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# A NOTE ON ARGUESIAN LATTICES

by ALAN DAY\*) (Received May 5, 1981)

#### §0 INTRODUCTION

In [4] Jónsson introduced the Arguesian lattice identity which reflects precisely Desargues' Theorem for projective spaces qua lattices. In this note we "geometrize" Jónsson's equation by proving that one can assume more geometrical facts about the lattice variables in the equational or the (equivalent) implicational form. An interesting consequence of this is the fact that 2-distributive (modular) lattices are Arguesian. This fact reflects the property that one dimensional projective spaces geometrically trivial.

## §1 PRELIMINARIES

The Arguesian identity and its equivalents have been developed by Jónsson et al in a series of papers (see especially [3], [6] and [7]). An early result, [5], was that Arguesian lattices are modular. Therefore we will assume throughout this paper that all lattices are modular.

Let L be a (modular) lattice. A triangle (or trilateral) in L is an arbitrary triple  $a = (a_0, a_1, a_2) \in L^3$ . For two triangles, a and b, in L we require certain derived polynomials and statements

#### 1. Definition. For a, b in L

(a)  $p(a, b) = (a_0 \lor b_0) \land (a_1 \lor b_1) \land (a_2 \lor b_2)$ 

- (b)  $c_i(a, b) = (a_j \lor a_k) \land (b_j \lor b_k), \quad \{i, j, k\} = \{0, 1, 2\}$
- (c)  $CP(a, b) : (a_0 \lor b_0) \land (a_1 \lor b_1) \leq a_2 \lor b_2$
- (d)  $AP(a, b) : c_2(a, b) \leq c_0(a, b) \lor c_1(a, b)$

CP(a, b) is an abbreviation of central perspectivity for the triangles (of points in a projective plane) a and b. AP(a, b) is an abbreviation for the axial perspectivity of these triangles (of points).

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**2. Definition.** A lattice  $(L; \vee, \wedge)$  is called:

(a) Arguesian if for any pair of triangles a, b in  $L, p(a, b) \leq a_0 \vee \vee (b_0 \wedge (b_1 \vee (c_2 \wedge (c_0 \vee c_1))))$ .

(b) Desarguean if for any pair of triangles a, b in L, CP(a, b) implies AP(a, b).

(c) Arguesian (Desarguean) for a given class of triangles if (a) (respectively (b)) holds for that particular class of triangles.

The following theorem is a synopsis of the pertient results that appear in aforementioned papers.

**3. Theorem.** ([3], [6], and [7]): Let L be a modular lattice, then the following are equivalent:

(1) L is Arguesian.

(2) L is Desarguean.

(3) L is Desarguean for triangles a, b in L satisfying (A):  $a_i \lor p = b_i \lor p = a_i \lor b_i$ , i = 0, 1, 2.

(4) L is Arguesian for triangles a, b in L satisfying (A):  $a_i \lor p = b_i \lor p = a_i \lor b_i$ , i = 0, 1, 2.

The other notion we need is that of an *n*-diamond. This notion reflects the geo metrical property of n + 1 points in general position in a projective space o-dimension  $\ge n - 1$ .

4. Definition. (i) A sequence  $d = (d_0, d_1, ..., d_n)$  in a (modular) lattice L is called an *n*-diamond if

$$(D_n 1) \bigvee_{\substack{j \neq i}}^{0, n} d_j = v, \qquad i \leq n$$
$$(D_n 2) d_i \wedge \bigvee_{\substack{k \neq i \\ k \neq j}}^{0, n} d_k = u, \qquad i \neq j \leq n.$$

We will call a sequence, d in L, a lower n-diamond (upper n-diamond) if it satisfies  $(D_n 2)$  (respectively  $(D_n 1)$ ).

(ii) A lattice is called *n*-distributive if it does not contain an *n*-diamond.

By Huhn [2] and Hermann & Huhn [1] we know that (modular) *n*-distributive lattices form an equational class of (modular) lattices and that an *n*-diamond is either non-trivial  $(d_i \neq d_j \text{ for } i \neq j)$  or completely trivial (i.e.  $d_0 = d_1 = \ldots = d_n$ ).

The last notion we need is that of an *independent set*  $\{a_1, \ldots, a_n\}$  in L.

$$\perp (a_1, \ldots, a_n) \quad \text{iff} \quad u = a_i \wedge \bigvee_{j=i}^{l, n} a_j \quad \text{for all } i \leq n.$$

#### §2 THE RESULTS

Our goal is to discover what other restrictions one can place on the triangles and still ensure the Desarguean implication. Basically, we are attempting to "regeometrize" the Desarguean implication by removing all "degenerate" cases of variable substitutions. Our first result is an easy reduction to independent triangles.

**1. Lemma.** For a modular lattice L, L is Arguesian if and only if L is Desarguean for triangles a and b in L satisfying (A)  $a_i \lor p = b_i \lor p$  (=  $a_i \lor b_i$ )  $i \le 2$  and (I)  $\perp(a_0, a_1, a_2)$  and  $\perp(b_0, b_1, b_2)$ .

**Proof:** The condition is clearly necessary so let a and b be arbitrary centrally perspective triangles in L satisfying (A). By defining

$$x'_i = (x_i \lor x_j) \land (x_i \lor x_k)$$

for  $x \in \{a, b\}$  and  $\{i, j, k\} = \{0, 1, 2\}$  the reader can easily check that a' and b' are centrally perspective triangles in L that satisfy (A) and (I). By the condition we infer AP(a', b') and since  $c'_i = c_i$  for all i we obtain AP(a, b).

Later results will be proven in a similar way as we add more conditions on our triangles. Let us first note however that we may always assume that the triangles a and b are *strictly* independent. That is  $x_i$  and  $x_j$  will be incomparable if  $i \neq j$ ,  $x \in \{a, b\}$ ,

**2. Lemma.** Let a and b be centrally perspective triangles in a modular lattice L. If either  $\{a_0, a_1, a_2\}$  or  $\{b_0, b_1, b_2\}$  contains comparable elements, then a and b are axially perspective.

Proof: By symmetry we need only consider the cases  $a_0 \leq a_1$ ,  $a_0 \leq a_2$  and  $a_0 \leq a_1$ . If  $a_0 \leq a_1$  then CP(a, b) gives us

$$a_0 \leq (a_0 \lor b_0) \land (a_1 \lor b_1) \leq a_2 \lor b_2$$

and

$$c_{0} \lor c_{1} = (a_{1} \lor a_{2}) \land (b_{1} \lor b_{2} \lor c_{1}) =$$

$$= (a_{1} \lor a_{2}) \land (b_{1} \lor ((b_{2} \lor a_{0} \lor a_{2}) \land (b_{0} \lor b_{2}))) \ge$$

$$\ge (a_{1} \lor a_{2}) \land (b_{1} \lor ((a_{0} \lor b_{0}) \land (a_{1} \lor b_{1}) \land (b_{0} \lor b_{2}))) =$$

$$= (a_{1} \lor a_{2}) \land (a_{1} \lor b_{1}) \land (b_{1} \lor ((a_{0} \lor b_{0}) \land (b_{0} \lor b_{2}))) \ge$$

$$\ge a_{1} \land (b_{1} \lor b_{0}) =$$

$$= (a_{0} \lor a_{1}) \land (b_{0} \lor b_{1}) =$$

$$= c_{2}.$$

If  $a_2 \leq a_0$  then CP(a, b) gives us

$$(a_2 \lor a_1 \lor b_1) \land (a_0 \lor b_0) \leq a_2 \lor b_2$$

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and

$$c_{0} \vee c_{1} = c_{0} \vee (a_{0} \wedge (b_{0} \vee b_{2})) =$$

$$= (a_{0} \wedge (b_{0} \vee b_{2})) \vee (a_{2} \cap b_{0} \vee b_{2})) \vee c_{0} =$$

$$= (a_{0} \wedge (b_{0} \vee b_{2})) \vee \{(a_{1} \vee a_{2}) \wedge [b_{1} \vee b_{2} \vee (a_{2} \wedge (b_{0} \vee b_{2}))]\} =$$

$$= (a_{0} \wedge (b_{0} \vee b_{2})) \vee \{(a_{1} \vee a_{2}) \wedge [b_{1} \vee ((b_{2} \vee a_{2}) \wedge (b_{0} \vee b_{2}))]\} \geq$$

$$\geq (a_{0} \wedge (b_{0} \vee b_{2})) \vee \{(a_{1} \vee a_{2}) \wedge [b_{1} \vee ((a_{2} \vee a_{1} \vee b_{1}) \wedge (a_{0} \vee b_{0}) \wedge (b_{0} \vee b_{2}))]\} =$$

$$= (a_{0} \wedge (b_{0} \vee b_{2})) \vee \{(a_{1} \vee a_{2}) \wedge [b_{1} \vee b_{0} \vee (a_{0} \wedge (b_{0} \vee b_{2}))]\} =$$

$$= [a_{1} \vee a_{2} \vee (a_{0} \wedge (b_{0} \vee b_{2}))] \wedge (b_{0} \vee b_{1} \vee (a_{0} \wedge (b_{0} \vee b_{2})) \geq$$

$$\geq [a_{1} \vee ((a_{2} \vee b_{0} \vee b_{2}) \wedge a_{0})] \wedge (b_{0} \vee b_{1}) \geq$$

$$\geq [a_{1} \vee [a_{0} \vee (b_{0} \vee ((a_{0} \vee b_{0}) \wedge (a_{1} \vee b_{1})))]] \wedge (b_{0} \vee b_{1}) =$$

$$= [a_{1} \vee (a_{0} \wedge (a_{1} \vee b_{0} \vee b_{1}))] \wedge (b_{0} \vee b_{1}) =$$

$$= c_{2}.$$

The case where  $a_0 \leq a_2$  is left for the reader.

Now let a and b be centrally perspective triangles in L satisfying (A) and (I) define

and(i)  $s_{\mathbf{a}} = (a_0 \land (a_2 \lor b_2)) \lor (a_1 \land (a_2 \lor b_2))$ 

(ii) 
$$s_{b} = (b_0 \land (a_2 \lor b_2)) \lor (b_1 \land (a_2 \lor b_2))$$

(iii)  $a'_i = a_i \lor s_a$  and  $b'_i = b_i \lor s_b$ 

Easy calculations give us that a' and b' are centrally perspective triangles satisfying (A) and (I). Moreover  $c'_2 = c_2$ ,  $c'_1 = c_1 \lor (c_0 \land (a_2 \lor b_2))$  and  $c'_0 = c_0 \lor (c_1 \land (a_2 \lor b_2))$ . Therefore AP(a, b) holds if and only if AP(a', b') holds.

The above allows us to add the conditions  $a_0 \wedge (a_2 \vee b_2) = a_1 \wedge (a_2 \vee b_2) = 0_a$ and  $b_0 \wedge (a_2 \vee b_2) = b_1 \wedge (a_2 \vee b_2) = 0_b$   $(0_a = a_0 \wedge a_1 = a_0 \wedge a_2 = a_1 \wedge a_2$  and similarly for  $0_b$  to (A) and (I). We are however more interested in certain consequences of these conditions.

**3. Theorem.** A modular lattice, L, is Arguesian if and only if L is Desarguean for triangles a and b in L satisfying (A), (I) and (U):  $p = (a_i \lor b_i) \land (a_j \lor b_j), i \neq j$ .

Proof: Using the notation above the theorem statement we get

$$a'_0 \wedge (a'_2 \vee b'_2) = 0_a \vee s_a \leq p'.$$

Therefore  $p' = p' \lor (a'_0 \land (a'_2 \lor b'_2)) = (p' \lor a'_0) \land (a'_2 \lor b'_2)$  since  $CP(a', b') = (a'_0 \lor b'_0) \land (a'_2 \lor b'_2)$  by (A).

We want one more condition on our triangles. Let a, b be triangles in L satisfying (A), (I) and (U) and define

(i)  $u_{\mathbf{a}} = \bigvee_{i}^{0,2} (a_{i} \wedge p), \ u_{b} = \bigvee_{i}^{0,2} (b_{i} \wedge p),$ (ii)  $a'_{i} = a_{i} \vee u_{\mathbf{a}}, \ b'_{i} = b_{i} \vee u_{b}.$ 

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We have by easy calculations:

(1)  $a'_i \vee b'_i = a_i \vee b_i$ (2) p' = p and a', b' satisfy (U) (3)  $a'_i \vee a'_j = a_i \vee a_j \vee (a_k \wedge (a_i \vee b_i)) = a_i \vee a_j \vee (a_k \wedge (a_j \vee b_j))$ (4)  $(a'_i \vee a'_j) \wedge (a'_i \vee a'_k) = [(a_i \vee a_j) \wedge (a_i \vee a_k)] \vee u_a$ (5) a, b satisfy (I) (6)  $a'_i \vee p' = a'_i \vee b'_i = b'_i \vee p'$  hence a, b satisfy (A) (7)  $c'_i = c_i \vee (c_j \wedge (a_k \vee b_k)) = c_i \vee (c_k \wedge (a_j \vee b_j))$ (8) AP(a, b) iff AP(a', b')

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Therefore it behooves one to determine what other special properties hold for these triangles.

**4. Lemma.**  $(a'_0, a'_1, a'_2, p')$  and  $(b'_0, b'_1, b'_2, p')$  are lower 3-diamonds with bottom element  $u_a$  (resp.  $u_b$ ).

Proof:

$$\begin{aligned} a'_{0} \wedge (a'_{1} \vee a'_{2}) &= (a_{0} \vee (a_{1} \wedge p) \vee (a_{2} \wedge p)) \wedge (a_{1} \vee a_{2} \vee (a_{0} \wedge p)) = \\ &= (a_{0} \wedge (a_{1} \vee a_{2})) \vee u_{e} = \\ &= u_{e} \end{aligned}$$

$$\begin{aligned} a'_{0} \wedge (a'_{1} \vee p') &= (a_{0} \vee (a_{1} \wedge p) \vee (a_{2} \wedge p)) \wedge (a_{1} \vee p) = \\ &= u_{e} \vee (a_{0} \wedge (a_{1} \vee p)) = \\ &= u_{e} \vee (a_{0} \wedge (a_{1} \vee b_{1}) \wedge (a_{1} \vee a_{0} \vee b_{0})) = \\ &= u_{e} \vee (a_{0} \wedge (a_{1} \vee b_{1})) = \\ &= u_{e} \vee (a_{0} \wedge (a_{1} \vee b_{1})) = \\ &= u_{e} \vee (a_{0} \wedge a_{1} \vee (a_{2} \wedge p)) = \\ &= (a_{2} \wedge p) \vee [(a_{0} \vee a_{1}) \wedge (a_{0} \vee b_{0}) \wedge (a_{1} \vee b_{1})] = \\ &= u_{e} \end{aligned}$$

**5. Lemma.** For each  $\{i, j, k\} = \{0, 1, 2\}, (a'_i, b'_i, c'_j, c'_k)$  is a lower 3-diamond. Proof:

$$\begin{aligned} c'_{j} \wedge (a'_{i} \vee b'_{i}) &= (a_{i} \vee b_{i}) \wedge (c_{j} \vee (c_{k} \wedge (a_{i} \vee b_{i}))) = \\ &= [c_{j} \wedge (a_{i} \vee b_{i})] \wedge [c_{k} \wedge (a_{i} \vee b_{i})] = \\ &= c'_{k} \wedge (a'_{i} \vee b'_{i}) \end{aligned}$$

$$\begin{aligned} a'_{i} \wedge (c'_{j} \vee c'_{k}) &= [a_{i} \vee (a_{j} \wedge p) \vee (a_{k} \wedge p)] \wedge [c_{j} \vee c_{k}] = \\ &= [a_{i} \vee (a_{j} \wedge (a_{i} \vee b_{i})) \vee (a_{k} \wedge (a_{i} \vee b_{i}))] \wedge [c_{j} \vee c_{k}] = \\ &= [a_{i} \vee (b_{j} \wedge (a_{i} \vee a_{j})) \vee (b_{i} \wedge (a_{i} \vee a_{k}))] \wedge [c_{j} \vee c_{k}] = \\ &= (b_{j} \wedge (a_{i} \vee a_{j})) \vee (b_{i} \wedge (a_{i} \vee a_{k})) \vee (a_{i} \wedge (c_{j} \vee c_{k})) \end{aligned}$$

$$= (b_j \land (a_i \lor a_j)) \lor (b_i \land (a_i \lor a_k)) \lor (a_i \land (b_i \lor b_k)) \lor (a_i \land (b_i \lor b_j)) =$$
  
=  $[c_j \land (a_i \lor b_i)] \lor [c_k \land (a_i \lor b_i)]$ 

The other calculations are left to the reader.

**6. Theorem.** A modular lattice, L, is Arguesian if and only if any triangles **a**, **b** in L satisfying:

(A)  $a_i \lor p = b_i \lor p = a_i \lor b_i$ 

(I)  $\perp (a_0, a_1, a_2)$  and  $\perp (b_0, b_1, b_2)$ 

(U)  $p = (a_i \lor b_i) \land (a_i \lor b_i), i \neq j$ 

(LF)  $(a_0, a_1, a_2, p)$ ,  $(b_0, b_1, b_2, p)$ ,  $(a_i, b_i, c_j, c_k)$  are lower 3-diamonds are axially perspective. Moreover  $c_i \lor c_j = c_i \lor c_k = c_j \lor c_k$ .

7. Corollary. 2-distributive (modular) lattices are Arguesian.

Proof: If L is 2-distributive and a and b are triangles in L satisfying (A), (I), (U) and (LF) then the 3-diamond generated by the lower 3-diamond  $(a_2, b_2, c_0, c_1)$  must be trivial. This gives

$$[c_0 \land (a_2 \lor b_2)] \lor [c_1 \land (a_2 \lor b_2)] = q_{02} \land q_{12} \land a \land b$$

where

$$q_{ij} = a_i \lor a_j \lor p$$
,  $a = a_0 \lor a_1 \lor a_2$  and  $b = b_0 \lor b_1 \lor b_2$ .

Therefore

$$c_0 \lor c_1 = c_0 \lor c_1 \lor (q_{01} \land q_{02} \land a \land b) =$$
  
=  $[c_0 \lor c_1 \lor (q_{01} \land q_{02})] \land a \land b =$   
=  $[c_0 \lor (q_{02} \land (c_1 \lor q_{01}))] \land a \land b$   
=  $(c_0 \lor q_{02}) \land a \land b =$   
=  $a \land b \ge$   
 $\ge c_2.$ 

### §3 DISCUSSION

One would obviously have liked to show that we can assume  $0_a = 0_b$  and a = b. This however is impossible for as Jónsson (private communication) has observed, 2-dimensional Hall-Dilworth gluings of vector space lattices can produce both Arguesian and non-Arguesian lattices. In the non Arguesian case one needs  $0_a \neq 0_b$ .

Andras Huhn has asked if there exists a partial (modular) lattice configuration whose exclusion would characterize Arguesian lattices. The author cannot answer this question at present. Its answer should be connected with the projectivity of

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the "correct" version of the Desarguean implication. Our version seems to lack that property.

Finally the author feels that one *should* be able to add  $a_i \lor a_j \lor p = a_i \lor a_k \lor p$  to the conditions on the triangles. He (obviously) has no proof of this at this time.

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