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NOTE ON CANONIZING ORDERING THEOREMS
FOR HALES JEWETT STRUCTURES

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This note presents a continuation of the research at the Winterschool 1982, see [1] .

Definition

Let (A, \leq) be a finite totally ordered set. A lexicographic tree is a set T of intervals of (A, \leq) satisfying the following rules:

- (L1) $A \in T$
- (L2) for all intervals I and J in T it follows that $I \cap J = \emptyset$ or $I \subseteq J$ or $J \subseteq I$.
- (L3) for every interval I in T which contains at least two elements there exist mutually disjoint subintervals I_0, \dots, I_ℓ in T , where $\ell \geq 1$, such that

$$I = I_0 \cup \dots \cup I_\ell .$$

Convention

Let T be a lexicographic tree on (A, \leq) . For every interval $I \in T$ let $i \text{ Succ}(I) \subseteq T$ denote the set of immediate successors of I .

Denote by $T^* = \{I \in T; i \text{ Succ}(I) \neq \emptyset\}$.

Definition

- (1) A quasi - ordering \leq_q is a transitive relation such that each two elements are comparable, but possibly $I <_q J$ and $J <_q I$ for different elements I and J .

(2) Let T be a lexicographic tree on (A, \leq) . We say that a quasi-ordering \leq_q on T^* extends the tree T iff $J \subseteq I$ implies that $I <_q J$, but not $J <_q I$.

Remark

If \leq_q is a quasi-ordering on T^* , then \leq_q induces an equivalence relation \approx_q on T^* by $I \approx_q J$ iff $I \leq_q J$ and $J \leq_q I$. The quasi-ordering \leq_q acts as a total ordering on the equivalence classes.

Notation

If \leq_q is a quasi-ordering on T^* , let $T(0), T(1), T(2), \dots$ be a monotonous (with respect to \leq_q) enumeration of the equivalence classes of \approx_q . Thus $T(0)$ is the least equivalence class and so forth. Recall that $T(0) = A$.

Definition (ordering-scheme)

Let A be a finite set. A 3-tuple $F = (\leq, T, \leq_q)$ is an ordering scheme for A iff

- (1) \leq is a total ordering on A ,
- (2) T is a lexicographic tree for (A, \leq) ,
- (3) \leq_q is a quasi-ordering on T^* which extends the tree T .

Next we show how an ordering scheme F on A can be used in order to define a total ordering on the set A^m of m -tuples over the set A .

Notation

Let $\vec{x} \in A^m$ and let X be a set of subsets of A . Then $\vec{x} \upharpoonright X$ denotes the maximal subword of \vec{x} consisting of all entries x_v which belong to some member of X .

Notation

Let $\vec{x} = (x_0, \dots, x_{m-1}) \in A^m$ and let $P = \{I_0, \dots, I_{\ell-1}\}$ be a partition of A into mutually disjoint and nonempty subsets. Then

$$\vec{x}_{/P} = (\hat{x}_0, \dots, \hat{x}_{m-1}) \in P^m$$

denotes the factorization of \vec{x} with respect to P , i.e. $\hat{x}_v = I$ iff $x_v \in I$ for every $I \in P$ and $v \in \{0, \dots, m-1\}$.

Notation

(1) if \leq_q is a quasi-ordering extending the lexicographic tree T then

$$i \text{ Succ}(T(s)) = \cup \{i \text{ Succ}(I) | I \in T(s)\}$$

denotes the union of the sets of immediate successors of intervals in $T(s)$.

As the intervals in T do not overlap (cf. (L2)) the ordering \leq on A can be extended to $i \text{ Succ}(T(s))$.

(2) By \leq_L we denote the lexicographic ordering on $(i \text{ Succ}(T(s)))^m$ with respect to the extension of \leq to $i \text{ Succ}(T(s))$.

Definition

Let A be a finite set and let $F = (\leq, T, \leq_q)$ be an ordering scheme for A . The ordering \leq_F on A^m , where m is a positive integer, is defined in the following way:

Let \vec{x} and \vec{y} be two different m -tuples in A^m .

$$\vec{x} \leq_F \vec{y} \text{ iff}$$

there exists a nonnegative integer s such that

$$\vec{x} \upharpoonright T(s)_{/i \text{ Succ}(T(s))} \leq_L \vec{y} \upharpoonright T(s)_{/i \text{ Succ}(T(s))}$$

and for every $i < s$ it follows that

$$\vec{x} \upharpoonright T(i)_{/i \text{ Succ}(T(i))} = \vec{y} \upharpoonright T(i)_{/i \text{ Succ}(T(i))}.$$

\leq_F is also called a canonical ordering of A^m (induced by the ordering scheme F).

We have the following results which solve the problem of the canonical set of

orderings of cubes stated in [1] :

Theorem 1 (irredundancy of \leq_F) :

Let A be a finite set. Then there exists a positive integer $m = m(A)$ with the following property:

For every pair F and F' of different ordering schemata for A there exist m -tuples \vec{x} and \vec{y} in A^m such that

$$\vec{x} <_F \vec{y} \quad \text{and} \quad \vec{y} <_{F'} \vec{x}$$

Theorem 2 (necessity of \leq_F) :

Let F be an ordering scheme for A . Let $m \leq n$ be positive integers. Then there exists an ordering \leq on A^n such that for every $f \in [A]_{\binom{n}{m}}$ holds:

$$\vec{x} \leq_F \vec{y} \quad \text{iff} \quad f \cdot \vec{x} \leq f \cdot \vec{y}$$

for all m -tuples \vec{x} and \vec{y} in A^m .

Theorem 3 (sufficiency of \leq_F) :

Let A be a finite set. For every nonnegative integer m there exists a nonnegative integer n with the following property:

for every total ordering on A^n there exists an m -parameter word $f \in [A]_{\binom{n}{m}}$ and there exists an ordering scheme F for A such that

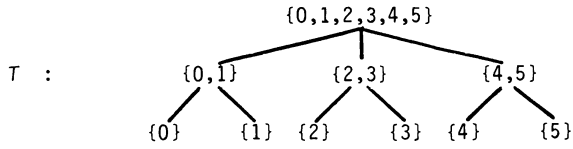
$$\vec{x} \leq_F \vec{y} \quad \text{iff} \quad f \cdot \vec{x} \leq f \cdot \vec{y}$$

for all m -tuples \vec{x} and \vec{y} in A^m .

It is interesting to note that the relatively complicated structure of unavoidable orderings of A^n appears only for larger alphabets A . For small alphabets the canonical set of orderings was determined earlier in [1] and it represents simpler results.

The first non-trivial example appears at $A = \{0,1,2,3,4,5\}$ and can be depicted

as follows:



$\leq :$ $0 < 1 < 2 < 3 < 4 < 5$

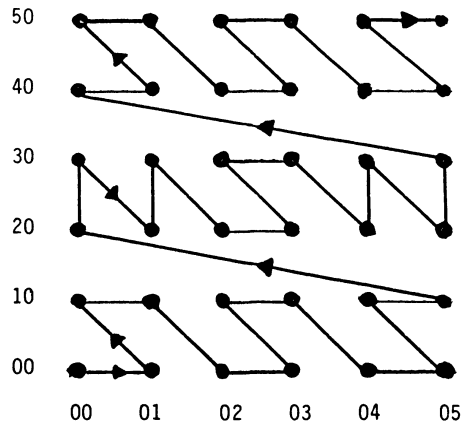
$\leq_q :$ $\{0,1\} \approx \{4,5\} <_q \{2,3\}$

(all other pairs are determined as $<_q$ is an extension of T) .

Consequently: $T(0) = \{0,1,2,3,4,5\}$

$T(1) = \{0,1,4,5\}$, $T(2) = \{2,3\}$.

The structure of the standard ordering \leq_F may be indicated by the following pairs of A^2 :



The proofs of the above theorems will appear elsewhere.

References

- [1] J. Nešetřil, H.J. Prömel, V. Rödl and B. Voigt: A canonical ordering theorem, a first attempt. In: "Proceedings of the 10th Winterschool", Supplemento ai Rendiconti del Circolo Matematico di Palermo, Serie II numero 2(1982), 193 - 197 .

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