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Mathematica Slovaca, Vol. 31 (1981), No. 1, 71--78

Persistent URL: <http://dml.cz/dmlcz/129345>

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THE ORIENTABILITY OF THE DIRECT PRODUCT OF GRAPHS

EVA GEDEONOVÁ

The covering graph $C(P)$ of a partially ordered set P is the graph whose vertices are the elements of P and whose edges are those pairs $\{a, b\}$, $a, b \in P$, for which a covers b or b covers a . The covering graph $C(P)$ of a partially ordered set P with some properties determines certain further properties of P . In some cases if $C(P)$ is a direct product of graphs G_1, G_2 , then the partially ordered set P is a direct product of partially ordered sets P_1, P_2 and $C(P_i) = G_i$, $i = 1, 2$. The following example shows a case in which this assertion does not hold. The covering graph of the partially ordered set P of Fig. 1 is the direct product of two twoelemented graphs but the partially ordered set is not a direct product of two twoelemented partially ordered sets.

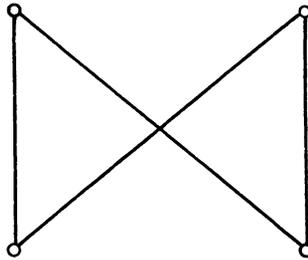


Fig. 1

The direct product $G_1 \times G_2$ of the graphs $G_1 = (V_1, H_1)$, $G_2 = (V_2, H_2)$ is the graph whose vertices are the elements of $V_1 \times V_2$ and whose edges are those pairs $\{(a_1, b_1), (a_2, b_2)\}$ $a_i \in V_1, b_i \in V_2, i = 1, 2$ satisfying either $a_1 = a_2$ and $\{b_1, b_2\} \in H_2$ or $\{a_1, a_2\} \in H_1$ and $b_1 = b_2$. By a graph isomorphism of graphs $G = (V, H)$ and $G' = (V', H')$ we mean a bijection $f: V \rightarrow V'$ of vertex sets such that $\{a, b\} \in H$ iff $\{f(a), f(b)\} \in H'$ for all $a, b \in V$. For vertices a and b of a graph G a path from a to b of length n is a sequence $a = c_0, c_1, \dots, c_n = b$ of vertices of G such that successive pairs in this sequence are joined by an edge of G . Let $d(a, b)$ denote the distance from a to b , i.e. the length of a shortest path from a to b . A graph is

connected if for all its vertices a, b there holds that a, b are connected by a path. Note that every graph isomorphism f of connected graphs is a distance isomorphism (i.e. $d(a, b) = d(f(a), f(b))$). A partially ordered set P is locally finite if for every $a, b \in P, a < b$ there is a finite maximal chain between a and b . If a locally finite partially ordered set P has the least element and all maximal chains in P between fixed endpoints have the same order, then we say that P is graded. In this case we define the height $h(a)$ of an element of P as the order of a maximal chain from the least element of P to a , minus one. For elements a and $b, a > b$, of a partially ordered set P we write $a \triangleright b$ or $b \triangleleft a$ (a covers b or b is covered by a) if $a \geq c > b$ implies $a = c$ for every element $c \in P$.

In the whole paper G_1, G_2 are graphs, P is a partially ordered set and if $f: G_1 \times G_2 \rightarrow C(P)$ is a graph isomorphism, then we denote elements of P by $f(a, b)$, where $a \in G_1, b \in G_2$. This is correct since f is a bijection. Note that if $d(x, y) < \infty, x, y \in G_1 \times G_2$, then $d(x, y) = d(f(x), f(y))$.

Theorem 1. *Let $f: G_1 \times G_2 \rightarrow C(P)$ be a graph isomorphism. Then there exist partially ordered sets P_1, P_2 such that G_i is graph isomorphic to $C(P_i), i = 1, 2$.*

Proof. Let $a_0 \in G_1, b_0 \in G_2$. If

$$P_1 = \{f(a, b_0), a \in G_1\}, \quad P_2 = \{f(a_0, b), b \in G_2\},$$

then $P_i \subset P, i = 1, 2$, hence P_i are partially ordered and it is easy to see that G_i and $C(P_i)$ are graph isomorphic.

Lemma 1. (Kotzig [2]). *Let a graph G be a direct product of the graphs G_1, G_2 , let $a, b \in G_1, c, d \in G_2$. Then*

$$d((a, c), (b, d)) = d(a, b) + d(c, d).$$

Definition 1. *Let $f: G_1 \times G_2 \rightarrow C(P)$ be a graph isomorphism. We say that f has the property \square if for every $a_1, a_2 \in G_1, b_1, b_2 \in G_2, d(a_1, a_2) = d(b_1, b_2) = 1$,*

$$f(a_1, b_1) \triangleleft f(a_2, b_1) \text{ implies } f(a_1, b_2) \triangleleft f(a_2, b_2)$$

and

$$f(a_1, b_1) \triangleleft f(a_1, b_2) \text{ implies } f(a_2, b_1) \triangleleft f(a_2, b_2).$$

Lemma 2. *Every graph isomorphism $f: G_1 \times G_2 \rightarrow C(L)$, where L is a lattice, has the property \square .*

Proof. $d(a_1, a_2) = d(b_1, b_2) = 1$ and $f(a_1, b_1) \triangleleft f(a_2, b_1)$. Since by Lemma 1 $d(f(a_2, b_1), f(a_2, b_2)) = 1, d(f(a_2, b_2), f(a_1, b_2)) = 1, d(f(a_1, b_2), f(a_1, b_1)) = 1$ and L is a lattice, there must be $f(a_1, b_2) \triangleleft f(a_2, b_2)$.

In the same way the second implication of Definition 1 can be proved.

Lemma 3. *Every graph isomorphism $f: G_1 \times G_2 \rightarrow C(P)$, where P is a graded partially ordered set, has the property \square .*

Proof Let $d(a_1, a_2) = d(b_1, b_2) = 1$, $a_1, a_2 \in G_1$, $b_1, b_2 \in G_2$ and let

$$f(a_1, b_1) \triangleleft f(a_2, b_1). \quad (1)$$

Since $d(f(a_1, b_2), f(a_2, b_2)) = 1$ there is either

$$f(a_1, b_2) \triangleleft f(a_2, b_2) \text{ or } f(a_2, b_2) \triangleleft f(a_1, b_2). \quad (2)$$

Let us suppose that

$$f(a_2, b_2) \triangleleft f(a_1, b_2). \quad (3)$$

Since $d(f(a_1, b_1), f(a_1, b_2)) = d(f(a_2, b_1), f(a_2, b_2)) = 1$ and P is a partially ordered set, it is easy to check that

$$f(a_1, b_1) \triangleleft f(a_1, b_2) \text{ and } f(a_2, b_2) \triangleleft f(a_2, b_1). \quad (4)$$

From (3) and (4) it follows that

$$h(f(a_1, b_1)) = h(f(a_2, b_2)) \text{ and } h(f(a_2, b_1)) = h(f(a_1, b_2)) \quad (5)$$

Let the least element of the partially ordered set P be $f(a_0, b_0)$. Since $h(f(a, b)) = d(f(a, b), f(a_0, b_0)) = d((a, b), (a_0, b_0))$, by Lemma 1 and by (5) we have

$$\begin{aligned} d(a_1, a_0) + d(b_1, b_0) &= d(a_2, a_0) + d(b_2, b_0), \\ d(a_2, a_0) + d(b_1, b_0) &= d(a_1, a_0) + d(b_2, b_0). \end{aligned}$$

From these equalities it follows that

$$d(a_1, a_0) - d(a_2, a_0) = d(a_2, a_0) - d(a_1, a_0).$$

Hence we have

$$d(a_1, a_0) = d(a_2, a_0). \quad (6)$$

Moreover, on the basis of $d(a_1, a_2) = 1$ we have either $f(a_1, b_0) \triangleleft f(a_2, b_0)$ or $f(a_2, b_0) \triangleleft f(a_1, b_0)$. If $f(a_1, b_0) \triangleleft f(a_2, b_0)$, then by Lemma 1 $d(a_2, a_0) = h(f(a_2, b_0)) = h(f(a_1, b_0)) + 1 = d(a_1, a_0) + 1$, which contradicts (6). Analogously $f(a_2, b_0) \triangleleft f(a_1, b_0)$ leads to contradiction, too.

Since supposition (3) does not hold, from (2) there follows the assertion. The second implication of Definition 1 can be proved in the same way.

Lemma 4. *Let a graph isomorphism $f: G_1 \times G_2 \rightarrow C(P)$ have the property \square . Then for any $a_1, a_2 \in G_1$, $b_1, b_2 \in G_2$, $d(a_1, a_2) < \infty$, $d(b_1, b_2) < \infty$*

$$f(a_1, b_1) \triangleleft f(a_2, b_1) \text{ implies } f(a_1, b_2) \triangleleft f(a_2, b_2)$$

and

$$f(a_1, b_1) \triangleleft f(a_1, b_2) \text{ implies } f(a_2, b_1) \triangleleft f(a_2, b_2).$$

Proof. We prove the first implication. If $d(b_1, b_2) = 1$, then the assertion follows by Definition 1, because $f(a_1, b_1) \triangleleft f(a_2, b_1)$ implies $d(a_1, a_2) = 1$. For the

second part of the induction we assume that the assertion is true for $d(b_1, b_2) = k$. If $d(b_1, b_2) = k + 1$, then there exists a path

$$b_1 = c_0, c_1, \dots, c_{k+1} = b_2, c_i \in G_2, 0 \leq i \leq k + 1.$$

Since $d(b_1, c_k) = k$, by the induction hypothesis $f(a_1, c_k) < f(a_2, c_k)$. But $d(c_k, c_{k+1}) = d(a_1, a_2) = 1$ and the graph isomorphism f has the property \square , which yields $f(a_1, b_2) < f(a_2, b_2)$.

The second statement of Lemma 4 can be proved analogously.

Lemma 5. *Let a graph isomorphism $f: G_1 \times G_2 \rightarrow C(P)$ have the property \square*
Let

$$f(a_1, b_1) \triangleleft f(a_2, b_2) \triangleleft \dots \triangleleft f(a_{k-1}, b_{k-1}) \triangleleft f(a_k, b_k), \quad (7)$$

$k \geq 3$ be a chain in the partially ordered set P . Then in P there exists a chain

$$f(a_1, b_1) = c_1 \triangleleft c_2 \triangleleft \dots \triangleleft c_k = f(a_k, b_k)$$

such that $c_2 = f(a_1, d)$ for some $d \in G_2$.

Proof. Let $k = 3$. In the case $a_1 = a_2$ there is nothing to prove. If $a_1 \neq a_2$, then $f(a_1, b_1) \triangleleft f(a_2, b_2)$ implies $b_1 = b_2$. Applying now the preceding Lemma we infer that

$$f(a_1, b_1) \triangleleft f(a_2, b_1) \text{ implies } f(a_1, b_3) \triangleleft f(a_2, b_3)$$

and

$$f(a_2, b_3) \triangleleft f(a_2, b_3) \text{ implies } f(a_1, b_2) \triangleleft f(a_1, b_3),$$

which means

$$f(a_1, b_1) = f(a_1, b_2) \triangleleft f(a_1, b_3) \triangleleft f(a_2, b_3).$$

Assume, as usual, for the second part of the induction that the statement of the Lemma is true for $k = n - 1$. If $k = n$, then the last three elements of the chain (7) are

$$f(a_{n-2}, b_{n-2}) \triangleleft f(a_{n-1}, b_{n-1}) \triangleleft f(a_n, b_n).$$

If $a_{n-2} = a_{n-1}$, then by the induction hypothesis there exists a chain with the required property. If $a_{n-2} \neq a_{n-1}$, then $b_{n-2} = b_{n-1}$ and analogously as in the first part of this proof it is easy to show that

$$f(a_{n-2}, b_{n-2}) = f(a_{n-2}, b_{n-1}) \triangleleft f(a_{n-2}, b_n) \triangleleft f(a_{n-1}, b_n).$$

The chain $f(a_1, b_1) \triangleleft \dots \triangleleft f(a_{n-2}, b_{n-2}) \triangleleft f(a_{n-2}, b_n)$ is of the length $n - 2$. Using the induction hypothesis, we can find a chain

$$f(a_1, b_1) = c_1 \triangleleft c_2 \triangleleft \dots \triangleleft c_{n-1} = f(a_{n-2}, b_n) \triangleleft f(a_{n-1}, b_n)$$

such that $c_2 = f(a_1, d)$ and the induction is completed.

Lemma 6. Let a graph isomorphism $f: G_1 \times G_2 \rightarrow C(P)$ have the property \square . If

$$f(a_1, b_1) \triangleleft (a_2, b_2) \triangleleft \dots \triangleleft f(a_{n-1}, b_{n-1}) \triangleleft f(a_n, b_n) = f(a_1, b_n), \quad (8)$$

$n > 1$, then $a_i = a_1$ for all i , $1 \leq i \leq n$.

Proof. We observe that $a_i \neq a_{i+1}$ for all i , $1 \leq i \leq n-1$ implies $b_i = b_{i+1}$ for all i . This means that $b_1 = b_n$, which is impossible. We conclude that there exists j , $1 \leq j \leq n-1$ such that

$$a_j = a_{j+1}. \quad (9)$$

Clearly, the statement of the Lemma is true for $n = 3$. Assume for the second part of the induction that the assertion is true for $k < n$. If $j = n-1$ (see(9)), then $a_{n-1} = a_n = a_1$. The chain

$$f(a_1, b_1) < f(a_2, b_2) < \dots < f(a_{n-1}, b_{n-1}) = f(a_1, b_{n-1})$$

has $n-1$ elements and by the induction hypothesis $a_i = a_1$ for all i , $1 \leq i \leq n-1$. Let us suppose that

$$a_j = a_{j+1}, \quad j < n-1.$$

We divide the chain (8) into two parts.

$$f(a_1, b_1) \triangleleft \dots \triangleleft f(a_j, b_j) \triangleleft f(a_{j+1}, b_{j+1}) = f(a_j, b_{j+1}), \quad (10)$$

$$f(a_{j+2}, b_{j+2}) \triangleleft \dots \triangleleft f(a_n, b_n) = f(a_1, b_n). \quad (11)$$

Applying now the preceding Lemma on the chain (10) we obtain the chain (12) with $j+1$ members

$$f(a_1, b_1) \triangleleft f(a_1, d) \triangleleft \dots \triangleleft f(a_j, b_{j+1}). \quad (12)$$

The chain

$$f(a_1, d) \triangleleft \dots \triangleleft f(a_j, b_{j+1}) \triangleleft f(a_{j+2}, b_{j+2}) \triangleleft \dots \triangleleft f(a_n, b_n) = f(a_1, b_n), \quad (13)$$

which we obtain from (12) by omitting the least element and from (11), has $n-1$ elements. Using the induction hypothesis we have $a_1 = a_j = a_{j+2} = \dots = a_{n-1}$. The chain (10) has $j+1$ elements ($j+1 < n$) and $a_j = a_1$, hence by the inductin hypothesis $a_1 = a_2 = a_3 \dots = a_{j-1}$.

Lemma 7. Let a graph isomorphism $f: G_1 \times G_2 \rightarrow C(P)$ have the property \square , let P have a locally finite length, let $f(a, b) < f(c, d)$. Then there exist elements

$$z_0, z_1, \dots, z_j \in G_2, \quad z_{j+1}, \dots, z_n \in G_1$$

such that

$$f(a, b) = f(a, z_0) \triangleleft f(a, z_1) \triangleleft \dots \triangleleft f(a, z_j) \triangleleft f(z_{j+1}, d) \triangleleft \dots \triangleleft f(z_n, d) = f(c, d).$$

Proof. We proceed by induction through $n - d(f(a, b), f(c, d))$. If $n = 1$, the result is obvious. For the second part of the induction assume that the statement of Lemma 5 holds for $d(f(a, b), f(c, d)) = n - 1$. If $d(f(a, b), f(c, d)) = n$, then there exists a chain

$$f(a, b) - f(x_0, y_0) \triangleleft f(x_1, y_1) \triangleleft \dots \triangleleft f(x_n, y_n) = f(c, d) \quad (14)$$

1. If $a = x_1$, then using the induction hypothesis we infer that there exist elements $z_0, \dots, z_j \in G_2, z_{j+1}, \dots, z_{n-1} \in G_1$ such that

$$f(a, y_1) = f(a, z_0) \triangleleft \dots \triangleleft f(a, z_j) \triangleleft f(z_{j+1}, d) \triangleleft \dots \triangleleft f(z_{n-1}, d) = f(c, d).$$

If we denote the elements of this chain by c_0, \dots, c_{n-1} , then $\{f(a, b), c_1, \dots, c_{n-1}\}$ is a chain with the required property.

2. If $a \neq x_1$, and $x_i \neq x_{i+1}$ for all $i, 0 \leq i \leq n - 1$, then $y_i = y_{i+1}$ for all $i, 0 \leq i \leq n - 1$, hence $b = y_0 = y_1 = \dots = d$. If we denote b as z_0 and x_i 's as z_i for all $i, 1 \leq i \leq n$, then the chain (14) has the required property.

If $a \neq x_1$ and there exists $i, 1 \leq i \leq n - 1$ such that $x_i = x_{i+1}$, then we divide the chain (14) into the chains

$$f(a, b) = f(x_0, y_0) \triangleleft \dots \triangleleft f(x_i, y_i) \triangleleft f(x_i, y_{i+1}) \quad (15)$$

$$f(x_{i+2}, y_{i+2}) \triangleleft \dots \triangleleft f(x_n, y_n) = f(c, d) \quad (16)$$

By Lemma 5 there exists a chain

$$f(a, b) \triangleleft f(a, t) \triangleleft \dots \triangleleft f(x_i, y_{i+1}), \quad (17)$$

which has the length $i + 1$. The elements of the chain (16) and (17) build a chain of the length n between $f(a, b)$ and $f(c, d)$ and as in the part 1 of this proof we can find a chain between $f(a, b), f(c, d)$ with the required property.

Theorem 2. ([1]). Let (M, \leq) be a quasiordered set. If Θ_1, Θ_2 are equivalences on the set M such that

- (i) $\Theta_1 \cap \Theta_2 = \omega$, where ω is the least equivalence on M ,
- (ii) $\Theta_1 \cup \Theta_2 = \iota$, where ι is the greatest equivalence on M ,
- (iii) $\Theta_1 \circ \Theta_2 = \Theta_2 \circ \Theta_1$.
- (iv) $c_1 \Theta_1 d_1, c_2 \Theta_2 d_2, d_1 \Theta_1 d_2, i \neq j, c_1 \leq c_2$ implies $d_1 \leq d_2$,

then (M, \leq) is isomorphic to the direct product of the quasiordered sets M/Θ_1 and M/Θ_2 . ($[a_1]_{\Theta_1} \leq [a_2]_{\Theta_1}$ iff $b_1 \leq b_2$ for some $b_i \in M, b_i \Theta_i a_i, i = 1, 2$)

Corollary 1. If (M, \leq) is a partially ordered set or a lattice and Θ_1, Θ_2 are equivalences on the set M with the properties (i), (ii), (iii), (iv) of Theorem 2, then $M/\Theta_1, M/\Theta_2$ are partially ordered sets, respectively lattices.

Theorem 3. Let a graph isomorphism $f: G_1 \times G_2 \rightarrow C(P)$ have the property \square ,

let the partially ordered set P have a locally finite length. Then P is isomorphic to the direct product of the partially ordered sets P_1, P_2 such that $C(P_i)$ is graph isomorphic to G_i , $i = 1, 2$.

Proof. Clearly, the relations

$$\Theta_1 = \{(f(a, b), f(a, d)), a \in G_1, b, d \in G_2\},$$

$$\Theta_2 = \{(f(a, b), f(c, b)), a, c \in G_1, b \in G_2\}$$

are equivalence relations on P . It is easy to see that Θ_1, Θ_2 have the properties (i), (ii), (iii).

Let $c_1, c_2, d_1, d_2 \in P$, $c_1 = f(a, b)$, $c_2 = f(c, d)$ and $f(a, b) \leq f(c, d)$. We suppose also $c_1 \Theta_1 d_1$, $c_2 \Theta_1 d_2$ and $d_1 \Theta_2 d_2$. Then $d_1 = f(a, x)$, $d_2 = f(c, x)$ for some $x \in G_2$.

In the case $a = c$ we have $d_1 = d_2$. If $a \neq c$, then $f(a, b) < f(c, d)$ and by Lemma 7 there exist elements $z_0, \dots, z_j \in G_2, z_{j+1}, \dots, z_n \in G_1$ such that

$$f(a, b) = f(a, z_0) \triangleleft f(a, z_1) \triangleleft \dots \triangleleft f(a, z_j) \triangleleft f(z_{j+1}, d) \triangleleft \dots \triangleleft f(z_n, d) = f(c, d).$$

If $z_{j+1} \neq a$, then $z_j = d$ and $f(a, d) < f(c, d)$

If $z_{j+1} = a$, then also $f(a, d) < f(c, d)$.

Since P has a locally finite length, there exists a finite maximal chain connecting $f(a, d)$ and $f(c, d)$. By Lemma 6 the chain has the following form

$$f(a, d) = f(x_0, d) \triangleleft f(x_1, d) \triangleleft \dots \triangleleft f(x_n, d) = f(c, d)$$

Since the graph isomorphism has the property \square , we have

$$f(a, x) = f(x_0, x) \triangleleft f(x_1, x) \triangleleft \dots \triangleleft f(x_n, x) = f(c, x)$$

(see Lemma 4), hence $d_1 \leq d_2$.

Analogously the second condition of (iv) of Theorem 2 can be proved.

If we denote P/Θ , as P_i , then from Theorem 2 and Corollary 1 it follows that P is isomorphic to $P_1 \times P_2$, P_1, P_2 are partially ordered sets.

Let $a_1 \in G_1, b_0 \in G_2$ be some fixed elements. We show that the mapping

$$g: C(P_1) \rightarrow C(\{f(a, b_0), a \in G_1\})$$

defined by $g([f(a, b)]\Theta_1) = f(a, b_0)$ is a graph isomorphism. If $[f(a, c)]\Theta_1 \triangleleft [f(b, d)]\Theta_1$, then there exist elements $x, y \in G_2$ such that $f(a, x) \triangleleft f(b, y)$ (see Theorem 2) But $a \neq b$, hence $x = y$ and $d(a, b) = 1$. From $f(a, x) \triangleleft f(b, x)$ and from the property \square of the graph isomorphism f it follows that $f(a, b_0) \triangleleft f(b, b_0)$ (see Lemma 4).

If $f(a, b_0) \triangleleft f(b, b_0)$, then $[f(a, b_0)]\Theta < [f(b, b_0)]\Theta_1$. It is easy to see that $C(\{f(a, b_0), a \in G_1\})$ is graph isomorphic to G_1 . Hence $C(P/\Theta_1)$ is graph isomorphic to G_1 .

Analogously $C(P/\Theta_2)$ and G_2 are graph isomorphic.

Note that the mapping g is also an isomorphism of partially ordered sets P_1 and $\{f(a, b_0), a \in G_1\}$.

Corollary 2. *If the supposition of the preceding Theorem are fulfilled and $a_0 \in G_1, b_0 \in G_2$, then*

$$A_1 = \{f(a, b_0), a \in G_1\}, \quad A_2 = \{f(a_0, b), b \in G_2\}$$

are partially ordered sets and A_i are isomorphic to P_i .

Theorem 4. *If $f: G_1 \times G_2 \rightarrow C(P)$ is a graph isomorphism and P is a graded partially ordered set (P is a lattice of a locally finite length), then there exist graded partially ordered sets P_1, P_2 (lattices P_1, P_2 of a locally finite length) such that P is isomorphic to $P_1 \times P_2$ and G_i is graph isomorphic to $C(P_i), i = 1, 2$.*

Proof. The statement of the Theorem is an easy consequence of Theorem 3, Corollary 1, Lemma 3, respectively Lemma 2.

Theorem 5. *Let a lattice H be a direct product of lattices A_1, A_2 and let there exist a graph isomorphism $f: C(H) \rightarrow C(L)$, where L is a lattice of a locally finite length. Then L is a direct product of lattices B_1, B_2 such $C(A_i)$ is graph isomorphic to $C(B_i), i = 1, 2$.*

Proof. Note that $C(H) = C(A_1) \times C(A_2)$ and apply the preceding Theorem.

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Received December 7, 1978

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ОБ ОРИЕНТИРОВАНИИ ПРЯМОГО ПРОИЗВЕДЕНИЯ ГРАФОВ

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Резюме

Граф, вершины которого являются элементами частично упорядоченного множества P и ребра суть те пары $\{a, b\}, a, b \in P$, где a покрывает b или b покрывает a , называется покрывающим графом $C(P)$ частично упорядоченного множества P . Пусть P является решеткой локально конечной длины, или P является частично упорядоченным множеством удовлетворяющим условию Дедекинда, тогда верно следующее утверждение: Если G_1, G_2 графы $i f: G_1 \times G_2 \rightarrow C(P)$ изоморфизм графов, то $P = P_1 \times P_2$, P_1, P_2 являются частично упорядоченными множествами и графы G_i и $C(P)$ суть изоморфны, $i = 1, 2$.