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## OPERATIONS ON GRAPHS DETERMINING CONGRUENCES ON GRAPHS

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The purpose of this paper is to characterize by means of concepts and operations of graph theory partitions of the elements of a finite modular lattice  $H$  that determine congruence relations on  $H$ . By the aid of the characterization we construct thereafter a class of congruence relations on graphs. We recall first some concepts of graph theory and apply thereafter them to the Hasse diagram of  $H$  in order to obtain the characterization.

We shall consider finite undirected and connected graphs  $G = (P(G), L(G))$  only without loops and multiple lines, where  $P(G)$  is the set of points of  $G$  and  $L(G)$  its set of lines.  $SP$  is a mapping  $P(G) \times P(G) \rightarrow 2^{P(G)}$  defined as follows:

$$SP(x, y) = \{z \mid z \in P(G) \text{ and } z \text{ is on a shortest path joining } x \text{ and } y \text{ in } G\}.$$

We shall call  $SP$  a binary operation on  $P(G)$ , although the mapping induced by the operation is a one-to-many mapping, as the name operation helps us to find some useful analogies we shall apply. In particular,  $\{x, y\} \subseteq SP(x, y)$  and  $SP(x, x) = \{x\}$ ,  $x, y \in P(G)$ . In general, let  $U$  and  $W$  be two subsets of  $P(G)$ , then  $SP(U, W)$  denotes the union of the sets  $SP(u, w)$ , where  $u \in U$  and  $w \in W$ ; formally  $SP(U, W) = \{z \mid z \in SP(u, w) \text{ for some } u \text{ and } w, u \in U \text{ and } w \in W\}$ . A set  $U \subset P(G)$  is called an ideal of  $G$ , if  $U \neq \emptyset$  and  $SP(U, U) = U$ . By the notation  $SP^n(x, y)$  we denote the operation  $SP(SP^{n-1}(x, y), SP^{n-1}(x, y))$ . Thus  $SP^2(x, y) = SP(SP(x, y), SP(x, y))$ . As we consider finite graphs only, there is for any pair  $x, y \in P(G)$  a value of  $n$  such that  $SP^n(x, y)$  is an ideal of  $G$ . The graph of Figure 1 illuminates the case where  $SP^2(x, y)$  is not an ideal of  $G$  but  $SP^3(x, y)$  is. It is important to construct from a pair  $x, y \in P(G)$  an ideal of  $G$  by means of sequential applying of the  $SP$ -operation and in order to use a brief notation,  $SU(x, y)$  denotes the ideal obtained from  $x, y$  by applying the  $SP$ -operation enough many times.

Ideals of graphs and the  $SP$ -operation were introduced in [4] and briefly considered in [5]. These concepts are natural generalizations of corresponding concepts defined for trees by Nebeský in [3].

In this paper we consider the Hasse diagram of a lattice  $H$  as an undirected graph and denote it by  $G_H$ . Lemma 1 and Theorem 1 are proved in a more general form than we need later. A lattice  $H$  is locally finite, if its every interval is finite.

**Lemma 1.** *Let  $H$  be a locally finite lattice. Then  $SU(x, y) = [x \wedge y, x \vee y]$  for any two elements  $x, y \in H$  if and only if  $H$  is modular.*

*Proof.* If  $H$  is modular, then according to the metric properties of finite modular lattices,  $x \wedge y, x \vee y \in SP(x, y)$  (see e.g. Draškovičová [1]). As the lengths of any two chains between  $a$  and  $b$  in a finite modular lattice are equivalent when  $a < b$ , each  $z \in [x \wedge y, x \vee y]$  belongs to a shortest path from  $x \wedge y$  to  $x \vee y$  and so  $z \in SU(x, y)$ . Obviously  $SU(x, y) \subseteq [x \wedge y, x \vee y]$ , and thus  $SU(x, y) = [x \wedge y, x \vee y]$ .

Let  $H$  satisfy the condition of the lemma for any pair  $x, y \in H$ . If  $H$  were non-modular, then it contains the well known non-modular sublattice (in Figure 2 the sublattice of elements  $a, b, c, d, e$ ), where the set  $\{a, b, c\} = SU(b, c) \neq [a, e] = [b \wedge c, b \vee c]$ . This completes the proof.

Now we are ready to prove the characterization.

**Theorem 1.** *Let  $H$  be a locally finite modular lattice and  $\mathfrak{C} = \{C_1, \dots, C_m\}$  a partition of its elements.  $\mathfrak{C}$  is a congruence partition of  $H$  with respect to the operations  $\vee$  and  $\wedge$  on  $H$  if and only if the condition (A) holds.*

(A) *If  $x, y \in C_i$  and  $a, b \in C_j$  in  $\mathfrak{C}$ , then  $SU(x, a) \cap C_k \neq \emptyset$  holds for some  $k$  in  $G_H$  if and only if  $SU(y, b) \cap C_k \neq \emptyset$  holds,  $1 \leq k \leq m$ .*

*Proof.* Assume that  $\mathfrak{C}$  is a partition of the points  $P(G_H)$  such that  $SU(x, a) \cap C_k \neq \emptyset$  if and only if  $SU(y, b) \cap C_k \neq \emptyset$ . We show that  $R$  is a latticecongruence on  $H$ , with the classes  $C_1, \dots, C_m$ . Clearly  $R$  is reflexive, symmetric and transitive. Thus it remains to show the compatibility of  $R$ , i.e. to show that  $xRy$  implies  $x \wedge zRy \wedge z$  and  $x \vee zRy \vee z$  for any  $z \in H$ . Moreover, if  $qRp \Leftrightarrow q \wedge pRq \vee p$ , we may assume that  $x \leq y$ .

Let  $x < y$  (the case  $x = y$  is trivial),  $xRy$  and  $z \in H$ . Thus  $x \vee z \leq y \vee z$ . We assume that in the partition  $\mathfrak{C}$  of  $Hx \vee z$  and  $y \vee z$  belong to different sets of  $\mathfrak{C}$ . As  $x \leq y \wedge (x \vee z) \leq y$  and  $yRx, y \wedge (x \vee z)Ry$  holds, too. The relations  $y \vee zRy \vee z$  and  $yRx$  imply that (A) holds for  $SU(y \vee z, y)$  and  $SU(y \vee z, x)$ .  $x \vee z \in SU(y \vee z, x)$  and we assume that  $x \vee z \in C_h$ . As  $x \vee z < y \vee z, x \vee z \notin SU(y \vee z, y)$ . Then according to (A),  $SU(y \vee z, y) \cap C_h \neq \emptyset$ , and let  $t$  be the greatest element of the set  $SU(y \vee z, y) \cap C_h$ ; such an element exists as  $SU(y \vee z, y)$  is finite and for any two elements of  $SU(y \vee z, y)$  (of  $C_h$ ),  $SU(y \vee z, y) \cap C_h$  contains the join of these elements. But  $x \vee z \vee t \in C_h$  and  $x \vee z \vee t \leq y \vee z$ , whence  $x \vee z \vee t \in SU(y \vee z, y)$ . Thus we can assume that  $x \vee z \leq t$ , and as  $t \in SU(y \vee z, y)$ ,  $t \geq y$ . But then  $y \vee x \vee z = y \vee z \leq t$ , whence  $y \vee z, x \vee z \in C_h$ , which is a contradiction. Hence  $y \vee z, x \vee z \in C_k$  for some value  $k$  of  $i$ . The proof is similar for  $y \wedge zRx \wedge z$ .

Conversely, we assume that  $\mathfrak{C}$  generates a latticecongruence on  $H$ . Let  $x, y \in C_i$ ,

$a, b \in C_j$  and  $i \neq j$ . Accordingly, we may assume that  $x \leq y$  and  $a \leq b$ . As  $R$  is a congruence relation on  $H$ ,  $y \vee bRx \vee a$ . If there is an element  $q \in C_k$ ,  $y \leq q \leq \leq y \vee b$ , then  $x \leq q \wedge (x \vee a) \leq x \vee a$ ,  $q \wedge (x \vee a) Rq$  and thus  $q \wedge (x \vee a) \in C_j$ . By applying this technique to the intervals  $[y \wedge b, y \vee b]$  and  $[x \vee a, x \wedge a]$  we see that the condition (A) holds for  $SU(y, b)$  and  $SU(x, a)$ . The proof is similar for  $SU(x, b)$  and  $SU(y, a)$ . This completes the proof.

As the example of Figure 2 shows, a partition of a non-modular lattice  $H$  satisfying the condition (A) need not be either a  $\wedge$ -congruence or a  $\vee$ -congruence on  $H$ .

In the next theorem we show how the condition (A) generalizes by a natural way the construction of compatible tolerances on graphs introduced by Zelinka in [7].

We call a binary, reflexive, symmetric and transitive relation  $R$  on a graph a  $SU$ -compatible congruence relation on  $G$  when  $aRb$  and  $xRy$  imply  $SU(a, x)RSU(b, y)$ . The notation  $SU(a, x)RSU(b, y)$  means that for any  $z \in SU(a, x)$  there is a point  $u \in SU(b, y)$  such that  $zRu$ , and for any  $w \in SU(b, y)$  there is a point  $v \in SU(a, x)$  such that  $vRw$ .

**Theorem 2.** *Let  $\mathfrak{C}$  be a partition of the pointset  $P(G)$  of a graph  $G = (P(G), L(G))$ . The relation  $R$  given by  $\mathfrak{C}$  determines a  $SU$ -compatible congruence relation on  $G$  if and only if  $\mathfrak{C}$  satisfies the condition (A).*

*Proof.* If  $\mathfrak{C}$  is a partition of  $P(G)$  such that the relation  $R$  given by  $\mathfrak{C}$  satisfies the condition (A), the  $SU$ -compatibility of  $R$  follows directly from (A). The converse proof follows similarly directly from the definition of the  $SU$ -compatibility.

By using the terminology of Theorem 2, we can say, according to Theorem 1, that  $R$  is a latticecongruence on a finite modular lattice  $H$  if and only if  $R$  is a  $SU$ -compatible congruence relation on  $G_H$ .

We obtain also a characterization of finite modular lattices as given in the next theorem.

**Theorem 3.** *Let  $H$  be a finite lattice and  $\mathfrak{C}$  a partition of  $H$  determining a  $SU$ -compatible relation  $R$  on  $G_H$ .  $H$  is modular if and only if each  $R$  defined above is a latticecongruence on  $H$ .*

*Proof.* If  $H$  is modular, then the assertion follows from Theorems 1 and 2. Thus let each  $R$  of the theorem be a congruence relation on  $H$ . If  $H$  is non-modular, it contains as a sublattice the lattice of the elements  $a, b, c, d, e$  in Figure 2, where the subset  $\{a, b, c\}$  of the partition  $\mathfrak{C} = \{\{a, b, c\}, \{d, e\}\}$  shows that  $\mathfrak{C}$  does not determine a congruence relation on  $H$  although  $R$  is  $SU$ -compatible on  $G_H$ .

As a model for constructing a  $SU$ -compatible congruence on  $G$  were latticecongruences on a finite modular lattice  $H$ . This model is used in the following theorem where an analogy is presented between  $SU$ -compatible congruences on  $G$  and congruences on algebras. Its proof is a direct copy of the corresponding proof for algebras given e.g. in [6, Thm. 96 and its supplement], and hence we omit it.

Theorem 4. Let  $G$  be a given graph.  $G$  is a Cartesian product of two non-trivial graphs  $G_1$  and  $G_2$ , i.e.  $G = G_1 \times G_2$ , if and only if there are two non-trivial SU-compatible congruences  $R_1, R_2 \in H(G)$  which are permutable and complements of each other in  $H(G)$ .  $H(G)$  is the lattice of SU-compatible congruences on  $G$ .

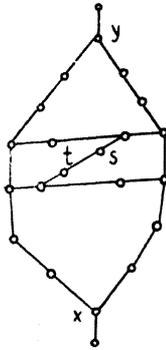


Fig. 1

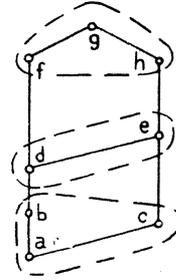


Fig. 2

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