle x, y \rangle \in R_F$ . Let  $\langle x, y \rangle \notin R_G$  and let  $F(x) \doteq F(y)$ . Let there be  $x_n, y_n \in u_n$  such that  $x \leq x_n < y_n \leq y$ , then  $F(x_n) \doteq F(y_n)$ . From the construction of F it follows that  $G(x_n) \doteq G(z) \doteq G(y_n)$  for all  $z \in (u_m \cap [x_n, y_n]), m \geq n$  and thus  $G(x_n) \doteq G(z) \doteq G(y_n)$  for all  $z \in [x_n, y_n]$ . It means  $\langle x_n, y_n \rangle \in R_G$  - a contradiction. Thus card( $[x, y] \cap u_n) \leq 1$  for all  $n \in FN$ . This implies  $\langle x, y \rangle \notin S$  where S is the relation of connectedness of  $R_G$ . Thus we have  $G(x) \neq G(y)$  and  $G(x) \doteq G(z)$  for  $z \in S''\{x\} \cap [x, y]$  and  $G(z) \doteq G(y)$  for  $z \in S''\{y\} \cap [x, y]$ . Consequently  $F_n(x) \neq F_n(y)$  for an  $n \in FN$  and so  $\langle x, y \rangle \notin R_F$ .

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