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## SOME RESULTS CONCERNING RECONSTRUCTION CONJECTURE

Václav Nýdl

**Abstract:** B.Manvel first showed that, for every  $k$ , there exist two finite nonisomorphic graphs with the same collections of  $k$ -point subgraphs. Here, we give some new results concerning Manvel's observation. We find the bounds of reconstructibility and nonreconstructibility of graphs from subgraphs for some classes of graphs /all graphs, all trees, all equivalences/.

### 0. Introduction

We consider finite undirected graphs without loops and multiple edges. More precisely: for a set  $X$  we denote  $P_2(X)$  the set of all 2-point subsets of  $X$ ; a graph is a couple  $G = \langle V(G), E(G) \rangle$ , where  $V(G)$  is a finite set and  $E(G) \subseteq P_2(V(G))$ .

A mapping  $f: V(G) \rightarrow V(H)$  is called the homomorphism from the graph  $G$  into the graph  $H$  if for every  $Z \in V(G)$   $f(Z) \in V(H)$ , and is called the isomorphism if  $f$  is a bijection and for every  $Z$   $f(Z) \in V(H)$  if and only if  $Z \in V(G)$ . We write  $G \cong H$  to indicate isomorphic graphs.

For every subset  $Y$  of the set  $V(G)$  of the graph  $G$  the induced graph  $G/Y = \langle Y, V(G) \cap P_2(Y) \rangle$  is defined. The number of induced graphs of the graph  $G$  isomorphic to the graph  $H$  is called the frequency of  $H$  in  $G$  and denoted by  $\text{frq}(H, G)$ .

We use homomorphisms of some special types. A homomorphism  $f: G \rightarrow H$  is called the monomorphism if  $f: G \rightarrow H/f(V(G))$  is an isomorphism and is called the semimonomorphism if for every component of connectivity  $C$  of the graph  $G$   $f: G/C \rightarrow H/f(V(C))$  is an isomorphism. A homomorphism  $f: G \rightarrow H$  is said to be covering if  $f(V(G)) = V(H)$ . It is obvious that every covering monomorphism has to be an isomorphism.

The number of components of connectivity of the graph  $G$  will be denoted by  $\text{cp}(G)$ . It is obvious that a semimonomorphism  $f: G \rightarrow H$  is a monomorphism if and only if  $\text{cp}(G) = \text{cp}(H/f(V(G)))$ .

We use some integral-valued functions:

card X ... denotes the number of elements of the set X,  
 $|G|$  ... = card(V(G)) for the graph G,  
 mono(G,H) ... denotes the number of monomorphisms from G into H,  
 semi(G,H) ... denotes the number of semimonomorphisms from G into H,  
 cov(G,H) ... denotes the number of covering semimonomorphisms  
 from G into H,  
 aut(G) ... denotes the number of automorphisms of the graph G  
 /we use the identity  $\text{mono}(G,H) = \text{frq}(G,H) \cdot \text{aut } G/$ .

### 1. The frequency and the similarity of graphs

**Definition 1.1.** Let G,H be two graphs such that  $|G| = |H|$ , let k be an integer. The graphs G,H are called k-similar / $\leq k$ -similar,  $\leq k$ -c-similar, respectively/ if for every graph R such that  $|R| = k$  /  $|R| \leq k$ ,  $|R| \leq k$  & R is connected, respectively/  $\text{frq}(R,G) = \text{frq}(R,H)$  holds. We use the notation  $G \overset{k}{\sim} H$  /  $G \overset{\leq k}{\sim} H$ ,  $G \overset{\leq k}{\sim}_c H$ , respectively/.

**Corollary 1.2.** /Kelly's lemma/. For any two graphs G,H and any integer k, if  $G \overset{k}{\sim} H$ , then  $G \overset{\leq k}{\sim} H$ .

Proof. See [1] pp. 229-230.

**Corollary 1.3.** /Reconstruction conjecture/. It is conjectured that for any two graphs G,H such that  $n = |G| = |H| \geq 2$  the implication "if  $G \overset{n-1}{\sim} H$ , then  $G \cong H$ " is true.

Now, we describe some "counting" rules for frequencies.

**Lemma 1.4.** If  $G \overset{\leq k}{\sim}_c H$ , then for every R such that  $|R| \leq k$   $\text{semi}(R,G) = \text{semi}(R,H)$ .

Proof. The equality follows immediately from the observation that  $\text{semi}(R,G) = \sum_{m=1}^{\text{cp}(R)} \text{mono}(C_m, G)$ , where  $C_m$  are the components of R, from the identity  $\text{mono}(C_m, G) = \text{frq}(C_m, G) \cdot \text{aut}(C_m)$  and from their analogues for the graph H.

**Lemma 1.5.** Let  $I = I_1 \cup I_2 \cup \dots \cup I_m \cup \dots$  be a set and let  $\{R_i, i \in I\}$  be a collection of graphs such that:

1/ for every m, if  $i \in I_m$ , then  $\text{cp}(R_i) = m$ ,

2/ for every graph R there is one and only one  $i \in I$  such that

$R_i \cong R$ .

Then for any two graphs R,G the identity  $\text{semi}(R,G) = \sum_{i \in I} \text{cov}(R, R_i) \cdot \text{frq}(R_i, G)$  holds.

Proof. Let  $P = P_2(V(G))$ . For  $Z \in P$  let  $\varphi(Z) = i$  so that  $G/Z \cong R_i$ . Obviously  $\text{frq}(R_i, G) = \text{card}(\varphi^{-1}(i))$ . And now we can write  $\text{semi}(R,G) = \sum_{Z \in P} \text{cov}(R, G/Z) = \sum_{Z \in P} \text{cov}(R, R_{\varphi(Z)}) = \sum_{i \in I} \text{cov}(R, R_i) \cdot \text{frq}(R_i, G)$ .

**Lemma 1.6.** If  $G, H$  are two graphs such that  $G \stackrel{\leq k}{\sim}_C H$ , then  $G \stackrel{\leq k}{\sim} H$ .

**Proof.** Let  $I, I_m, R_1$  be the same as in Lemma 1.5. We prove by induction that for every  $j \leq k$  the proposition  $A(j)$ : "if  $q \in I_j$  and  $|R_q| \leq k$ , then  $\text{frq}(R_q, G) = \text{frq}(R_q, H)$ " is true.

1/  $A(1)$  is true because of assumption  $G \stackrel{\leq k}{\sim}_C H$ .

2/  $A(1), A(2), \dots, A(j-1)$  are supposed to be true. We introduce  $Q_G = \sum_{i \in I_1 \cup I_2 \cup \dots \cup I_{j-1}} \text{cov}(R_q, R_1) \cdot \text{frq}(R_1, G)$ , and analogically  $Q_H$ .

Using Lemma 1.5. we obtain  $\text{semi}(R_q, G) = Q_G + \text{aut}(R_q) \cdot \text{frq}(R_q, G)$  and  $\text{semi}(R_q, H) = Q_H + \text{aut}(R_q) \cdot \text{frq}(R_q, H)$ . But  $Q_G = Q_H$  because for every  $i \in I_1 \cup I_2 \cup \dots \cup I_{j-1}$   $\text{frq}(R_1, G) = \text{frq}(R_1, H)$  /if we suppose  $|R| \leq k/$ . Moreover,  $\text{semi}(R_q, G) = \text{semi}(R_q, H)$  according to Lemma 1.4. Thus, we have  $\text{frq}(R_q, G) = [\text{semi}(R_q, G) - Q_G] / \text{aut}(R_q) = [\text{semi}(R_q, H) - Q_H] / \text{aut}(R_q) = \text{frq}(R_q, H)$ .

**Theorem 1.7.** For any two graphs  $G, H$  and for any integer  $k$ , the following three properties are equivalent

1/  $G \stackrel{k}{\sim} H$ , 2/  $G \stackrel{\leq k}{\sim}_C H$ , 3/  $G \stackrel{\leq k}{\sim} H$ .

**Proof.** The theorem is the summary of Corollary 1.2. and

Lemma 1.6.

**Corollary 1.8.** The reconstruction conjecture is true for disconnected graphs.

**Proof.** Let  $G, H$  be two disconnected graphs such that  $n = |G| = |H| > 2$  and let  $G \stackrel{n-1}{\sim} H$ . Using Theorem 1.7. we get  $G \stackrel{(n-1)}{\sim}_C H$  and, since  $G, H$  are disconnected, even  $G \stackrel{\leq n}{\sim}_C H$ . Now, by Theorem 1.7.,  $G \stackrel{n}{\sim} H$ , i.e.  $G \simeq H$ .

## 2. Bounds of reconstructibility and nonreconstructibility

Let  $N$  be the set of all natural numbers. For every subset  $M$  of  $N$  we define  $\max M = +\infty$ . Let us denote  $N^* = N \cup \{+\infty\}$ .

**Definition 2.1.** Let  $\mathcal{F}$  be a subclass of the class of all finite graphs. We define the mapping  $u_{\mathcal{F}}: N \rightarrow N^*$  as  $u_{\mathcal{F}}(n) = \max \{m; (\forall F_1, F_2 \in \mathcal{F}) (|F_1| = |F_2| \leq m \& F_1 \stackrel{n}{\sim} F_2) \Rightarrow F_1 \simeq F_2\}$ .

**Corollary 2.2.** We denote  $\mathcal{G}$  the class of all finite graphs. B. Manvel showed in [2] that for every  $n \in N$  the inequality  $u_{\mathcal{G}}(n) < +\infty$  holds. Further, the reconstruction conjecture can be written in the form  $u_{\mathcal{G}}(n) \geq n+1$  for  $n \geq 2$ .

**Proposition 2.3.** Let  $\mathcal{T}$  be the class of all finite trees. Then, for every  $n > 1$ ,  $n+1 \leq u_{\mathcal{T}}(n) < 2n$ .

**Proof.** The first inequality expresses the fact that the reconstruction conjecture is true for the case of trees. The second

one was proved in [5] where, for every  $n > 1$ , we constructed two nonisomorphic trees  $T_1, T_2$  having  $2n$  elements such that  $T_1 \overset{n}{\sim} T_2$ .

**Proposition 2.4.** If  $\mathcal{G}$  is the class of all finite graphs, then, for every  $n > 1$ , the inequality  $u_{\mathcal{G}}(n) < \min(2n, 3n/2 + 15/2)$  holds.

**Proof.** To prove the inequality we use Proposition 2.3. and the construction from [5] where, for every  $k \geq 2$ , we constructed two nonisomorphic graphs  $G_1, G_2$  having  $3k + 6$  elements such that  $G_1 \overset{2k}{\sim} G_2$ .

**Corollary 2.5.** V.Müller in [3] showed that for every  $\rho$ ,  $1 < \rho < 2$ , there exist a class  $\mathcal{R}$  and a number  $n_{\rho}$  such that for every  $n \in \mathbb{N}$   $u_{\mathcal{R}}(n) > \rho \cdot n$  and moreover,  $\mathcal{R}$  contains asymptotically the most graphs on  $n$  elements.

**Remark 2.6.** It was proved in [4] that for every  $n \in \mathbb{N}$  in the class  $\mathcal{E}$  of all finite equivalences the inequalities

$n \cdot (\ln n - 1) \leq u_{\mathcal{E}}(n) < (n+1) \cdot 2^{n-1}$  hold /here  $\ln$  denotes the logarithmus naturalis/.

**Problem 2.7.** Prove that, for every sufficiently "rich" class  $\mathcal{F}$  of finite graphs, the inequality  $u_{\mathcal{F}}(n) < +\infty$  holds for every  $n \in \mathbb{N}$ .

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