

Ladislav Nebeský

A new proof of a characterization of the set of all geodesics in a connected graph

Czechoslovak Mathematical Journal, Vol. 48 (1998), No. 4, 809–813

Persistent URL: <http://dml.cz/dmlcz/127456>

Terms of use:

© Institute of Mathematics AS CR, 1998

Institute of Mathematics of the Czech Academy of Sciences provides access to digitized documents strictly for personal use. Each copy of any part of this document must contain these *Terms of use*.



This document has been digitized, optimized for electronic delivery and stamped with digital signature within the project *DML-CZ: The Czech Digital Mathematics Library* <http://dml.cz>

A NEW PROOF OF A CHARACTERIZATION OF THE SET OF
ALL GEODESICS IN A CONNECTED GRAPH

LADISLAV NEBESKÝ, Praha

(Received September 13, 1996)

In [2], the present author gave a characterization of the set of all geodesics (or shortest paths) in a connected graph G . More precisely, he gave a necessary and sufficient condition for a set of paths in G to be the set of all geodesics in G . The proof of necessity is easy and was omitted in [2]. But the proof of sufficiency given there was rather long.

This characterization was partially modified in [3]; the proof given there was also long (in fact, the characterization was derived from a more general theorem proved there). In the present paper we present its new and shorter proof. The proof utilizes a new lemma, which yields a deeper insight into the idea of the characterization.

Let G be a (finite undirected) connected graph (without loops and multiple edges), and let V , E and D denote its vertex set, its edge set and its diameter, respectively. If $u, v \in V$, then $d(u, v)$ denotes the distance between u and v in G . (The letters g, h, \dots, n will be used to denote integers).

We denote by Σ_N the set of all sequences

$$(0) \quad u_0, \dots, u_g$$

where $u_0, \dots, u_g \in V$ and $g \geq 0$. Similarly as in [2] and [3], instead of (0) we will write $u_0 \dots u_g$. Let $\alpha = v_0 \dots v_h$, where $v_0, \dots, v_h \in V$ and $h \geq 0$. We write $A\alpha = v_0$, $Z\alpha = v_h$, $\|\alpha\| = h$ and

$$\bar{\alpha} = v_h \dots v_0.$$

Let $\beta = x_0 \dots x_i$ and $\gamma = y_0 \dots y_j$, where $x_0, \dots, x_i, y_0, \dots, y_j \in V$, $i \geq 0$ and $j \geq 0$; we write

$$\beta\gamma = x_0 \dots x_i y_0 \dots y_j.$$

We denote by $*$ the empty sequence in the sense that $*\delta = \delta = \delta*$ for each $\delta \in \Sigma_N$. Put $** = *$ and $\bar{*} = *$. Define $\Sigma = \Sigma_N \cup \{*\}$.

Let $\pi \in \Sigma_N$. We say that π is a path in G if there exist $k \geq 0$ and mutually distinct $w_0, \dots, w_k \in V$ such that $\pi = w_0 \dots w_k$ and if $k \geq 1$, then

$$\{w_0, w_1\}, \dots, \{w_{k-1}, w_k\} \in E.$$

The set of all paths in G will be denoted by \mathcal{P} . If $\mathcal{Q} \subseteq \mathcal{P}$ and $m \geq 0$, then we define

$$\mathcal{Q}(m) = \{\omega \in \mathcal{Q}; d(A\omega, Z\omega) = m\}.$$

Obviously, if $\pi \in \mathcal{P}$, then $d(A\pi, Z\pi) \leq \|\pi\|$.

Let $\tau \in \Sigma_N$. We say that τ is a *geodesic* (or a shortest path) in G if $\tau \in \mathcal{P}$ and $d(A\tau, Z\tau) = \|\tau\|$.

The following theorem gives a characterization of the set of all geodesics in G .

Theorem ([3]). *Let $\mathcal{R} \subseteq \mathcal{P}$, and let Γ denote the set of all geodesics in G . Then the statements (1) and (2) are equivalent:*

(1) $\mathcal{R} = \Gamma$.

(2) \mathcal{R} satisfies the following properties **A**(\mathcal{R}) – **G**(\mathcal{R}) (for all $u, v, x, y \in V$ and all $\varphi, \psi \in \Sigma$):

A(\mathcal{R}) if $uv\varphi x \in \mathcal{R}$, then $\{u, x\} \notin E$;

B(\mathcal{R}) if $uv\varphi x \in \mathcal{R}$, then $x\bar{\varphi}vu \in \mathcal{R}$;

C(\mathcal{R}) if $uv\varphi x \in \mathcal{R}$, then $v\varphi x \in \mathcal{R}$;

D(\mathcal{R}) if $uv\varphi x, v\psi x \in \mathcal{R}$, then $uv\psi x \in \mathcal{R}$;

E(\mathcal{R}) if $uv\varphi x, v\psi y \in \mathcal{R}$ and $\{x, y\} \in E$, then $v\varphi xy \in \mathcal{R}$;

F(\mathcal{R}) if $uv\varphi x \in \mathcal{R}$, $\{x, y\} \in E$, $uv\varrho y \notin \mathcal{R}$ for all $\varrho \in \Sigma$ and $u\sigma y x \notin \mathcal{R}$ for all $\sigma \in \Sigma$, then $v\varphi xy \in \mathcal{R}$;

G(\mathcal{R}) there exists $\xi \in \mathcal{R}$ such that $A\xi = u$ and $Z\xi = x$.

We will present a new proof of the theorem. The proof that (1) implies (2) is not complicated and will be omitted here. We only prove that (2) implies (1).

The next lemma yields a deeper insight into the theorem and suggest a new method for proving it.

Lemma. *Let $u_0, u_1, \dots, u_{g+h-1} \in V$, where $\min(g, h) \geq 2$. Denote $j = \min(g, h)$ and $u_{g+h} = u_0, u_{g+h+1} = u_1, \dots, u_{g+h+j} = u_j$. Moreover, denote*

$$\alpha_i = u_i u_{i+1} \dots u_{i+g} \text{ and } \beta_i = u_{i+g} u_{i+g+1} \dots u_{i+g+h}$$

for each $i, 0 \leq i \leq j$.

Let $\mathcal{Q}, \mathcal{T} \subseteq \mathcal{P}$. Assume that \mathcal{Q} satisfies $\mathbf{B}(\mathcal{Q}) - \mathbf{F}(\mathcal{Q})$ and \mathcal{T} satisfies $\mathbf{B}(\mathcal{T}) - \mathbf{E}(\mathcal{T})$. Next, assume that the following conditions I–IV hold for all $i, 0 \leq i \leq j$, and all $\varphi, \psi \in \Sigma$:

- I if $\alpha_i \in \mathcal{Q}$ and $\beta_i \notin \mathcal{T}$, then $\alpha_i \in \mathcal{T}$;
- II if $u_i u_{i+1} \varphi u_{i+g+1} \in \mathcal{Q}$, $\beta_i \in \mathcal{T}$ and $\alpha_{i+1} \notin \mathcal{Q}$, then $u_i u_{i+1} \varphi u_{i+g+1} \in \mathcal{T}$;
- III if $u_i u_{i+1} \varphi u_{i+g} \in \mathcal{T}$ and $u_i u_{i+1} \psi u_{i+g} \in \mathcal{Q}$, then $u_{i+1} \varphi u_{i+g} \in \mathcal{Q}$;
- IV if $u_{i+g} u_{i+g+1} \varphi u_{i+g+h} \in \mathcal{T}$ and $u_{i+g} u_{i+g+1} \psi u_{i+g+h} \in \mathcal{Q}$, then $u_{i+g+1} \varphi u_{i+g+h} \in \mathcal{Q}$.

Finally, assume that $\alpha_0 \in \mathcal{Q}$ and $\beta_0 \in \mathcal{T}$. Then $\beta_0 \in \mathcal{Q}$.

P r o o f. Suppose, to the contrary, that $\beta_0 \notin \mathcal{Q}$. Then $\beta_0 \in \mathcal{T} - \mathcal{Q}$. First, we will show that

$$(3) \quad \text{either } \alpha_j \notin \mathcal{Q} \text{ or } \beta_j \notin \mathcal{T}.$$

If $g = h$, then (3) immediately follows from the fact that $\beta_0 \notin \mathcal{Q}$. Next, let $g > h$. Then $j = h$. Suppose that $\alpha_j \in \mathcal{Q}$. Applying $\mathbf{B}(\mathcal{Q})$ to α_0 and $\mathbf{C}(\mathcal{Q})$ to α_j , we get

$$u_g u_{g-1} \dots u_0, u_{g-1} u_g \dots u_{g+h} \in \mathcal{Q}$$

Recall that $u_{g+h} = u_0$. By $\mathbf{D}(\mathcal{Q})$,

$$u_g u_{g-1} u_g \dots u_{g+h} \in \mathcal{Q}.$$

Thus $\mathcal{Q} - \mathcal{P} \neq \emptyset$, a contradiction. We get $\alpha_j \notin \mathcal{Q}$ and therefore (3) holds. Finally, let $h > g$. Then $j = g$. In a similar way, we get $\beta_j \notin \mathcal{T}$. Thus (3) holds again.

Recall that $\alpha_0 \in \mathcal{Q}$ and $\beta_0 \in \mathcal{T} - \mathcal{Q}$. Combining this fact with (3), we see that there exists $k, 0 \leq k < j$ such that

$$(4) \quad \alpha_k \in \mathcal{Q} \text{ and } \beta_k \in \mathcal{T} - \mathcal{Q}$$

and

$$(5) \quad \alpha_{k+1} \notin \mathcal{Q} \text{ or } \beta_{k+1} \notin \mathcal{T} - \mathcal{Q}.$$

Denote $\alpha = \alpha_k, \alpha' = \alpha_{k+1}, \beta = \beta_k, \beta' = \beta_{k+1}$. Clearly, $\alpha, \beta \in \mathcal{P}$. We distinguish two cases:

1. Let $\alpha' \in \mathcal{Q}$. If $\alpha' \notin \mathcal{T}$, then I implies that $\beta' \in \mathcal{T}$. If $\alpha' \in \mathcal{T}$, then combining $\mathbf{B}(\mathcal{T})$ and $\mathbf{E}(\mathcal{T})$, we get $\beta' \in \mathcal{T}$ again. By (5), $\beta' \notin \mathcal{T} - \mathcal{Q}$. Hence $\beta' \in \mathcal{Q}$. By (4), $\alpha \in \mathcal{Q}$. Combining $\mathbf{B}(\mathcal{Q})$ with $\mathbf{E}(\mathcal{Q})$, we get $\beta \in \mathcal{Q}$, which contradicts (4).

2. Let $\alpha' \notin \mathcal{Q}$. Denote $u = u_k, v = u_{k+1}, x = u_{k+g}$ and $y = u_{k+g+1}$. Clearly, there exist $\varphi, \tau \in \Sigma$ such that $\alpha = uv\varphi x$ and $\beta = xy\tau u$. Hence $\alpha' = v\varphi xy$. We distinguish two subcases:

2.1. Let $u\sigma yx \notin \mathcal{Q}$ for all $\sigma \in \Sigma$. By virtue of $\mathbf{F}(\mathcal{Q})$, there exists $\varrho \in \Sigma$ such that $uv\varrho y \in \mathcal{Q}$. Since $\beta \in \mathcal{T}$ and $\alpha' \notin \mathcal{Q}$, II implies that $uv\varrho y \in \mathcal{T}$. By $\mathbf{B}(\mathcal{T})$, $y\bar{\varrho}vu \in \mathcal{T}$. Recall that $xy\tau u \in \mathcal{T}$. By $\mathbf{D}(\mathcal{T})$, $xy\bar{\varrho}vu \in \mathcal{T}$ and by $\mathbf{B}(\mathcal{T})$, $uv\varrho yx \in \mathcal{T}$. Since $uv\varphi x \in \mathcal{Q}$, III implies that $v\varrho yx \in \mathcal{Q}$. By $\mathbf{D}(\mathcal{Q})$, $uv\varrho yx \in \mathcal{Q}$ and by $\mathbf{B}(\mathcal{Q})$, $xy\bar{\varrho}vu \in \mathcal{Q}$. Since $xy\tau u \in \mathcal{T}$, IV implies that $y\tau u \in \mathcal{Q}$. By $\mathbf{D}(\mathcal{Q})$, $\beta = xy\tau u \in \mathcal{Q}$, which contradicts (4).

2.2. Let there exist $\sigma \in \Sigma$ such that $u\sigma yx \in \mathcal{Q}$. By $\mathbf{B}(\mathcal{Q})$, $xy\bar{\sigma}u \in \mathcal{Q}$. Since $\beta \in \mathcal{T}$, IV implies that $y\tau u \in \mathcal{Q}$. By $\mathbf{D}(\mathcal{Q})$, $\beta \in \mathcal{Q}$, which contradicts (4) again.

Thus $\beta_0 \in \mathcal{Q}$, which completes the proof of the lemma. \square

Proof of the theorem. We will only prove that (2) implies (1). Now, let (2) hold. We will prove that $\Gamma(n) \subseteq \mathcal{R}(n)$ and $\mathcal{R}(n) \subseteq \Gamma(n)$ for every $n \geq 0$. We proceed by induction on n . The fact that $\Gamma(0) = \mathcal{R}(0)$ follows from $\mathbf{G}(\mathcal{R})$. The fact that $\Gamma(1) = \mathcal{R}(1)$ follows from $\mathbf{G}(\mathcal{R})$ and $\mathbf{A}(\mathcal{R})$. Let $n \geq 2$. Assume that

$$(6) \quad \Gamma(m) = \mathcal{R}(m) \text{ for each } m, 0 \leq m < n.$$

The case when $D < n$ is trivial. Suppose that $D \geq n$.

Consider an arbitrary $\omega \in \Gamma(n)$. Put $\mathcal{Q} = \mathcal{R}, \mathcal{T} = \Gamma$ and $h = n$. Obviously, $\mathbf{G}(\mathcal{Q})$ holds. There exist $u_0, \dots, u_{g+h-1} \in V$, where $g \geq h$, such that

$$(7) \quad u_0 u_1 \dots u_g \in \mathcal{Q} \text{ and } \omega = u_g u_{g+1} \dots u_{g+h}, \text{ where } u_{g+h} = u_0.$$

By virtue of (6), I–IV hold. According to the lemma, $\omega \in \mathcal{R}$. We have proved that $\Gamma(n) \subseteq \mathcal{R}(n)$.

Consider an arbitrary $\omega \in \mathcal{R}(n)$. Put $\mathcal{Q} = \Gamma, \mathcal{T} = \mathcal{R}$ and $g = n$. There exist $u_0, \dots, u_{g+h-1} \in V$, where $h \geq g$, such that (7) holds. Combining (6) with the fact that $\Gamma(n) \subseteq \mathcal{R}(n)$, we see that I–IV hold. According to the lemma, $\omega \in \Gamma$. We have proved that $\mathcal{R}(n) \subseteq \Gamma(n)$, which completes the proof of the theorem. \square

Remark 1. The fact that V is finite was not utilized in any point of our proof.

Remark 2. A different way of characterizing the set of all geodesics in a connected graph can be found in [4].

Remark 3. Some types of graphs can be characterized by counting geodesics. For this topic, see [1].

References

- [1] *P. V. Ceccherini*: Studying structures by counting geodesics. In: Combinatorial Designs and Applications (W. D. Wallis et al., eds.). Marcel Dekker, New York and Basel, 1990, pp. 15–32.
- [2] *L. Nebeský*: A characterization of the set of all shortest paths in a connected graph. *Math. Bohemica* 119 (1994), 15–20.
- [3] *L. Nebeský*: On the set of all shortest paths of a given length in a connected graph. *Czechoslovak Math. Journal* 46 (121) (1996), 155–160.
- [4] *L. Nebeský*: Geodesics and steps in a connected graph. *Czechoslovak Math. Journal* 47 (122) (1997), 149–161.

Author's address: Nám. J. Palacha 2, 116 38 Praha 1, Czech Republic (Filozofická fakulta Univerzity Karlovy).