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# MODAL OPERATORS ON BOUNDED COMMUTATIVE RESIDUATED $\ell$-MONOIDS 

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#### Abstract

Bounded commutative residuated lattice ordered monoids ( $R \ell$-monoids) are a common generalization of, e.g., Heyting algebras and $B L$-algebras, i.e., algebras of intuitionistic logic and basic fuzzy logic, respectively. Modal operators (special cases of closure operators) on Heyting algebras were studied in [MacNAB, D. S.: Modal operators on Heyting algebras, Algebra Universalis 12 (1981), 5-29] and on $M V$-algebras in [HARLENDEROVÁ, M.-RACHU゚NEK, J.: Modal operators on MV-algebras, Math. Bohem. 131 (2006), 39-48]. In the paper we generalize the notion of a modal operator for general bounded commutative $R \ell$-monoids and investigate their properties also for certain derived algebras.


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Commutative residuated lattice ordered monoids ( $R \ell$-monoids) are duals to commutative $D R \ell$-monoids which were introduced by S wamy [16] as a common generalization of Abelian lattice ordered groups and Brouwerian algebras. By [11], [12], [13], also algebras of logics behind fuzzy reasoning can be considered as particular cases of bounded commutative $R \ell$-monoids. Namely from this point of view, $M V$-algebras, an algebraic counterpart of the Lukasiewicz infinitevalued propositional logic, are precisely bounded commutative $R \ell$-monoids satisfying the double negation law. Further, $B L$-algebras, an algebraic semantics of the Hájek basic fuzzy logic, are just bounded commutative $R \ell$-monoids isomorphic to subdirect products of linearly ordered commutative $R \ell$-monoids. Heyting algebras (duals to Brouwerian algebras), i.e. algebras of intuitionistic logic, are characterized as bounded commutative $R \ell$-monoids with idempotent multiplication.

[^0]Modal operators (special cases of closure operators) on Heyting algebras were introduced and studied by Macnab in [10]. Analogously, modal operators on $M V$-algebras were introduced in [7] recently.

In this paper we define modal operators for arbitrary bounded commutative $R \ell$-monoids and we study their properties in the class of normal $R \ell$-monoids in particular.

For concepts and results relating to $M V$-algebras, $B L$-algebras and Heyting algebras see for instance [3], [6], [1].

Definition 1. A bounded commutative $R \ell$-monoid is an algebra $M=(M ; \cdot$, $\wedge, \rightarrow, 0,1)$ of type $\langle 2,2,2,2,0,0\rangle$ satisfying the following conditions.
(i) $(M ; \odot, 1)$ is a commutative monoid.
(ii) $(M ; \vee, \wedge, 0,1)$ is a bounded lattice.
(iii) $x \odot y \leq z$ if and only if $x \leq y \rightarrow z$, for any $x, y, z \in M$.
(iv) $x \odot(x \rightarrow y)=x \wedge y$, for any $x, y \in M$.

Bounded commutative $R \ell$-monoids are special cases of residuated lattices, more precisely (see for instance [4]), they are exactly commutative integral generalized $B L$-algebras in the sense of [2] and [8].

In what follows, by an $R \ell$-monoid we will mean a bounded commutative $R \ell$-monoid.

Let us define on any $R \ell$-monoid $M$ the unary operation of negation "-" by $x^{-}:=x \rightarrow 0$ for any $x \in M$. Further, we put $x \oplus y:=\left(\begin{array}{ll}x & y^{-}\end{array}\right)^{-}$for any $x, y \in M$.

Algebras of the above mentioned propositional logics can be characterized in the class of all $R \ell$-monoids as follows: An $R \ell$-monoid $M$ is
a) a $B L$-algebra ([13]) if and only if $M$ satisfies the identity of pre-linearity $(x \rightarrow y) \vee(y \rightarrow x)=1 ;$
b) an $M V$-algebra ([11], [12]) if and only if $M$ fulfills the double negation law $x^{--}=x ;$
c) a Heyting algebra ([16]) if and only if the operations ". " and " $\wedge$ " coincide on $M$.

Lemma 1. ([16], [15]) In any bounded commutative $R \ell$-monoid $M$ we have for any $x, y \in M$ :
(1) $x \leq y \Longleftrightarrow x \rightarrow y=1$.
(2) $x \odot y \leq x \wedge y \leq x, y$.
(3) $x \leq y \Longrightarrow x \odot z \leq y \odot z$.
(4) $x \leq y \Longrightarrow z \rightarrow x \leq z \rightarrow y, y \rightarrow z \leq x \rightarrow z$.
(5) $(x \odot y) \rightarrow z=x \rightarrow(y \rightarrow z)=y \rightarrow(x \rightarrow z)$.
(6) $(x \rightarrow y) \odot(y \rightarrow z) \leq x \rightarrow z$.
(7) $1^{--}=1,0^{--}=0$.
(8) $x \leq x^{--}, x^{-}=x^{---}$.
(9) $x \leq y \Longrightarrow y^{-} \leq x^{-}$.
(10) $(x \vee y)^{-}=x^{-} \wedge y^{-}$.
(11) $(x \wedge y)^{--}=x^{--} \wedge y^{--}$.
(12) $(x \odot y)^{-}=y \rightarrow x^{-}=y^{--} \rightarrow x^{-}=x \rightarrow y^{-}=x^{--} \rightarrow y^{-}$.
(13) $(x \odot y)^{--} \geq x^{--} \odot y^{--}$.
(14) $(x \rightarrow y)^{--}=x^{--} \rightarrow y^{--}$.

Remark 2. It is obvious that $x \oplus z \leq y \oplus z$ holds for any $x, y, z \in M$ such that $x \leq y$. Further by [14, Lemma 2.11], $x^{--} \oplus y^{--}=x \oplus y$ for any $x, y \in M$, hence also $x \oplus y=x^{--} \oplus y=x \oplus y^{--}=x^{--} \oplus y^{--}$.

Definition 2. Let $M$ be an $R \ell$-monoid. A mapping $f: M \longrightarrow M$ is called a modal operator on $M$ if, for any $x, y \in M$,

1. $x \leq f(x)$;
2. $f(f(x))=f(x)$;
3. $f(x \odot y)=f(x) \odot f(y)$.

If, moreover, for any $x, y \in M$,
4. $f(x \oplus y)=f(x \oplus f(y))$,
then $f$ is called a strong modal operator on $M$.
Proposition 3. If $f$ is a modal operator on an $R \ell$-monoid $M$ and $x, y \in M$, then
(i) $x \leq y \Longrightarrow f(x) \leq f(y)$;
(ii) $f(x \rightarrow y) \leq f(x) \rightarrow f(y)=f(f(x) \rightarrow f(y))=x \rightarrow f(y)=f(x \rightarrow f(y))$;
(iii) $f(x) \leq(x \rightarrow f(0)) \rightarrow f(0)$;
(iv) $f(x) \odot x^{-} \leq f(0)$;
(v) $x \oplus f(0) \geq f\left(x^{--}\right) \geq f(x)$;
(vi) $f(x \vee y)=f(x \vee f(y))=f(f(x) \vee f(y))$.

Proof.
(i) $x \leq y \Longrightarrow f(x \wedge y)=f(x) \Longrightarrow f(y \odot(y \rightarrow x))=f(x) \Longrightarrow$ $f(y) \odot f(y \rightarrow x)=f(x) \Longrightarrow f(x) \leq f(y)$.
(ii) By (i), $f(x) \odot f(x \rightarrow y)=f(x \odot(x \rightarrow y))=f(x \wedge y) \leq f(y)$.

This implies

$$
f(x \rightarrow y) \leq f(x) \rightarrow f(y)
$$

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From this we get

$$
\begin{aligned}
f(f(x) \rightarrow f(y)) & \leq f(f(x)) \rightarrow f(f(y))=f(x) \rightarrow f(y) \leq x \rightarrow f(y) \\
& \leq f(x \rightarrow f(y)) \leq f(x) \rightarrow f(f(y)) \\
& =f(x) \rightarrow f(y) \leq f(f(x) \rightarrow f(y))
\end{aligned}
$$

therefore

$$
f(x) \rightarrow f(y)=f(f(x) \rightarrow f(y))=x \rightarrow f(y)=f(x \rightarrow f(y))
$$

(iii) By use of (ii) and (i), we have

$$
\begin{aligned}
f(x) \odot(f(x) \rightarrow f(0))=f(x) \wedge f(0)=f(0) & \Longrightarrow f(x) \leq(f(x) \rightarrow f(0)) \rightarrow f(0) \\
& \Longrightarrow f(x) \leq(x \rightarrow f(0)) \rightarrow f(0)
\end{aligned}
$$

(iv) By (ii), we obtain

$$
0 \leq f(0) \Longrightarrow x^{-}=x \rightarrow 0 \leq x \rightarrow f(0)=f(x) \rightarrow f(0)
$$

thus

$$
f(x) \odot x^{-} \leq f(x) \odot(f(x) \rightarrow f(0))=f(x) \wedge f(0)=f(0)
$$

(v) According to Remark 2, Lemma 1(12), (8) and the part (ii) consecutively,

$$
\begin{aligned}
x \oplus f(0) & =x^{--} \oplus f(0)=\left(x^{---} \odot f(0)^{-}\right)^{-}=x^{---} \rightarrow f(0)^{--} \\
& =x^{-} \rightarrow f(0)^{--}=f\left(x^{-} \rightarrow f(0)^{--}\right) \geq f\left(x^{-} \rightarrow f(0)\right) \geq f\left(x^{-} \rightarrow 0\right) \\
& =f\left(x^{--}\right) \geq f(x)
\end{aligned}
$$

Hence

$$
x \oplus f(0) \geq f\left(x^{--}\right) \geq f(x)
$$

(vi) $f(x \vee y) \leq f(x \vee f(y)) \leq f(f(x) \vee f(y))=f(f(x \vee y))=f(x \vee y)$.

Remark 4. By the definition of a modal operator and Proposition 3(i) every modal operator on an $R \ell$-monoid $M$ is a closure operator on the lattice ( $M ; \vee, \wedge$ ).

Remark 5. M. Galatos and C. Tsinakis introduced in [5] the notion of a nucleus of a residuated lattice $L$ as a closure operator $\gamma$ on $L$ satisfying $\gamma(a) \gamma(b) \leq \gamma(a b)$, to represent generalizations of $M V$-algebras (dropping integrality, commutativity and the existence of bounds) by means of $\ell$-groups and nuclei of negative cones of $\ell$-groups. From this point of view, a modal operator $f$ on an $R \ell$-monoid $M$ is a nucleus of $M$ satisfying $f(x) \odot f(y) \geq f(x \odot y)$.

Proposition 6. If $f$ is a strong modal operator on an $R \ell$-monoid $M$ and $x, y \in M$, then
(vii) $f(x \oplus y)=f(f(x) \oplus f(y))$;
(viii) $x \oplus f(0)=f\left(x^{--}\right)$.

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Proof. Let us suppose that $f$ is a strong modal operator. Then
(vii) $f(x \oplus y)=f(x \oplus f(y))=f(f(x) \oplus f(y))$;
(viii) By Proposition 3(v), $f(x \oplus f(0))=f(x \oplus 0)=f\left(x^{--}\right)$implies $f\left(x^{--}\right)=$ $f(x \oplus f(0)) \geq x \oplus f(0) \geq f\left(x^{--}\right)$.

Theorem 7. Let $M$ be an $R \ell$-monoid and $f: M \longrightarrow M$ be a mapping. Then $f$ is a modal operator on $M$ if and only if for any $x, y \in M$ it is satisfied:
(a) $x \rightarrow f(y)=f(x) \rightarrow f(y)$;
(b) $f(x) \odot f(y) \geq f(x \odot y)$.

Proof. Let a mapping fulfil conditions (a) and (b).

1. For any $x \in M$ we have $x \rightarrow f(x)=f(x) \rightarrow f(x)=1$. Therefore $x \leq f(x)$.
2. For all $x \in M$ it holds $1=f(x) \rightarrow f(x)=f(f(x)) \rightarrow f(x)$. This implies $f(f(x)) \leq f(x)$. Therefore, by $1, f(f(x))=f(x)$.
3. For any $x, y \in M$ it is true
$x \odot y \leq f(x \odot y) \Longrightarrow y \leq x \rightarrow f(x \odot y)=f(x) \rightarrow f(x \odot y) \Longrightarrow y \odot f(x) \leq$ $f(x \odot y) \Longrightarrow f(x) \leq y \rightarrow f(x \odot y)=f(y) \rightarrow f(x \odot y) \Longrightarrow f(x) \odot f(y) \leq$ $f(x \odot y) \Longrightarrow f(x) \odot f(y)=f(x \odot y)$.
The converse implication is obvious.
Corollary 8. If $M$ is an $R \ell$-monoid and $f: M \longrightarrow M$ is a mapping, then $f$ is a nucleus of $M$ if and only if $f$ satisfies (a) of Theorem 7 and it is isotone.

Remark 9. If $M$ is a Heyting algebra and $x, y \in M$, then $f(x) \odot f(y)=$ $f(x) \wedge f(y) \geq f(x \wedge y)=f(x \odot y)$. Therefore, by Theorem $7, f$ is a modal operator on $M$ iff it satisfies condition (a) (see also [10]).

We say that an $R \ell$-monoid $M$ is normal if $M$ satisfies the identity

$$
(x \odot y)^{--}=x^{--} \odot y^{--}
$$

Remark 10. By [15, Proposition 5], every $B L$-algebra and every Heyting algebra is normal, hence the variety of normal $R \ell$-monoids is considerably wide.

Let $M$ be an $R \ell$-monoid. For arbitrary element $a \in M$ we denote by $\varphi_{a}: M \longrightarrow M$ the mapping such that $\varphi_{a}(x)=a \oplus x$ for every $x \in M$.

Denote by

$$
I(M)=\{a \in M: a \odot a=a\}
$$

the set of all multiplicative idempotents in an $R \ell$-monoid $M$. It is obvious that $0,1 \in I(M)$. By [9, Lemma 2.8.3], $a \odot x=a \wedge x$ holds for any $a \in I(M), x \in M$. Further, if $M$ is a normal $R \ell$-monoid and $a \in I(M)$, then also $a^{--} \in I(M)$.

Theorem 11. If $M$ is a normal $R \ell$-monoid and $a \in M$, then $\varphi_{a}$ is a strong modal operator on $M$ if and only if $a^{-}, a^{--} \in I(M)$.

Proof.
a) Let $a, x, y \in M, a^{-}, a^{--} \in I(M)$.

1. $\varphi_{a}(x)=a \oplus x=\left(a^{-} \odot x^{-}\right)^{-} \geq x^{--} \geq x$.
2. $\varphi_{a}\left(\varphi_{a}(x)\right)=a \oplus(a \oplus x)=a \oplus\left(a^{-} \odot x^{-}\right)^{-}=\left(a^{-} \odot\left(a \quad \cdot x^{-}\right)^{--}\right)^{-}$ $\left(a^{-} \odot\left(a^{-} \odot x^{-}\right)\right)^{-}=\left(\left(a^{-} \odot a^{-}\right) \odot x^{-}\right)^{-}=\left(a^{-} \cdot x^{-}\right)^{-} \quad a \oplus x=\varphi_{a}(x)$.
3. We first prove that $a \oplus x=(a \vee x)^{--}$.

By Lemma $1(10)$, we obtain $a \oplus x-\left(a^{-} \odot x^{-}\right)^{-}-\left(a^{-} \wedge x^{-}\right)$ $\left((a \vee x)^{-}\right)^{-}-(a \vee x)^{--}$.

We will now prove condition 3 from the definition of a modal operator.
We have

$$
\begin{aligned}
\varphi_{a}(x) \odot \varphi_{a}(y) & =(a \oplus x) \odot(a \oplus y)=\left(a^{--} \oplus x\right) \odot\left(a^{--} \oplus y\right) \\
& =\left(a^{--} \vee x\right)^{--} \odot\left(a^{--} \vee y\right)^{--}=\left(\left(a^{--} \vee x\right) \cdot\left(a^{--} \vee y\right)\right)^{--} \\
& =\left(\left(a^{--} \odot a^{---}\right) \vee\left(x \odot a^{--}\right) \vee\left(a^{--} \odot y\right) \vee(x \cdot y)\right)^{--} \\
& =\left(a^{--} \vee(x \odot y)\right)^{--}=a^{--} \oplus(x \odot y)=a \oplus(x \cdot y) \\
& =\varphi_{a}(x \odot y) .
\end{aligned}
$$

4. According to [14, Proposition 2.10], $(M ; \oplus)$ is a commutative semigroup.

For this reason

$$
\begin{aligned}
\varphi_{a}(x \oplus y) & =a \oplus(x \oplus y)=a^{--} \oplus(x \oplus y)=(a \oplus a) \oplus(x \oplus y) \\
& =a \oplus(x \oplus(a \oplus y))=\varphi_{a}\left(x \oplus \varphi_{a}(y)\right)
\end{aligned}
$$

b) Let $\varphi_{a}$ be a strong modal operator on $M$. Then on account of condition 3, we have $a \oplus(x \odot y)=(a \oplus x) \odot(a \oplus y)$. Then for $x=y=0$ we obtain $a \oplus(0 \cdot 0)-$ $(a \oplus 0) \odot(a \oplus 0)$, hence $a \oplus 0=(a \oplus 0) \odot(a \oplus 0)$. Since $a \oplus 0-a^{-} \quad$ (see [14, Lemma 2.11]), we conclude that $a^{--}=a^{--} \odot a^{--}$, which yields $a^{--} \in I(M)$

From condition 4 we have $a \oplus(x \oplus y)=a \oplus(x \oplus(a \oplus y))$. Then for $x=y=0$ it follows that $a^{--}=a \oplus 0=a \oplus(0 \oplus 0)=a \oplus(0 \oplus(a \oplus 0))-(a \oplus 0) \oplus a^{--}--$ $a^{--} \oplus a^{--}$, thus $a^{--}=a^{--} \oplus a^{--}$. From this $a^{--}=\left(a^{-} \cdot a^{-}\right)^{-}$, hence $a^{-}=\left(a^{-} \odot a^{-}\right)^{--}$. Since $M$ is normal, we obtain $a^{-}=a^{-} \cdot a^{-}$and so $a^{-} \in I(M)$.

Remark 12. If $M$ is an $M V$-algebra and $a \in M$, then $a \in I(M)$ if and only if $a^{-}, a^{--} \in I(M)$. Concurrently, by [7], in any $M V$-algebra it is true that $\varphi_{a}$ is a modal operator on $M$ if and only if $\varphi_{a}$ is a strong modal operator on $M$ (namely if and only if $a \in I(M)$ ). The question, whether $\varphi_{a}$ is a modal operator on $M$ if and only if it is a strong modal operator also for any normal $R \ell$-monoid $M$, remains open.

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Corollary 13. Let $M$ be a normal $R \ell$-monoid and $f$ be a modal operator on $M$ such that $f(x)=f\left(x^{--}\right)$for all $x \in M$. Then $f$ is strong if and only if $f=\varphi_{f(0)}$ and $f(0)^{-} \in I(M)$.
Proof. Suppose that a modal operator $f$ on $M$ satisfies the condition $f(x)=$ $f\left(x^{--}\right)$for every $x \in M$. Then by Proposition 6 and Theorem 11, $f$ is strong if and only if $f=\varphi_{a}$ for some $a \in M$ such that $a^{-}, a^{--} \in I(M)$.

If $f$ is strong and $x \in M$, then $f(x)=f\left(x^{--}\right)=f(0) \oplus x$. Hence $f=\varphi_{f(0)}$ and we have $f(0), f(0)^{-} \in I(M)$.

For any modal operator $f$ we have $f(0)^{--} \in I(M)$. In fact, $f(0)^{--}=$ $f(0 \odot 0)^{--}=(f(0) \odot f(0))^{--}=f(0)^{--} \odot f(0)^{--}$. Hence, if $f=\varphi_{f(0)}$ and $f(0)^{-} \in I(M)$, then by Theorem 11, $f$ is strong.
Corollary 14. Specially for MV-algebras, we obtain (see [7]): If $M$ is an $M V$-algebra and $f$ is a modal operator on $M$, then $f$ is strong if and only if $f=\varphi_{f(0)}$.

Let $M$ be an $R \ell$-monoid and $a \in M$. Consider mappings $\psi_{a}: M \longrightarrow M$ and $\chi_{a}: M \longrightarrow M$ such that $\psi_{a}(x):=a \rightarrow x$ and $\chi_{a}(x):=(x \rightarrow a) \rightarrow a$ for every $x \in M$. These mappings are significant modal operators in Heyting algebras (see [10]). We will now deal with the mappings $\psi_{a}$ and $\chi_{a}$ in arbitrary $R \ell$-monoids.

Proposition 15. If $M$ is an $R \ell$-monoid and $a \in I(M)$, then for any $x, y \in M$

$$
x \rightarrow \psi_{a}(y)=\psi_{a}(x) \rightarrow \psi_{a}(y)
$$

Proof. By the definition of $\psi_{a}, x \rightarrow \psi_{a}(y)=x \rightarrow(a \rightarrow y)$ and $\psi_{a}(x) \rightarrow$ $\psi_{a}(y)=(a \rightarrow x) \rightarrow(a \rightarrow y)$. At the same time, by Lemma $1(5),(a \rightarrow x) \rightarrow$ $(a \rightarrow y)=((a \rightarrow x) \odot a) \rightarrow y=(a \wedge x) \rightarrow y=(a \odot x) \rightarrow y=x \rightarrow(a \rightarrow y)$, whence the assertion follows.

From Theorem 7 and Proposition 15 we obtain as an immediate consequence the following claim.

Corollary 16. Let $M$ be an $R \ell$-monoid and $a \in I(M)$. Then $\psi_{a}$ is a modal operator on $M$ if and only if for any $x, y \in M$

$$
\psi_{a}(x) \odot \psi_{a}(y) \geq \psi_{a}(x \odot y)
$$

Lemma 17. If $M$ is an $R \ell$-monoid and $a \in M$, then for any $x, y \in M$

$$
x \rightarrow \chi_{a}(y) \leq \chi_{a}(x) \rightarrow \chi_{a}(y)
$$

Proof. By the definition of $\chi_{a}$ and by Lemma $1(5), x \rightarrow \chi_{a}(y)=x \rightarrow$ $((y \rightarrow a) \rightarrow a)=(y \rightarrow a) \rightarrow(x \rightarrow a), \chi_{a}(x) \rightarrow \chi_{a}(y)=((x \rightarrow a) \rightarrow a) \rightarrow$ $((y \rightarrow a) \rightarrow a)$. Since by [14, Łemma 2.3], $(y \rightarrow a) \rightarrow(x \rightarrow a) \leq((x \rightarrow a) \rightarrow a)$ $\rightarrow((y \rightarrow a) \rightarrow a)$, we have $x \rightarrow \chi_{a}(y) \leq \chi_{a}(x) \rightarrow \chi_{a}(y)$.

For any $R \ell$-monoid $M$, let us denote by $B(M)$ the set of all elements from $M$ having the complement in the lattice $(M ; \vee, \wedge, 0,1)$. Note that $0,1 \in B(M)$. If $a \in B(M)$ then its complement $a^{\prime}$ is equal to the element $a^{-}$. By $[9$, Lemma 2.8.8], $B(M) \subseteq I(M)$.

Proposition 18. Let $M$ be an $R \ell$-monoid and $a \in B(M)$. Then for any $x, y \in M$

$$
x \rightarrow \chi_{a}(y)=\chi_{a}(x) \rightarrow \chi_{a}(y)
$$

Proof. Let $a \in B(M), x, y \in M$. Then

$$
\begin{gathered}
x \rightarrow \chi_{a}(y)=x \rightarrow((y \rightarrow a) \rightarrow a)=(y \rightarrow a) \rightarrow(x \rightarrow a), \\
\chi_{a}(x) \rightarrow \chi_{a}(y)=((x \rightarrow a) \rightarrow a) \rightarrow((y \rightarrow a) \rightarrow a) \\
=(y \rightarrow a) \rightarrow(((x \rightarrow a) \rightarrow a) \rightarrow a), \\
x \rightarrow a=x \rightarrow a^{--}=\left(x \odot a^{-}\right)^{-}, \\
(x \rightarrow a) \rightarrow a=\left((x \rightarrow a) \odot a^{-}\right)^{-}=\left(\left(x \odot a^{-}\right)^{-} \odot a^{-}\right)^{-}=\left(\left(x \wedge a^{-}\right)^{-} \wedge a^{-}\right)^{-} \\
=\left(\left(x \wedge a^{-}\right) \vee a\right)^{--}=\left((x \vee a) \wedge\left(a^{-} \vee a\right)\right)^{--}=(x \vee a)^{--}-x \oplus a, \\
((x \rightarrow a) \rightarrow a) \rightarrow a
\end{gathered} \begin{aligned}
&=\left(((x \rightarrow a) \rightarrow a) \odot a^{-}\right)^{-}=\left((x \oplus a) \odot a^{-}\right)^{-}=(x \oplus a) \rightarrow a \\
&=(x \vee a)^{--} \rightarrow a^{--}=((x \vee a) \rightarrow a)^{--}=\left((x \vee a) \cdot a^{-}\right)^{-} \\
&=\left(\left(x \odot a^{-}\right) \vee\left(a \odot a^{-}\right)\right)^{-}=\left(x \odot a^{-}\right)^{-}-x \rightarrow a .
\end{aligned}
$$

Hence

$$
\begin{aligned}
\chi_{a}(x) \rightarrow \chi_{a}(y) & =(y \rightarrow a) \rightarrow(((x \rightarrow a) \rightarrow a) \rightarrow a) \\
& =(y \rightarrow a) \rightarrow(x \rightarrow a)=x \rightarrow \chi_{a}(y) .
\end{aligned}
$$

Corollary 19. Let $M$ be an $R \ell$-monoid and $a \in B(M)$. Then $\chi_{a}$ is a modal operator on $M$ if and only if for any $x, y \in M$

$$
\chi_{a}(x) \odot \chi_{a}(y) \geq \chi_{a}(x \odot y)
$$

Let $M$ be an $R \ell$-monoid and $f$ be a modal operator on $M$. Then $\operatorname{Fix}(f)-$ $\{x \in M: f(x)=x\}$ will denote the set of all fixed elements of the operator $f$. By the definition of a modal operator it is obvious that $\operatorname{Fix}(f) \quad \operatorname{Im}(f)$.

Since $f$ is a closure operator on the lattice $(M ; \vee, \wedge)$, we infer that $\left(\operatorname{Fix}(f) ; \vee_{F}, \wedge\right)$, where $y \vee_{F} z=f(y \vee z)$ and " $\wedge$ " is the restriction of the corresponding operation from $M$ on $\operatorname{Fix}(f)$, is a lattice.
Theorem 20. If $f$ is a modal operator on an $R \ell$-monoid $M$, then $\operatorname{Fix}(f)$ is closed under the operations " $\odot$ " and " $\rightarrow$ " and $\operatorname{Fix}(f)-\left(\operatorname{Fix}(f) ; \cdot, \vee_{F}\right.$, $\wedge, \rightarrow, f(0), 1)$ is an $R \ell$-monoid.

Proof.
(i) Let $x, y \in \operatorname{Fix}(f)$. Then $f(x \odot y)=f(x) \odot f(y)=x \odot y$, thus $x \odot y \in \operatorname{Fix}(f)$.
(ii) $\left(\operatorname{Fix}(f) ; \vee_{F}, \wedge, f(0), 1\right)$ is a bounded lattice.
(iii) If $y, z \in \operatorname{Fix}(f)$, then by Proposition 3 we have $y \rightarrow z=f(y) \rightarrow f(z)=$ $f(f(y) \rightarrow f(z))=f(y \rightarrow z)$, hence $y \rightarrow z \in \operatorname{Fix}(f)$.

Therefore, if $x, y, z \in \operatorname{Fix}(f)$, then $x \odot y, y \rightarrow z \in \operatorname{Fix}(f)$ and for this reason $x \odot y \leq z$ holds in $\operatorname{Fix}(f)$ if and only if $x \leq y \rightarrow z$.
(iv) By foregoing, $\operatorname{Fix}(f)$ also satisfies the identity $x \odot(x \rightarrow y)=x \wedge y$.

Remark 21. The above theorem strengthens general Lemma 3.3 of [5] proved for any residuated lattices in our special case of bounded (commutative) $R \ell$-monoids.

Theorem 22. Let $M$ be an $R \ell$-monoid, $a \in I(M)$ and

$$
I(a):=[0, a]=\{x \in M: 0 \leq x \leq a\}
$$

For any $x, y \in I(a)$ we set $x \odot_{a} y=x \odot y$ and $x \rightarrow_{a} y:=(x \rightarrow y) \wedge a$. Then $I(a)=\left(I(a) ; \odot_{a}, \vee, \wedge, \rightarrow_{a}, 0, a\right)$ is an $R \ell$-monoid.

Proof.
(i) If $x, y \in I(a)$, then $x \odot y \in I(a)$ and $x \odot a=x \wedge a=x$, hence $\left(I(a) ; \odot_{a}, a\right)$ is a commutative monoid.
(ii) Obviously, $(I(a) ; \vee, \wedge, 0, a)$ is a bounded lattice.
(iii) Let $x, y \in I(a)$. It holds that $x \rightarrow y$ is the greatest element $z \in M$ such that $x \odot z \leq y$. Therefore $(x \rightarrow y) \wedge a$ is the greatest element in $I(a)$ with this property. That means, $x \odot_{a} z \leq y$ if and only if $z \leq(x \rightarrow y) \wedge a=x \rightarrow_{a} y$ for every $z \in I(a)$.
(iv) For any $x, y \in I(a)$ we have $x \odot_{a}\left(x \rightarrow_{a} y\right)=x \odot((x \rightarrow y) \wedge a)=x \odot$ $(x \rightarrow y) \odot a=(x \wedge y) \wedge a=x \wedge y$.

Remark 23. If for any $x, y \in I(a)$ we denote by $x^{-a}$ the negation of an element $x$ and by $x \oplus_{a} y$ the sum of elements $x$ and $y$ in the $R \ell$-monoid $I(a)$, then it holds

$$
x^{-a}=x^{-} \wedge a, \quad x \oplus_{a} y=(x \oplus y) \wedge a
$$

Indeed

$$
\begin{aligned}
x^{-a} & =x \rightarrow a 0=(x \rightarrow 0) \wedge a=x^{-} \wedge a \\
x \oplus_{a} y & =\left(x^{-a} \odot y^{-a}\right)^{-} \wedge a=\left(x^{-} \odot a \odot y^{-} \odot a\right)^{-} \wedge a=\left(x^{-} \odot y^{-} \odot a\right)^{-} \wedge a \\
& =\left(a \rightarrow\left(x^{-} \odot y^{-}\right)^{-}\right) \odot a=a \wedge\left(x^{-} \odot y^{-}\right)^{-}=a \wedge(x \oplus y)
\end{aligned}
$$

Now, let $M$ be arbitrary $R \ell$-monoid (still bounded and commutative), $a \in I(M)$ and let $f$ be a modal operator on $M$. Let us consider a mapping $f^{a}: I(a) \longrightarrow I(a)$ such that $f^{a}(x)=f(x) \wedge a(=f(x) \odot a)$, for every $x \in I(a)$.

Theorem 24. Let $M$ be an $R \ell$-monoid, $a \in I(M)$ and $f$ be a modal and strong modal, respectively, operator on $M$. Then $f^{a}$ is a modal and strong modal, respectively, operator on the $R \ell$-monoid $I(a)$.

Proof. Assume $x, y \in I(a)$.

1. $x \leq a$ and $x \leq f(x)$, hence $x \leq a \wedge f(x)=f^{a}(x)$.
2. $f^{a}\left(f^{a}(x)\right)=f(f(x) \wedge a) \wedge a=f(f(x) \odot a) \wedge a=(f(f(x)) \odot f(a)) \wedge a$

$$
=f(x) \wedge f(a) \wedge a=f(x) \wedge a=f^{a}(x)
$$

3. $\quad f^{a}(x \odot y)=f(x \odot y) \wedge a=f(x) \odot f(y) \odot a \odot a=(f(x) \wedge a) \odot(f(y) \wedge a)$ $=f^{a}(x) \odot f^{a}(y)$.
4. Let $f$ be strong. Then

$$
\begin{aligned}
f^{a}\left(x \oplus_{a} f^{a}(y)\right) & =f^{a}((x \oplus(f(y) \wedge a)) \wedge a)=f((x \oplus(f(y) \wedge a)) \wedge a) \wedge a \\
& =f(x \oplus(f(y) \wedge a)) \wedge f(a) \wedge a=f(x \oplus f(f(y) \wedge a)) \wedge a \\
& =f(x \oplus((f(f(y)) \wedge f(a))) \wedge a=f(x \oplus(f(y) \wedge f(a))) \wedge a \\
& =f(x \oplus f(y \wedge a)) \wedge a=f(x \oplus f(y)) \wedge a=f(x \oplus y) \wedge a \\
& =f^{a}(x \oplus y)
\end{aligned}
$$

## Theorem 25.

a) Let $M$ be an $R \ell$-monoid, let $f$ be a modal operator on $M$ and $\hat{f}=\left.f\right|_{I(M)}$. Then $I(M)$ is a subalgebra of the reduct $(M ; \odot, \vee, \wedge, 0,1)$ and $\hat{f}$ is a mapping of $I(M)$ into $I(M)$ satisfying conditions $1,2,3$ from the definition of a modal operator.
b) Let $M$ be a normal $R \ell$-monoid and let $x^{-} \in I(M)$ for each $x \in I(M)$. Then $I(M)$ is closed also under the operation " $\oplus$ ". Moreover, if $f$ is a strong modal operator on $M$, then $\hat{f}$ satisfies condition 4 from the definition of a strong modal operator.
c) Let $M$ be a BL-algebra. Then $I(M)$ is a subalgebra of the algebra $M$ which is a Heyting algebra. If $f$ is a modal operator on $M$, then $\hat{f}$ is a modal operator on the Heyting algebra $I(M)$. If $x^{-} \in I(M)$ holds for each $x \in I(M)$ and $f$ is a strong modal operator on $M$, then $\hat{f}$ is a strong modal operator on $I(M)$.

Proof.
a) Let $M$ be an $R \ell$-monoid and $x, y \in I(M)$. Then

$$
(x \odot y) \odot(x \odot y)=(x \odot x) \odot(y \odot y)=x \odot y
$$

thus $x \odot y=x \wedge y \in I(M)$. Further, $(x \vee y) \odot(x \vee y)=(x \odot x) \vee(y \odot x) \vee(x \odot y) \vee(y \odot y)=x \vee y \vee(x \odot y)=x \vee y$, therefore also $x \vee y \in I(M)$.

Obviously, $0,1 \in I(M)$.
Finally, if $f$ is a modal operator on $M$, then for each $x \in I(M)$ we have

$$
f(x)=f(x \odot x)=f(x) \odot f(x)
$$

It follows that $f(x) \in I(M)$. Therefore $\hat{f}$ is a mapping of $I(M)$ into $I(M)$ satisfying conditions $1-3$.
b) If $x^{-} \in I(M)$ holds for every $x \in I(M)$, then (similarly to the third part of the proof of Theorem 11) for any $x, y \in I(M)$ we obtain $x \oplus y=(x \vee y)^{--}$, and hence provided $M$ is normal we have

$$
\begin{aligned}
(x \oplus y) \odot(x \oplus y) & =(x \vee y)^{--} \odot(x \vee y)^{--}=((x \vee y) \odot(x \vee y))^{--}=(x \vee y)^{--} \\
& =x \oplus y
\end{aligned}
$$

therefore $x \oplus y \in I(M)$.
At the same time it is obvious that if $f$ is a strong modal operator on $M$, then $\hat{f}$ fulfills condition 4 as well.
c) By [13], an $R \ell$-monoid $M$ is a $B L$-algebra if and only if $M$ is isomorphic to a subdirect product of $R \ell$-chains ( $=B L$-chains). Let now a $B L$-algebra $M$ be a subdirect product of $B L$-chains $M_{\alpha}, \alpha \in \Gamma$. If $a \in M$, then $a=\left(a_{\alpha} ; \alpha \in \Gamma\right) \in$ $I(M)$ if and only if $a_{\alpha} \in I\left(M_{\alpha}\right)$ for each $\alpha \in \Gamma$. Let $x=\left(x_{\alpha} ; \alpha \in \Gamma\right), y=$ $\left(y_{\alpha} ; \alpha \in \Gamma\right) \in I(M)$. Then $x_{\alpha} \rightarrow y_{\alpha}=1$ for $y_{\alpha} \geq x_{\alpha}$ and $x_{\alpha} \rightarrow y_{\alpha}=y_{\alpha}$ for $x_{\alpha}>y_{\alpha}$. Whence $\left(x_{\alpha} \rightarrow y_{\alpha} ; \alpha \in \Gamma\right) \in I(M)$ and it is equal to the element $x \rightarrow y$. By [13], furthermore, $I(M)$ is a Heyting algebra.

Then it is clear that $\hat{f}$ is a modal operator on $I(M)$ for any modal operator $f$ on $M$. Moreover, by [15, Proposition 5], every $B L$-algebra is a normal $R \ell$-monoid. Therefore, if $x^{-} \in I(M)$ for each $x \in I(M)$, then $\hat{f}$ is a strong modal operator on the Heyting algebra $I(M)$ for every strong modal operator $f$ on $M$.

Remark 26. For any $a \in M$, also mappings $\pi_{a}: M \longrightarrow M$ (in our notation) defined by $\pi_{a}(x)=a \vee x$ for each $x \in M$ were introduced and studied for Heyting algebras in [10]. Evidently, if $M$ is an arbitrary $R \ell$-monoid, then $\pi_{a}$ satisfies conditions 1 and 2 from the definition of a modal operator on $M$. This begs the question if $\pi_{a}$ fulfills condition 3 from this definition as well and in which cases $\pi_{a}=\varphi_{a}$ holds, respectively.
a) If $M$ is a Heyting algebra then $x \odot y=x \wedge y$ for any $x, y \in M$. From the distributivity of the lattice $(M ; \vee, \wedge)$ it follows that condition 3 is satisfied for any $a \in M$. At the same time, $a \oplus x=(a \vee x)^{--}$, hence $\pi_{a}$ need not generally be equal to $\varphi_{a}$. For example, $\pi_{0}(x)=x, \varphi_{0}(x)=x^{--}$.
b) If $M$ is an $M V$-algebra, then $a \vee x=a \oplus x$ holds for any $a \in I(M)$ and $x \in M$, and $a \oplus(x \odot y)=(a \oplus x) \odot(a \oplus y)$. Therefore, we have $\varphi_{a}=\pi_{a}$ for each $a \in I(M)$ and hence, for each $a \in I(M)$, moreover $\pi_{a}$ is a strong modal operator on $M$.

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