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A Note on Reconstructing of Finite Trees from Small Subtrees

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Some basic observations on the reconstruction of a finite tree from its subtrees are given. It is conjectured that every tree with n vertices is determined by the collection formed of all its subtrees with the number of vertices $q = 1, 2, \dots, k$ provided k is greater than $n/2$.

1. Preliminaries

All graphs considered are finite, simple and undirected. If $G = (X, E)$ is a graph and Y is a subset of the set of vertices X then G/Y denotes the induced graph with the set of vertices Y and the set of edges formed of all edges from E contained in Y . We use the symbol \cong to denote the isomorphism of graphs.

For every two graphs H, G we define the frequency $frq(H, G)$ as the number of induced graphs of G isomorphic to H . Four types of similarity can be defined.

Definition. Let \mathcal{C} be a class of graphs and let k be an integer. For two graphs G_1, G_2 we define $G_1 \sim^k G_2$ ($G_1 \sim^{\leq k} G_2$, respectively) iff $frq(H, G_1) = frq(H, G_2)$ for every graph H on k (on $\leq k$, respectively) vertices, $G_1 \sim_{\mathcal{C}}^k G_2$ ($G_1 \sim_{\mathcal{C}}^{\leq k} G_2$, respectively) iff $frq(H, G_1) = frq(H, G_2)$ for every graph H from \mathcal{C} having k (having $\leq k$, respectively) vertices.

2. Reconstructing trees

In the reconstruction theory we are interested in the implication " $G_1 \sim G_2 \Rightarrow G_1 \cong G_2$ " where \sim denotes some of the similarity types. There are some positive answers in the case of trees (the class of all trees will be denoted by \mathcal{T}).

Kelly [4] proved in 1957 that the implication holds for trees G_1, G_2 on n vertices in the case of similarity \sim^{n-1} , i.e. that trees are reconstructible from one vertex deleted subgraphs. Further, Harary and Palmer [3] proved in 1966 that the implica-

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tion holds for trees G_1, G_2 on n vertices in the case of similarity $\sim_{\mathcal{F}}^{n-1}$, i.e. that trees are reconstructible from endvertex deleted subtrees. In 1976, Giles [1] showed that the implication holds when G_1, G_2 are trees on n vertices ($n > 4$) and the type of similarity is $\sim_{\mathcal{F}}^{n-2}$, i.e. that trees are reconstructible from two-vertex deleted subtrees (with a trivial exception). He also studied k -vertex deleted subtrees reconstruction.

The reconstruction of a graph from "small" subgraphs was first investigated by Manvel [5]. The following theorem we regard as basic.

Theorem 1. *Let k, n be integers, T_1, T_2 trees on n vertices. The following three properties are equivalent*

- (i) $T_1 \sim^k T_2$,
- (ii) $T_1 \sim^{\geq k} T_2$,
- (iii) $T_1 \sim^{\leq k} T_2$.

Proof. Apply Theorem 1.7 from [8] to the class of trees.

One may try to add the fourth property to Theorem 1, namely the property (iv) $T_1 \sim^k T_2$. We are going to show (in Theorem 2) that it is not possible.

Let n_1, \dots, n_s be integers and n their sum. We define the "star" graph $St(n_1, \dots, n_s)$ as a tree consisting of s paths of lengths n_1, \dots, n_s all "emanating" from one common vertex. Obviously, this tree has $n + 1$ vertices.

Theorem 2. *Let $k \geq 5$ be an integer. Then for every $n \geq 2k - 2$ there exist two trees T_1, T_2 on n vertices such that $T_1 \sim^k T_2$ but not $T_1 \sim^{\leq k} T_2$.*

Proof. Take $q = n - k + 1$, $T = St(1, k - 4, q)$ and suppose $T = (X, E)$ where $X = \{x_1, \dots, x_{n-2}\} \cup \{a\}$, $E = \{\{x_i, x_{i+1}\}; i = 1, \dots, n - 3\} \cup \{a, x_{q+1}\}$. Now, for $x, y \notin X$ let $X_1 = X \cup \{x\}$, $X_2 = X \cup \{y\}$, $E_1 = E \cup \{x, x_2\}$, $E_2 = E \cup \{x_{n-2}, y\}$ and finally $T_j = (X_j, E_j)$ for $j = 1, 2$. The trees T, T_1, T_2 are shown in Figure 1 for the case $k = 5, n = 8$.

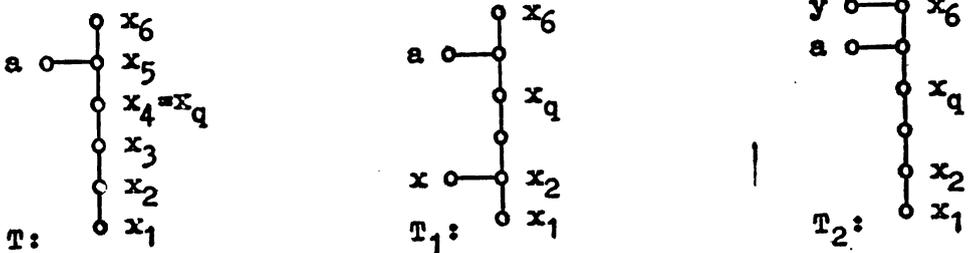


Fig. 1.

To prove the required properties we put $M_j = \{K \subset X_j; T_j/K \text{ is a tree and card } K = k\}$ and we define $f: M_1 \rightarrow M_2$ as follows:

- (1) if $x \notin K$ then $f(K) = K$,
- (2) if $x \in K$ et $x_1 \notin K$ then $f(K) = \{x_{q-1}, \dots, x_{n-2}\} \cup \{y\}$,
- (3) if $x \in K$ et $x_1 \in K$ then $f(K) = \{x_q, \dots, x_{n-2}\} \cup \{y, a\}$.

It can be easily shown that f is a bijection and for every $K \in M_1$ it is true that $T_1/K \cong T_2/f(K)$. Thus, we have $T_1 \sim_{\mathcal{F}}^k T_2$. On the other hand, since T_1 has just two vertices of degree 3 while T_2 has only one such a vertex, we have $frq(St(1, 1, 1), T_1) = 2$ but $frq(St(1, 1, 1), T_2) = 1$ and thus $T_1 \not\sim_{\mathcal{F}}^k T_2$ does not hold.

The following theorem shows that the diameter of a tree cannot be reconstructed from “small” subtrees.

Theorem 3. *For every $k \geq 1$ there exist two trees T_1, T_2 on $3k + 1$ vertices such that $T_1 \sim_{\mathcal{F}}^k T_2$ and $\text{diam } T_1 \neq \text{diam } T_2$.*

Proof. Let $T = St(k - 1, k, k) = (X, E)$ where $X = \{x_1, \dots, x_{2k}\} \cup \{y_1, \dots, y_k\}$ and $E = \{\{x_i, x_{i+1}\}; i = 1, \dots, 2k - 1\} \cup \{\{y_i, y_{i+1}\}; i = 1, \dots, k - 1\} \cup \{x_k, y_1\}$. Now, let $T_1 = St(k, k, k) = (X_1, E_1)$ and $T_2 = St(k - 1, k, k + 1) = (X_2, E_2)$ where $X_1 = X \cup \{x\}$, $X_2 = X \cup \{w\}$, $E_1 = E \cup \{x, x_1\}$, $E_2 = E \cup \{x_{2k}, w\}$.

We take $M_j = \{K \subset X_j; \text{card } K \leq k \text{ and } T_j/K \text{ is a tree}\}$ and we define $f: M_1 \rightarrow M_2$ as follows:

- (1) if $x \notin K$ then $f(K) = K$,
- (2) if $x \in K$ then $K = \{x_1, \dots, x_r\} \cup \{x\}$, where $r < k$, and then we put $f(K) = \{x_{2k-r+1}, \dots, x_{2k}\} \cup \{w\}$.

It is clear that for every $K \in M_1$ we have $T_1/K \cong T_2/f(K)$ but, on the other hand, $\text{diam } T_1 = 2k \neq 2k + 1 = \text{diam } T_2$.

3. Bounds of reconstructability

To analyse the problem in more details we define the concept of a function rec .

Definition. Let \mathcal{F} be a class of graphs. The integer valued function $rec_{\mathcal{F}}$ is defined as $rec_{\mathcal{F}}(n) = \min \{k; \text{ for every two graphs } G_1, G_2 \text{ on } n \text{ vertices from the class } \mathcal{F} \text{ it is true that } G_1 \sim_{\mathcal{F}}^{\leq k} G_2 \Rightarrow G_1 \cong G_2\}$.

Müller [6] proved that there exists a class \mathcal{H} containing asymptotically the most of graphs with n vertices (in the sense of limit) such that for every $r > 1/2$ there exists n_r such that for $n > n_r$ the inequality $rec_{\mathcal{H}}(n) < r \cdot n$ holds. On the other hand, we have proved (see [7]) that for the class \mathcal{G} of all graphs and for every $n \geq 15$ $rec_{\mathcal{G}}(n) > (2n/3) - 6$.

Since Müller's class \mathcal{H} does not contain the class \mathcal{T} of all trees nothing can be derived for $rec_{\mathcal{T}}$.

Theorem 4. For every $n \geq 6$ $rec_{\mathcal{T}}(n) > n/2$.

Proof. It was shown in [7] that for $T_1 = St(k - 1, k - 1, 1)$, $T_2 = St(k - 2, k, 1)$ with $2k$ vertices ($k \geq 3$) it is true that $T_1 \sim^k T_2$. Obviously T_1, T_2 are nonisomorphic.

Checking all the cases of trees with ≤ 10 vertices (according [2]) we obtained the exact values of $rec_{\mathcal{T}}(n)$ (see Table 1). Taking into account this result and supposing that our construction from Theorem 4 is “the best possible” we formulate:

Table 1 $t(n)$ denotes the number of nonisomorphic trees on n vertices)

n	4	5	6	7	8	9	10
$t(n)$	2	3	6	11	23	47	106
$rec_{\mathcal{T}}(n)$	3	3	4	4	5	5	6

Conjecture. For every $n \geq 4$ $rec_{\mathcal{T}}(n) = \lfloor n/2 \rfloor + 1$, where $\lfloor \]$ denotes the integral part of a number.

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