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## GRAPHS AND BETWEENESS\*

MILAN SEKANINA

1. In this paper, a graph  $(G, \varrho)$  is always connected, undirected, without loops and multiple edges,  $G \neq \emptyset$ . Thus, if  $a, b$  are vertices of  $(G, \varrho)$ ,  $\{a, b\} \in \varrho$  exactly when  $a, b$  are connected by an edge. We often write  $G$  instead of  $(G, \varrho)$ . If  $(G, \varrho)$  is a graph and  $M \subset G$ , then  $(M, \varrho/M)$  means a full subgraph of  $(G, \varrho)$ , i.e.  $a, b \in M$ ,  $\{a, b\} \in \varrho/M$  just if  $\{a, b\} \in \varrho$ . Often  $\varrho$  is used instead of  $\varrho/M$ . Let  $(M, \varrho)$  be a full subgraph of  $(G, \varrho)$ . Then  $\mathcal{K}(M)$  means the decomposition of  $(M, \varrho)$  into connected components.  $\mu$  means the usual metric in  $(G, \varrho)$ , i.e.  $\mu(a, b)$  is the number of edges in a shortest path connecting the vertices  $a$  and  $b$ .  $\mathcal{E}(G, \varrho)$  is the system of all 2-components of the graph  $(G, \varrho)$ . Here a 2-component is a maximal full subgraph of  $(G, \varrho)$  containing for any two distinct vertices  $a, b$  belonging to it at least one circle in which  $a$  and  $b$  are lying. By [6], § 15 one easily sees that the following assertion is true.

**Proposition.** *If  $X, Y \in \mathcal{E}(G, \varrho)$ ,  $X \neq Y$ , then  $\text{card}(X \cap Y) \leq 1$ . If  $X \in \mathcal{E}(G, \varrho)$  and  $Y \in \mathcal{K}(G - X)$ , there is exactly one  $y \in X$  such that  $\mu(y, Y) = 1$ . We shall call  $y$  the projection of  $Y$  in  $X$ .*

We shall say that a vertex  $b$  of  $(G, \varrho)$  lies between vertices  $a$  and  $c$  when  $b$  belongs to any path connecting  $a$  with  $c$ . We write  $[a, b, c]$  in this case.

1.1. a) For  $x, y \in G$  we have  $[x, x, x]$ ,  $[x, x, y]$ ,  $[x, y, y]$ .

b) For  $x, y, z \in G$  we have

$$[x, z, y] \Rightarrow [y, z, x].$$

1.2. Let  $(M, \varrho_1)$  be a connected subgraph of a graph  $(G, \varrho)$ ,  $a, b, c \in M$ ,  $[a, b, c]$  in  $(G, \varrho)$ . Then  $[a, b, c]$  in  $(M, \varrho_1)$ .

1.3. Let  $(G, \varrho)$ ,  $(G_1, \varrho_1)$  be two graphs,  $f: G \rightarrow G_1$  a map such that

$$[x, z, y] \text{ in } (G, \varrho) \Rightarrow [f(x), f(z), f(y)] \text{ in } (G_1, \varrho_1).$$

Then  $f$  is called a b-mapping.

\* The main results of this paper were presented at the Scientific Colloquium of the Technical High School, Ilmenau, October 1973.

1.4. Let  $f: (G, \varrho) \rightarrow (G_1, \varrho_1)$ ,  $g: (G_1, \varrho_1) \rightarrow (G_2, \varrho_2)$  be b-mappings. Then  $gf$  is a b-mapping, too.

1.5. Graphs  $(G, \varrho)$  together with the class of all b-mappings form a category. This category will be denoted by  $\mathcal{G}$ .

1.5 is evident by 1.4 and the fact that identity mappings are b-mappings.

1.6. Notions of the category theory will be used in the sense of reference [3]. Especially, for a category  $\mathcal{C}$ ,  $[a, b]_{\mathcal{C}}$  is the set of all morphisms from  $a$  to  $b$ .

The following assertion is clear.

1.7. In a 2-connected graph  $(G, \varrho)$   $[a, b, c]$  holds only in cases described by 1.1.a. Therefore  $[(G, \varrho), (G, \varrho)]_{\mathcal{G}} = G^G$  (the set of all mappings of  $G$  in  $G$ ).

1.8. Let  $X \in \mathcal{E}(G, \varrho)$ ,  $Y \in \mathcal{K}(G - X)$ ,  $y$  the projection of  $Y$  in  $X$ . Let  $f: G \rightarrow G$  be defined as follows:

$$f(x) = y \quad \text{for } x \in Y,$$

$$f(x) = x \quad \text{otherwise.}$$

Then  $f$  is a b-mapping.

Proof. Let  $[a, b, c]$  in  $(G, \varrho)$ .

1. If  $\{a, b, c\} \cap Y = \emptyset$ , then clearly  $[f(a), f(b), f(c)]$ :

2. Let  $\text{card } \{a, b, c\} \cap Y \geq 1$ . If  $\{a, c\} \subset Y$ , then by connectivity of  $Y$  we have  $b \in Y$  and hence  $[f(a), f(b), f(c)]$  for  $\text{card } \{a, b, c\} \cap Y \geq 2$  in general, therefore  $[f(a), f(b), f(c)]$ . For the connectivity of  $X$  we cannot have  $a, c \in X$ ,  $b \in Y$ . Hence let  $a \in Y$ ,  $b, c \in X$  (the case  $a, b \in X$ ,  $c \in Y$  is dual). Then  $f(a) = y$ ,  $f(b) = b$ ,  $f(c) = c$ . Let  $y = c_1, c_2, \dots, c$ ,  $y = d_1, d_2, \dots, d_s$ ,  $a$  be some paths in  $(G, \varrho)$ . Then  $a, d_s, \dots, d_2, y, c_2, \dots, c$  is a path from  $a$  to  $c$ . As  $d_s, \dots, d_2 \in Y$ ,  $b$  must be an element of  $\{y, c_2, \dots, c\}$ . Therefore  $[f(a), f(b), f(c)]$ .

1.9. Let  $X \in \mathcal{E}(G, \varrho)$ . Let  $\mathcal{Y}$  be a system of some  $Y_i \in \mathcal{K}(G - X)$ . Let  $y_i$  be always the projection of  $Y_i$  in  $X$ . Define  $f: G \rightarrow G$  so that

$$f(x) = y_i \quad \text{for } x \in Y_i \quad (\text{for all } i),$$

$$f(x) = x \quad \text{otherwise.}$$

Then  $f$  is a b-mapping.

The proof follows from 1.8. by composing suitable b-mappings.

1.10. Let  $a, b \in (G, \varrho)$ . Then there exists such a b-mapping  $f: G \rightarrow G$  for which  $f(G) = \{a, b\}$ ,  $f(a) = a$ ,  $f(b) = b$ .

Proof. If there exists a 2-component of  $(G, \varrho)$  containing both  $a$  and  $b$ , we get 1.10 immediately by 1.2, 1.7, 1.8.

Suppose  $a \in X \in \mathcal{E}(G, \varrho)$ ,  $b \notin X$ ,  $b \in Y \in \mathcal{K}(G - X)$ . For  $x \in Y$  define  $f(x) = b$ ,  $f(x) = a$  otherwise. Let  $[x, y, z]$  in  $(G, \varrho)$ . Everything will be proved if

$f(x) - f(z) \Rightarrow f(x) = f(y)$ . Suppose we have  $f(x) = f(z) \neq f(y)$ . We have to distinguish two cases.

1)  $f(x) = f(z) = a$ ,  $f(y) = b$ . Then  $x, z \in G - Y$ ,  $y \in Y$ , a contradiction with the connectivity of  $G - Y$ .

The second case

2)  $f(x) = f(z) = b$ ,  $f(y) = a$  is similar.

1.11. (Corollary). Let  $T$  be a two-vertex tree with vertices  $a_1, b_1$ . Let  $a, b \in \in (G, \rho)$ ,  $a \neq b$ . There exists such a b-mapping  $f: G \rightarrow T$  for which

$$f(a) = a_1, \quad f(b) = b_1.$$

1.12. Let  $a, b, a_1, b_1 \in (G, \rho)$ ,  $a \neq b$ . Then there exists such a b-mapping  $f: G \rightarrow G$  such that  $f(G) = \{a_1, b_1\}$ ,  $f(a) = a_1$ ,  $f(b) = b_1$ .

2. A tree is a graph  $(G, \rho)$  without circles (i.e.  $(G, \rho)$  has only one-element 2-components). Thus, one-vertex trees are considered, too.

2.1. Let  $(G, \rho)$  be a tree,  $a \neq b \neq c \neq a$  are vertices of  $G$ ,  $[a, b, c]$  in  $(G, \rho)$ . Then there exists a b-mapping  $f: G \rightarrow G$  such that  $f(G) = \{a, b, c\}$ ,  $f(a) = a$ ,  $f(b) = b$ ,  $f(c) = c$ .

Proof. Let  $a = a_0, a_1, \dots, a_i = b, \dots, a_s = c$  be the path connecting  $a$  with  $c$ . If  $x \in X \in \mathcal{K}(G - \{a_0, \dots, a_s\})$  and  $\mu(X, a_i) = 1$ , we put  $f(x) = a_i$ . Let  $f(a_i) = a_i$  for  $i = 0, 1, \dots, s$ . Similarly as in 1.8. one sees that  $f$  is a b-mapping. Let  $g: \{a_0, \dots, a_s\} \rightarrow \{a_0, a_i, a_s\}$  be defined by  $g(a_0) = g(a_i) = \dots$

$g(a_{i-1}) = a_0$ ,  $g(a_i) = a_i$ ,  $g(a_{i+1}) = \dots = g(a_s) = a_s$ .  $g$  is a b-mapping of  $(\{a_0, \dots, a_s\}, \rho)$  into itself and  $gf$  is a demanded mapping.

2.2. (Corollary). Let  $T'$  be a three-vertex tree with the vertices  $a_1, b_1, c_1$  for which  $[a_1, b_1, c_1]$  in  $T'$ . Let  $(G, \rho)$  be a tree,  $a, b, c \in G$ ,  $a \neq b \neq c \neq a$  and  $[a, b, c]$  in  $(G, \rho)$ . There exists a b-mapping  $f$  from  $(G, \rho)$  to  $T'$  such that  $f(a) = a_1$ ,  $f(b) = b_1$ ,  $f(c) = c_1$ .

2.3. (Corollary). Let  $[a, b, c]$ ,  $[a', b', c']$  in a tree  $(G, \rho)$ ,  $a \neq b \neq c \neq a$ ,  $a' \neq b' \neq c' \neq a'$ . There exists such a b-mapping  $f: G \rightarrow G$  for which  $f(G) = \{a', b', c'\}$  and  $f(a) = a'$ ,  $f(b) = b'$ ,  $f(c) = c'$ .

In [4] the following two propositions have been proved.

2.4. Let  $(G, \rho)$  be a tree,  $a, b, c \in G$ . Then there exists exactly one  $d \in G$  such that  $[a, d, b]$ ,  $[b, d, c]$ ,  $[a, d, c]$ . We write  $\omega(a, b, c) = d$  and the ternary operation  $\omega$  is called the intersection operation.

2.5. Let  $(G, \rho)$  be a tree. Then

$$1) [a, b, c] \Leftrightarrow \omega(a, b, c) = b.$$

$$2) \omega(a, a, b) = \omega(a, b, a) = \omega(b, a, a) = a.$$

3)  $\omega$  is a symmetrical operation.

Similarly as in 2.1. the following can be proved.

2.6. Let  $(G, \varrho)$  be a tree. Let  $a, b, c \in G$  and  $\omega(a, b, c) \notin \{a, b, c\}$ . There exists such a b-mapping  $f: G \rightarrow G$  for which  $f(G) = \{a, b, c, \omega(a, b, c)\}$ ,  $f(a) = a$ ,  $f(b) = b$ ,  $f(c) = c$ ,  $f(\omega(a, b, c)) = \omega(a, b, c)$ .

2.7. (Corollary). Let  $(G, \varrho)$  be a tree,  $a, b, c, a_1, b_1, c_1 \in G$ ,  $\omega(a, b, c) \notin \{a, b, c\}$ ,  $\omega(a_1, b_1, c_1) \notin \{a_1, b_1, c_1\}$ . There exists a b-mapping  $f: G \rightarrow G$  such that  $f(G) = \{a_1, b_1, c_1, \omega(a_1, b_1, c_1)\}$ ,  $f(a) = a_1$ ,  $f(b) = b_1$ ,  $f(c) = c_1$ ,  $f(\omega(a, b, c)) = \omega(a_1, b_1, c_1)$ .

3. The notions from the theory of universal algebras are taken from [2]. The category of all algebras with one fundamental operation of a fixed arity (generally denoted by  $\alpha$ ) will be denoted by  $\mathcal{A}_\alpha$ . In considerations on morphisms of  $\mathcal{A}_\alpha$  (and of  $\mathcal{G}$ , as well) we identify these mappings with the carrying mappings of the supports of relevant structures.

3.1. Let  $(G, \varrho), (G_1, \varrho_1)$  be trees,  $\omega$  denote (for both of them) the intersection operation. Then

$$[(G, \omega), (G_1, \omega)]_{\mathcal{A}_\alpha} = [(G, \varrho), (G_1, \varrho_1)]_{\mathcal{G}}.$$

Proof. Let  $f \in [(G, \omega), (G_1, \omega)]_{\mathcal{A}_\alpha}$  and  $[a, b, c]$  in  $(G, \varrho)$ .  $\omega(a, b, c) = b$  therefore  $f(b) = \omega(f(a), f(b), f(c))$ , and so  $[f(a), f(b), f(c)]$  in  $(G_1, \varrho_1)$ .

Let  $f \in [(G, \varrho), (G_1, \varrho_1)]_{\mathcal{G}}$  and  $\omega(a, b, c) = d$  in  $(G, \omega)$ . We have  $[a, d, b]$ ,  $[b, d, c]$ ,  $[a, d, c]$ , hence  $[f(a), f(d), f(b)]$ ,  $[f(b), f(d), f(c)]$ ,  $[f(a), f(d), f(c)]$  and therefore  $\omega(f(a), f(b), f(c)) = f(d)$ .

3.2. Let  $[(G, \varrho), (G, \varrho)]_{\mathcal{G}} \subset [(G, \alpha), (G, \alpha)]_{\mathcal{A}_\alpha}$ . Then  $\alpha$  is an idempotent operation.

This follows from the fact that  $[(G, \varrho), (G, \varrho)]_{\mathcal{G}}$  contains all constant mappings.

3.3. Let  $[(G, \varrho), (G, \varrho)]_{\mathcal{G}} \subset [(G, \alpha), (G, \alpha)]_{\mathcal{A}_\alpha}$ , where  $\alpha$  is a binary operation. Then  $\alpha$  is a projection.

Proof. Let  $\{a, b\} \subset G$ ,  $a \neq b$ . By 1.10.  $\{a, b\}$  is a subalgebra in  $(G, \alpha)$ . Let, e.g.  $\alpha(a, b) = a$ . Let  $a_1, b_1 \in G$ . By 1.12  $f(a) = a_1, f(b) = b_1$  for certain  $f \in [(G, \varrho), (G, \varrho)]_{\mathcal{G}}$ . Therefore  $\alpha(a_1, b_1) = a_1$  and so  $\alpha$  is a projection.

3.4. Let  $(G, \varrho)$  be a tree,  $\alpha$  essentially ternary on  $G$ , and  $[(G, \varrho), (G, \varrho)]_{\mathcal{G}} \subset [(G, \alpha), (G, \alpha)]_{\mathcal{A}_\alpha}$ . If  $\text{card } G = 2$ , then  $(G, \alpha)$  is a Post algebra, if  $\text{card } G > 2$ , we have  $\alpha = \omega$ . In both cases,

$$[(G, \varrho), (G, \varrho)]_{\mathcal{G}} = [(G, \alpha), (G, \alpha)]_{\mathcal{A}_\alpha}.$$

Proof. If  $\text{card } G = 2$ ,  $[(G, \varrho), (G, \varrho)]_{\mathcal{G}} = G^G$  and this implies by [2], § 6 that  $(G, \alpha)$  is a Post algebra. Let  $\text{card } G > 2$ ,  $a, b, c \in G$ ,  $a \neq b \neq c \neq a$   $[a, b, c]$  in  $(G, \varrho)$ . By 2.1.  $\{a, b, c\}$  is a subalgebra in  $(G, \alpha)$ . Suppose  $\alpha(a, b, c) = a$ . In the sequel, we use 1.12 and 2.3 without any further reference. E.g.,  $\alpha(a, b, c) = a$  by 1.12 implies  $\alpha(a, b, b) = a$ ,  $\alpha(a, a, c) = a$ . As  $\alpha(x, y, y)$ ,  $\alpha(x, x, y)$  are projections, we have  $\alpha(x, y, y) = x$ ,  $\alpha(x, x, y) = x$ .

By 2.3. we have  $\alpha(c, b, a) = c$ .

1. Suppose  $\alpha(b, a, c) = a$ , whence  $\alpha(x, y, x) = y$ .

1.1. If  $\alpha(a, c, b) = a$ , then  $\alpha(x, y, x) = x$ , a contradiction.

1.2. If  $\alpha(a, c, b) = b$ , then  $\alpha(x, y, y) = y$ , a contradiction.

1.3. If  $\alpha(a, c, b) = c$ , then  $\alpha(x, y, y) = y$  again.

2. Suppose  $\alpha(b, a, c) = c$ . Then  $\alpha(x, x, y) = y$ , a contradiction. Therefore  $\alpha(b, a, c) = b$ , hence  $\alpha(b, c, a) = b$  and  $\alpha(x, y, x) = x$ , therefore  $\alpha(c, a, b)$  can be  $c$  or  $b$ , but  $\alpha(x, y, y) = x$  implies  $\alpha(c, a, b) = c$ . Therefore, if  $x_1, y_1, z_1 \in \{a, b, c\}$ , we have  $\alpha(x_1, y_1, z_1) = x_1$ . By 2.3 this is simultaneously true for all triples  $a', b', c'$  with  $[a', b', c']$ .

Let now  $a', b', c' \in G$ ,  $\omega(a', b', c') \notin \{a', b', c'\}$ . By 2.6  $\{a', b', c', \omega(a', b', c')\}$  is a subalgebra. Suppose  $\alpha(a', b', c') = b'$ . Then (e.g., by 1.8)  $\alpha(a', \omega(a', b', c'), c') - \omega(a', b', c')$ , which is a contradiction, because  $[a', \omega(a', b', c'), c']$ . In the same way it turns out that  $\alpha(a', b', c') = \omega(a', b', c')$ ,  $\alpha(a', b', c') = c'$  are contradictory, too. Therefore  $\alpha(a', b', c') = a'$ , and so  $\alpha(x, y, z) = x$ , which is a contradiction to the supposition on  $\alpha$ .

Similarly for  $\alpha(a, b, c) = c$  and for

$$\alpha(a, b, c) = b, \quad \alpha(b, a, c) = a.$$

Suppose  $\alpha(a, b, c) = b$ ,  $\alpha(b, a, c) = c$ .  $\alpha(a, b, c) = b$  implies  $\alpha(x, x, y) = x$ ,  $\alpha(b, a, c) = c$  implies  $\alpha(x, x, y) = y$ , a contradiction.

Therefore we must have  $\alpha(a, b, c) = \alpha(b, a, c) = b$ . Then  $\alpha(x, x, y) = \alpha(y, x, x) = \alpha(x, y, x) = x$  and this implies  $\alpha(a, c, b) = b$ . Thus for  $x_1, y_1, z_1 \in \{a, b, c\}$  we have  $\alpha(x_1, y_1, z_1) = \omega(x_1, y_1, z_1)$  and this holds for all triples  $a', b', c'$  with  $[a', b', c']$  and by upper arguments for all triples in general.

3.5. If  $(G, \rho)$  is not a tree,  $\alpha$  a ternary operation on  $G$  and  $[(G, \rho), (G, \rho)]_{\mathcal{G}} \subset [(G, \alpha), (G, \alpha)]_{\mathcal{A}}$ , then  $\alpha$  is a projection.

Proof. Let  $X \in \mathcal{E}(G, \rho)$ ,  $\text{card } X \geq 3$ . Let  $f: G \rightarrow X$  from 1.6 for  $\mathcal{Y} = \mathcal{X}(G - X)$ . By 1.2. and 1.7.  $[(X, \rho), (X, \rho)]_{\mathcal{G}} = X^X$ . Therefore  $X$  is a subalgebra in  $(G, \alpha)$  and  $\alpha$  restricted to  $X$  (notation  $\alpha/X$ ) is a projection.

Let, e.g.  $\alpha/X$  be the projection to the first coordinate, i.e. for  $a, b, c \in X$  we have  $\alpha(a, b, c) = a$ . Therefore  $\alpha(x, x, y) = \alpha(x, y, x) = \alpha(x, y, y) = x$  on  $X$ . But as  $\text{card } X \geq 3$  and  $\alpha(x, x, y)$ ,  $\alpha(x, y, x)$ ,  $\alpha(x, y, y)$  are projections on the whole  $G$  by 3.3., the mentioned equalities hold on  $G$ . Suppose we have  $a, b, c \in G$  so that  $a \neq \alpha(a, b, c)$ . Let  $f$  be a b-mapping from 1.10. for which  $\alpha(a, b, c)$  stands for  $b$ .

It is  $\alpha(a, b, c) = f(\alpha(a, b, c)) = \alpha(f(a), f(b), f(c)) = \alpha(a, f(b), f(c))$ . But  $\text{card } \{a, f(b), f(c)\} \leq 2$  and therefore  $\alpha(a, f(b), f(c)) = a$ , a contradiction.

3.6. Let  $(G, \rho)$  be a graph and  $\alpha$  be an essentially  $n$ -ary operation on  $G$  for some  $n \geq 3$  such that  $\alpha(x_1, \dots, x_n) = x_1$ , whenever  $\text{card } \{x_1, \dots, x_n\} < n$ . Then there cannot hold

$$[(G, \varrho), (G, \varrho)]_{\mathcal{G}} \subset [(G, \alpha), (G, \alpha)]_{\mathcal{A}_\alpha}.$$

**Proof.** Suppose the upper inclusion to be true. As  $\alpha$  is essentially  $n$ -ary,  $\alpha$  is not a projection and there exist  $a_1, \dots, a_n \in G$  such that  $\alpha(a_1, \dots, a_n) \neq a_1$ . Let  $f$  be a  $b$ -mapping from 1.10, where  $\alpha(a_1, \dots, a_n)$  stands for  $b$ . We have  $\alpha(f(a_1), \dots, f(a_n)) = f(\alpha(a_1, \dots, a_n)) = \alpha(a_1, \dots, a_n)$ . But  $\text{card} \{f(a_1), \dots, f(a_n)\} \leq 2$ , i.e.  $\alpha(f(a_1), \dots, f(a_n)) = f(a_1) = a_1$ , a contradiction.

By 3.1.—3.6, taking in account Lemma 20 from [5] we get the final result of this section.

3.7. Let  $[(G, \varrho), (G, \varrho)]_{\mathcal{G}} \subset [(G, \alpha), (G, \alpha)]_{\mathcal{A}_\alpha}$ ,  $\text{card } G \geq 3$ , where  $\alpha$  is an operation on  $G$ . Then

1) if  $(G, \varrho)$  is a tree and  $\alpha$  is essentially  $n$ -ary for certain  $n \geq 2$ ,  $\omega$  is an algebraic operation in the algebra  $(G, \alpha)$  and  $[(G, \varrho), (G, \varrho)]_{\mathcal{G}} = [(G, \alpha), (G, \alpha)]_{\mathcal{A}_\alpha}$ .

2) if  $(G, \varrho)$  is not a tree, then  $\alpha$  is a projection.

4. Let  $\mathcal{T}$  be the full subcategory of  $\mathcal{G}$  consisting of all trees. Let  $F$  be the embedding of  $\mathcal{T}$  in  $\mathcal{A}_\alpha$  where  $\alpha$  is ternary given by  $F[(G, \varrho)] = (G, \omega)$  and  $F(f) = f$  for mappings. Let now  $(G, \varrho)$  be an arbitrary graph from  $\mathcal{G}$ . Let  $F'(G, \varrho)$  be the algebra from  $\mathcal{A}_\alpha$  generated by the set  $\mathbf{G}$  which is equal to  $G$  if  $(G, \varrho)$  is a tree and to  $G \times \{(G, \varrho)\}$  in other cases (we write  $\mathbf{x}$  instead of  $x$  or  $\langle x, (G, \varrho) \rangle$  in the sequel) with the following defining relations

$$\begin{aligned} \alpha(\mathbf{x}, \mathbf{y}, \mathbf{z}) &= \mathbf{v} \quad \text{iff } x, y, z, v \in G \quad \text{and} \\ [x, v, y], [y, v, z], [x, v, z] &\text{ in } (G, \varrho) \end{aligned} \quad (*)$$

4.1. If  $(G, \varrho)$  is a tree,  $F'(G, \varrho) = (G, \omega)$ . Clear.

4.2. For  $w \in F'(G, \varrho)$  we have  $\alpha(w, w, w) = w$  iff  $w \in \mathbf{G}$  (the equality is meant as an equality of elements in  $F'(G, \varrho)$  not an identity of terms).

**Proof.** (\*) and 1.1 imply  $\alpha(\mathbf{x}, \mathbf{x}, \mathbf{x}) = \mathbf{x}$  for  $x \in G$ . Let  $w$  be a term from  $F'(G, \varrho)$  having the minimal length among all terms which are equivalent to  $w$ . If this length is not 1, then terms equivalent to  $\alpha(w, w, w)$  are of the form  $\alpha(w_1, w_2, w_3)$ , where  $w_1, w_2, w_3$  are terms equivalent to  $w$ , therefore of a length greater or equal to the length of  $w$  and so the length of  $\alpha(w_1, w_2, w_3)$  is greater than the length of  $w$ . Therefore  $\alpha(w, w, w) \neq w$ .

4.3. Let  $f \in [F'(G, \varrho), F'(G_1, \varrho_1)]_{\mathcal{A}_\alpha}$ . Then  $f(\mathbf{G}) \subset \mathbf{G}_1$ .

The proof follows immediately by 4.2.

4.4. Let  $f \in [F'(G, \varrho), F'(G_1, \varrho_1)]_{\mathcal{A}_\alpha}$ . Define  $f' : G \rightarrow G_1$  by

$$f'(x) = y = f(\mathbf{x}) = \mathbf{y}. \quad \text{Then } f' \in [(G, \varrho), (G, \varrho_1)]_{\mathcal{G}}.$$

**Proof.** Let  $[x, y, z]$  in  $(G, \varrho)$ . Then  $\alpha(\mathbf{x}, \mathbf{y}, \mathbf{z}) = \mathbf{y}$  and so  $\alpha(f(\mathbf{x}), f(\mathbf{y}), f(\mathbf{z})) = f(\mathbf{y})$  in  $F'(G_1, \varrho_1)$ . Hence  $[f'(x), f'(y), f'(z)]$  in  $(G_1, \varrho_1)$ .

4.5. Let  $f_1 \in [(G, \varrho), (G_1, \varrho_1)]_{\mathcal{G}}$ . Then there exists exactly one  $f \in [F'(G, \varrho), F'(G_1, \varrho_1)]_{\mathcal{A}_\alpha}$  such that  $f' = f_1$  (notation as in 4.4).

Proof. Let  $f_1^*(\mathbf{x}) = \mathbf{y}$  iff  $f_1(x) = y$ . By (\*) and by the definition of the  $\mathbf{b}$ -mapping  $f^*$  preserves all defining relations. Therefore there exists exactly one  $f$  with the demanded properties.

$f$  from 4.5. will be denoted by  $F'(f_1)$ . An immediate consequence of 4.4 and 4.5 is

4.6. The functor  $F' : \mathcal{G} \rightarrow \mathcal{A}_\alpha$  is a full embedding, which extends  $F : \mathcal{F} \rightarrow \mathcal{A}_\alpha$ .

4.7. It can be easily proved that  $F'$  is the left Kan extension of  $F$  (see [1], chapter X).

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